

- Lab2 is available

- Lab exam 1 — 28th Aug.
Wed.

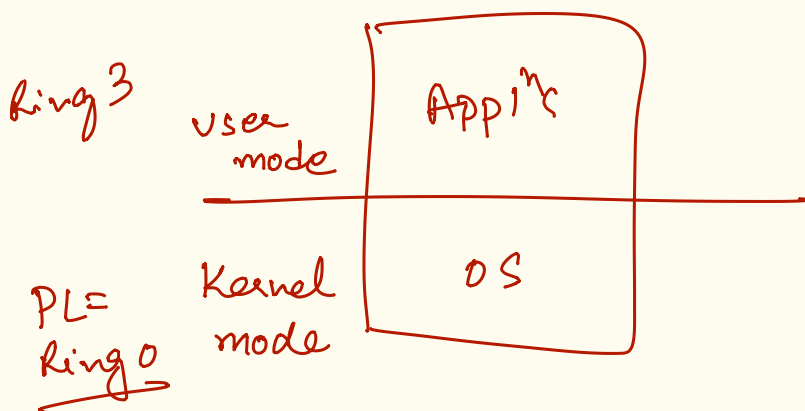
8:30 am
— 11 am

SL2

two main building blocks
of OS design

(i) privileged levels of execution
via the ISA.

- min. PL per instruction
for correct execution.



(ii) interrupts

world is non-deterministic

IO is non-deterministic / non-scheduled.

