

Persistence

the file abstraction

- name, organization (PATH)
- size, timestamps
- links
- access control information
- linear ~~be~~ sequence of bytes.

Q how to design of the file system?

requirements

the file abstraction implementation

reliability

efficiency!

deal w/ diverse storage devices.

CPU — faster than memory (RAM) access

faster than storage / disk

① Quiz 2 (pen-paper)

29th Oct. Wed

830 am

② Lab Quiz

5th Nov. Wed.

830 am

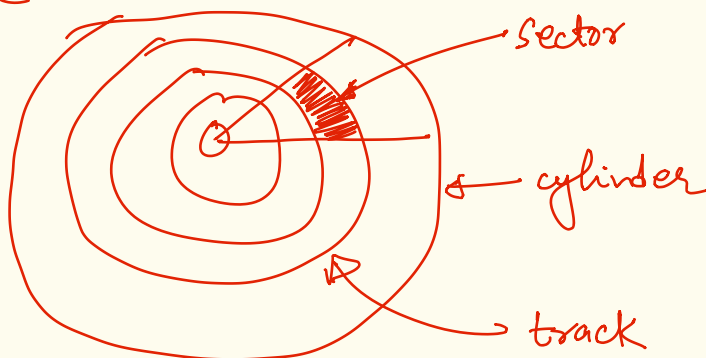
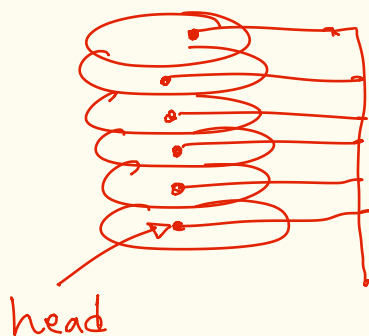
- implementing / coding

- Lab 1/2/3, Practice Labs

- SL 2

① what is n/w view?

(a) HDD — Hard disk drive



* 1500 rpm

each access $\Rightarrow$

C, t, S

— cylinder
track
sector

(*) ~ moving parts (electro-mechanical)

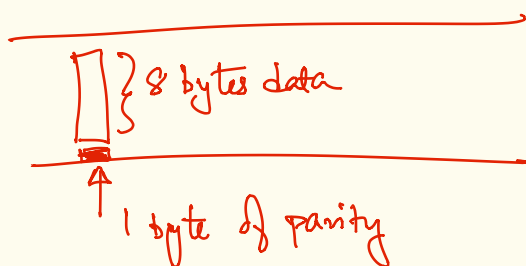
~ magnetization on the cylinder to encode 0s & 1s

(b) Flash drives / USB disk

- NOR/NAND flashes - / gates
- semi-conductor based - / no-moving parts
- banks & pages of physical storage layout
- access $\Rightarrow$ bank index & page number
r/w is at page granularity
- fixed # writes per storage location.

(c) Tape drives

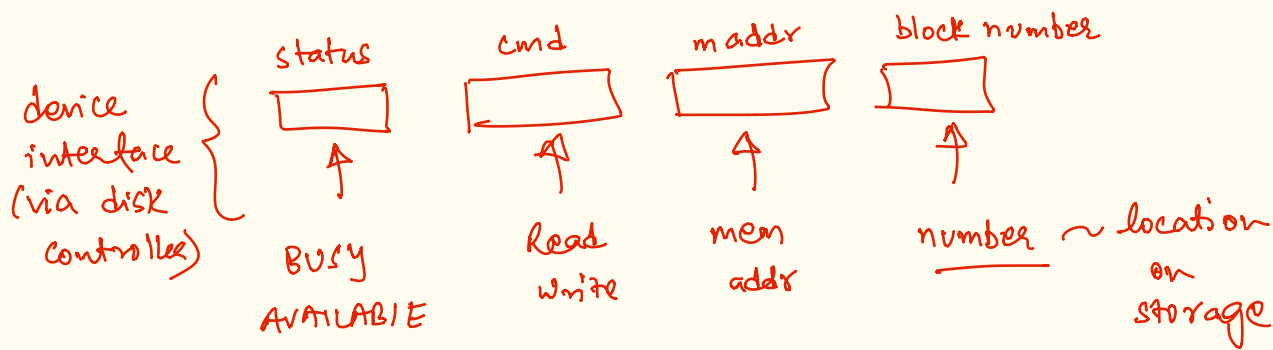
access $\leftarrow$ offset!



