

signals & scheduling.

recap: fork, exec, copy-on-write.

(i) signals

- OS-assisted / s/w defined mechanism to deliver events / interrupts / control signals to processes.
- a signal / pending signals are part of OS state!
- signals are processed before a process is scheduled on the CPU.

types of signals.

+ override or not

+ type of default action

+ terminate

+ dump core (memory contents of process to a file)

(*) signal handling

+ default OS action

+ user-mode handler per signal

(via registration)

(*) signal delivery

depends on the system calls.

SIG KILL

9

terminate cannot override

SIG INT

2

interrupt

terminate

SIG SEGV

15 11

override

override, dump core

SIG STOP

19

cannot override

change process state to blocking

SIG CONT

18

cannot override

change status

SIG CHLD

can override.

(signal on child process terminate)

ignore

sigchild

```
void handlechildexit() {
```

```
    pid = waitpid(-1, &status, WNOHANG);
```

```
}
```

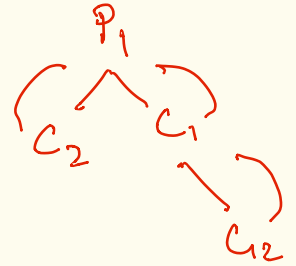
```
main() { signal(C handlechildexit, SIGCHLD);
```

```
    fork();
```

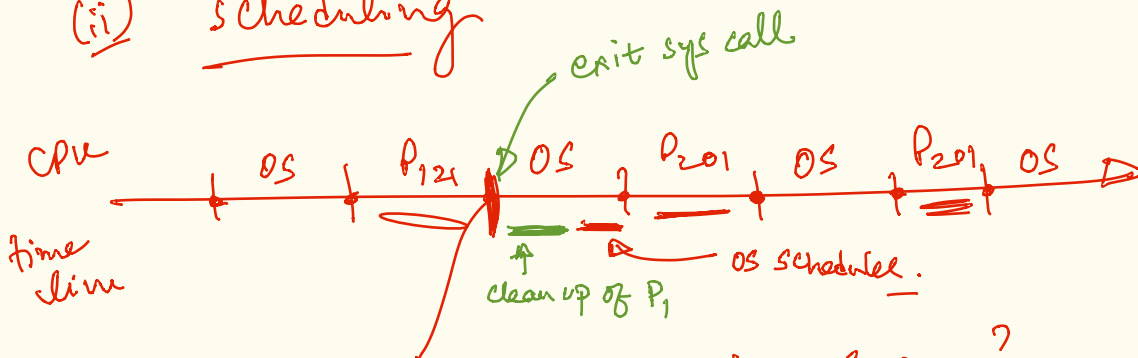
```
    fork();
```

```
    // long running code
```

```
}
```



(ii) scheduling



- when does this transition happen?
- how does this transition happen?
- what does the OS do in its time on the CPU?

⊗ list of actions that a scheduler has to do to switch from P_1 to P_2 ...

(i) an event that triggers the scheduler — interrupt system call

(ii) ~~context switch.~~ context switch.

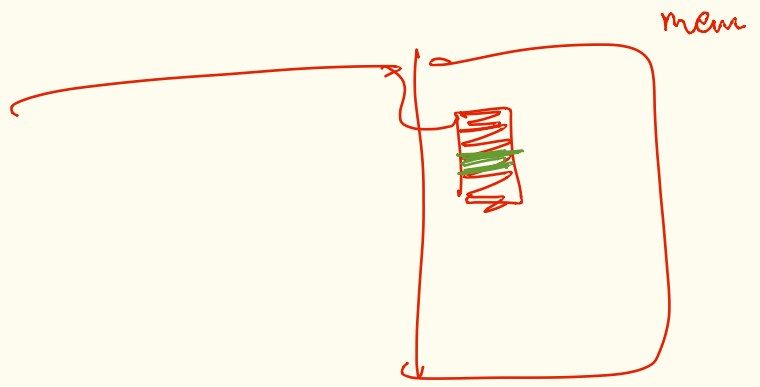
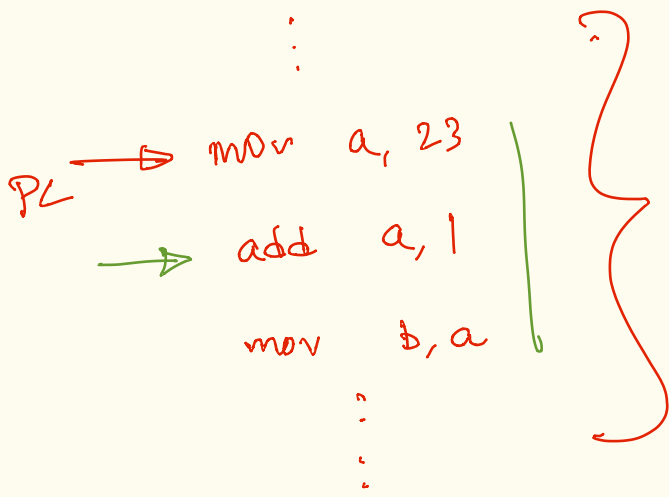
↓ save context of outgoing process.
↓ restore context of incoming process.

scheduling

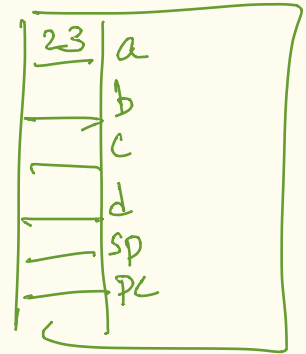
policy

(ii)

choose the next process to execute on CPU.

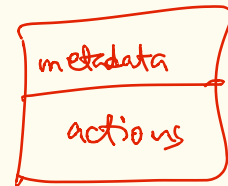


⊗ Context of a process ⇒ all CPU state (registers).



(iii) Scheduling mechanism

- ⊗ ready queue is the set of all schedulable processes.
- ⊗ multi-CPU setup



single ready queue
 needs synchronization
 → per CPU ready queue