interrupts

- hardware — IO devices / device controllers.
- software

(*) interrupt delivery

stopping of current seq. of instr. on CPU

switch to kernel mode

handle interrupt

schedule process/work

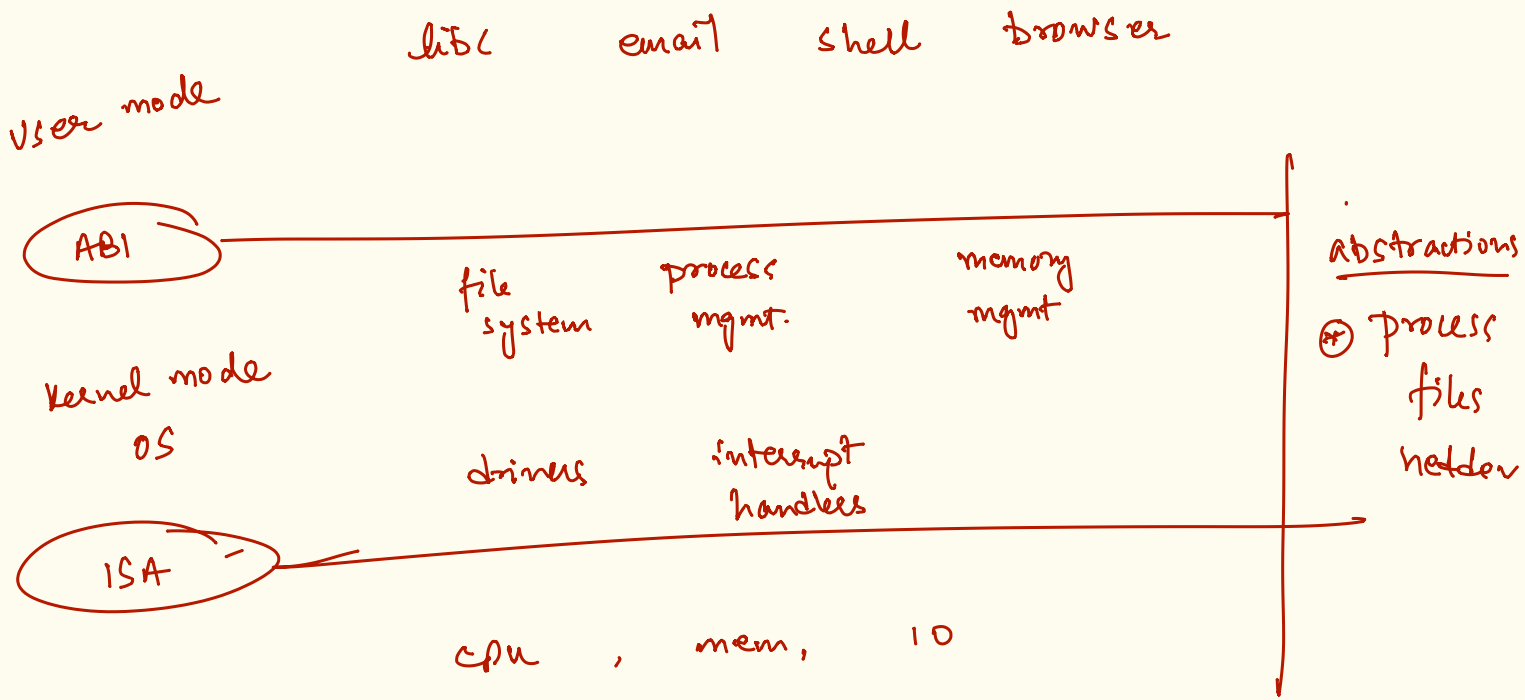
software

implicit — exceptions
div. by zero
seg fault

explicit — system calls

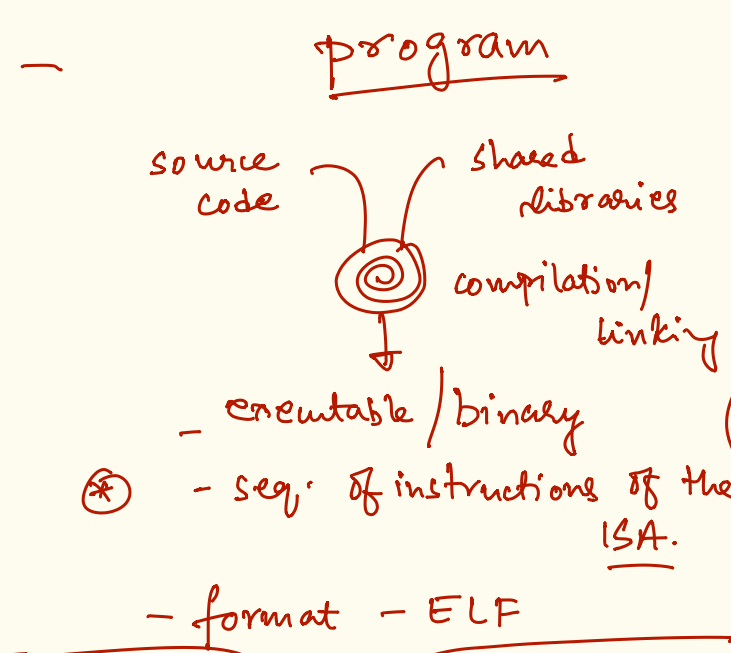
eg: x86

int 0x80



the process vs abstraction.

(user-mode)



- program in execution
- instance of a program
- entity that can consume resources / OS functionality
- ~ program is set of inst^r.
- loaded in memory & pointed to by the PC.

(Q) how to invoke the stored-program (von Neumann model)

- where to load program in memory?
- when to execute on CPU?
- how long to execute on CPU?
- how to share the CPU?

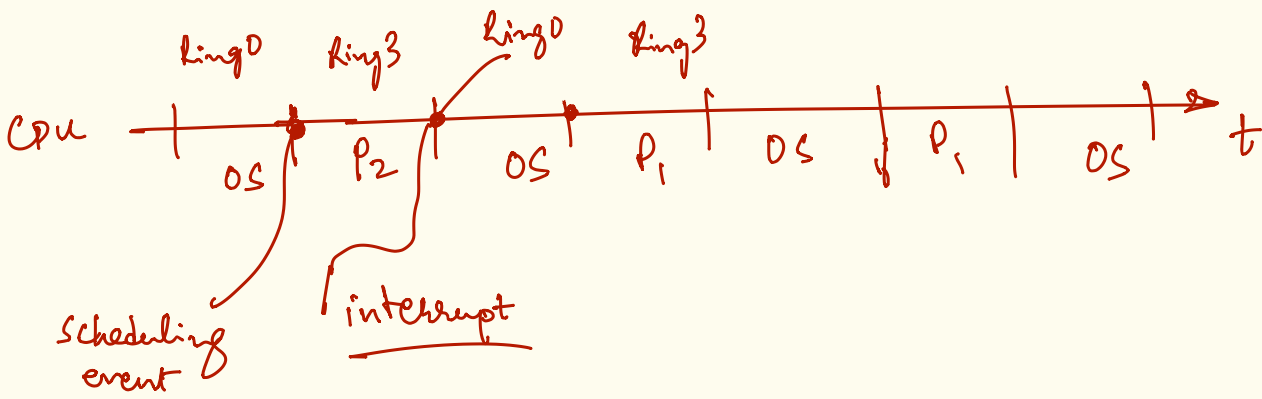
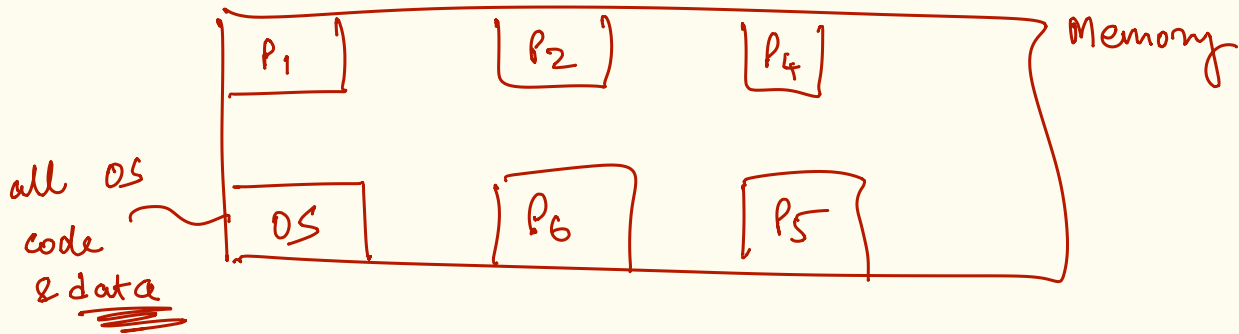
~ process abstraction decouples program development and its execution.

