

Lecture 10 — system calls & memory mgmt.

system call mechanism.

- ✓ + how to invoke?
- ✓ + how to switch mode of execution?
- ✓ + how to handle arguments & return values?
- ✓ + how to identify which system call?
- ✓ + how to handle context of execution?

* Lab 3 available
xv6-based
system calls,
process abstraction
memory abstraction

⊗ xv6-book |
reference

(i) How assisted mechanism for invocation
is an instruction (s) of an ISA.

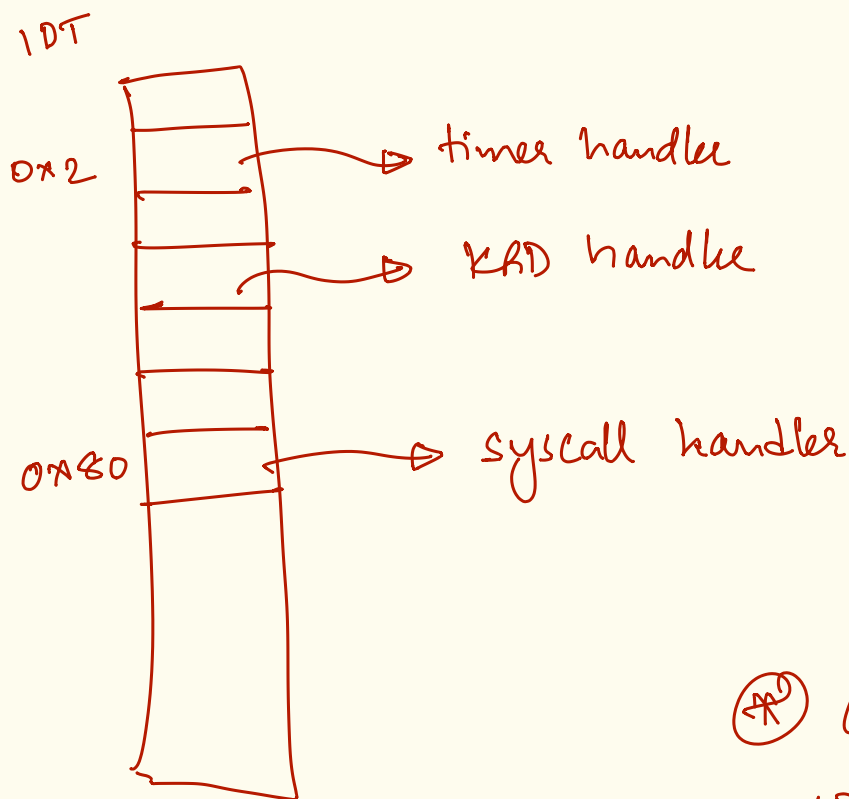
x86: int 0x80 ~ explicit s/w interrupt

- instruction argument - interrupt vector
- + s/w privilege level to Ring 0
 - ⊗ + save context of user process on kernel stack
 - + s/w to using kernel stack
 - + jump to interrupt handler. ~ depends on IDTR & interrupt vector

(ii) how to jump to the interrupt handler?

IDT — interrupt descriptor table

— array of function pointers (interrupt handling functions)



(*) how to get IDT?

IDTR
↓
register to store base address of the IDT.

(*) OS maintains IDT & IDTR.

(iii) how to identify the system call and its arguments?

— via syscall number

— before ~~invoking~~ invoking system call setup syscall number & arguments in CPU regs.

eax ← syscall number

ebx
cx
dx
⋮

} arguments.

int 0x80

} user-space wrapper fⁿ.

sets this up this invocation

Open: {
 mov eax, SYS_OPEN
 {mov ebx
 ecx
 :
 int 0x80

(iv) system call handler.

syscall - handler {

fn ptr = syscall-table [^{value of} eax];
~~from~~ on kernel stack

call fn ptr;

iret; — return from interrupt.

} — instruction of the ISA.

↓
 return values are via register (eax).

⑧ xv6 mechanism details.

traps.h — list of all interrupts

vectors.S — IDT entries for each interrupt vector

trapasm.S ~ alltraps & label/address entry point for all interrupts

trap.C — generic IDT handler

syscall.c — syscall handler

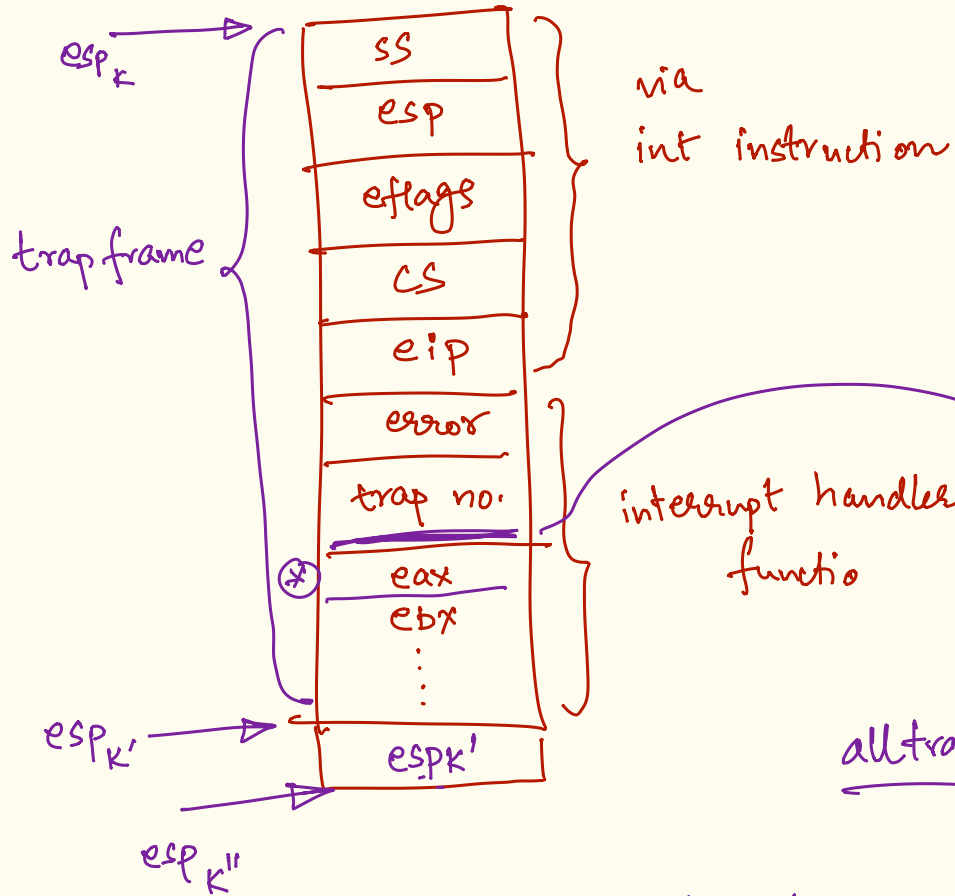
xv6 (x86) syscall / interrupt invocation state

kernel stack

Kernel
 esp_k — stack ptr.

info. is
stored in

task segment descriptor.



trap (trapframe * tf) { push esp @ * call trap