Disk / Storage controller - h/w element.

- └ provides a higher / uniform abstraction
e.g.: block
- a storage device / disk is a ordered set
of of blocks
- each block is ~~a~~ fixed size.

canonical device interface

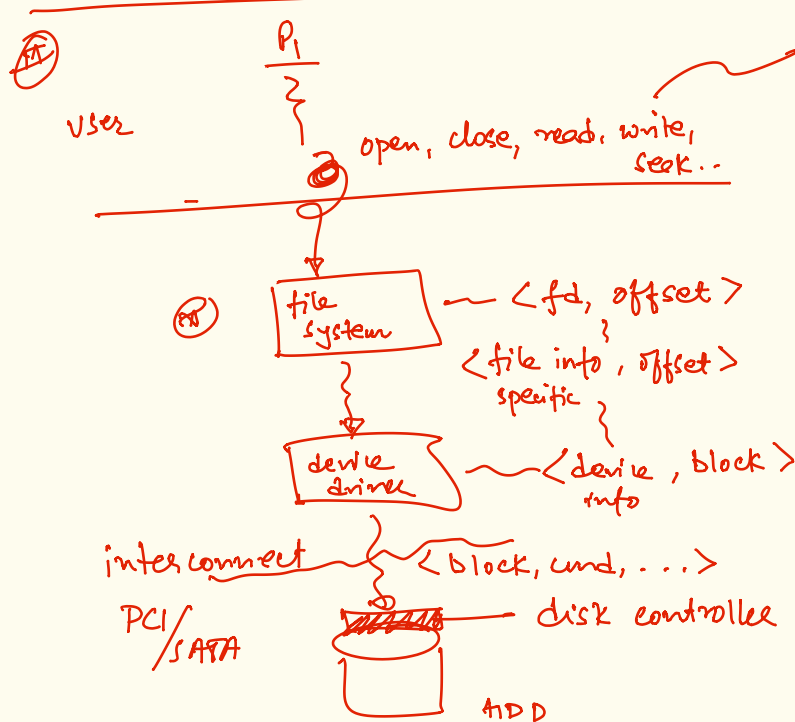


device driver!

```

while (STATUS == BUSY);
cmd ← R;
maddr ← &buffer;
block ← 23;
STATUS ← BUSY;
    
```

cmds for disk controller to work on

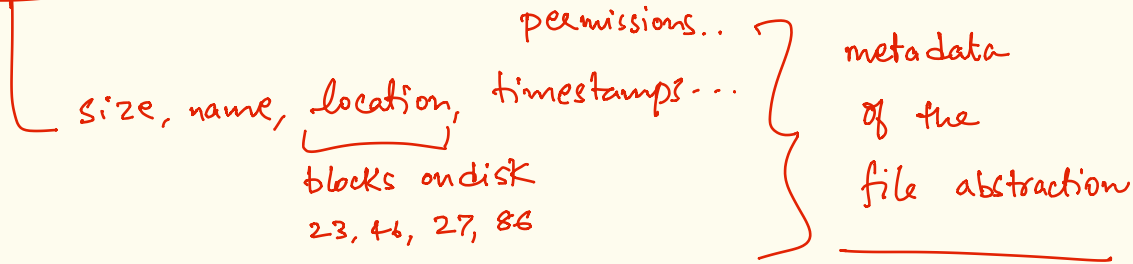


filename
fd - file descriptor

```

fd = open("/home/foo.txt", R);
read(fd, buf, size);
read(fd, buf, size);
    
```

file system design



Q where is the inode stored?

inode ← index node
metadata object for each file.

① has to be a persistent storage medium
on disk!

② disk metadata ~~~~~ status of disk usage
on disk

disk (linear order of blocks)