resource allocation, setup, scheduling, termination

⑧ for every abstraction OS provides, the OS stores meta-data / information regarding the abstraction's implementation.

e.g: cab-hailing service

└ # cabs, # drivers, locations of cabs, traffic, billing...



⑧ metadata of process. / PCB — process control block.

- PC, state
- pid, Ppid
- context ~~~~~ state / registers of the CPU due to instructions of the process.
- open files
- memory allocation regions
- signals

Q1) when a program is loaded / process is created will it always be executing?

- NO!

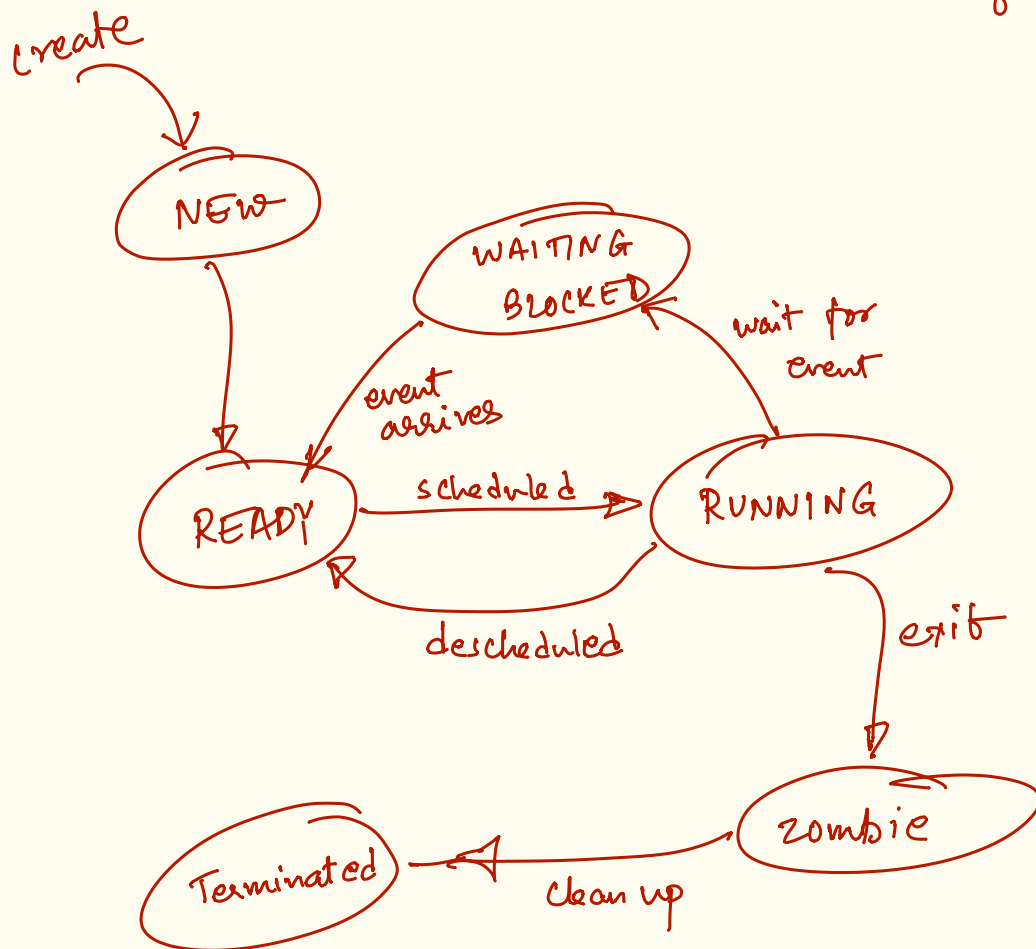
- waiting for its turn (scheduling decision)

- waiting for an event (IO completion, lock to be free)
(blocked)

- done execution

- being setup

Q2) process state that captures current execution state of a process.



Q3) fork, exec, wait, waitpid