

the system call mechanism

- abstraction for OS services
- the system call interface
list/set of all fⁿs/entry points
to invoke service

- fork, exec, wait, exit, read, write,
open, close,

lseek

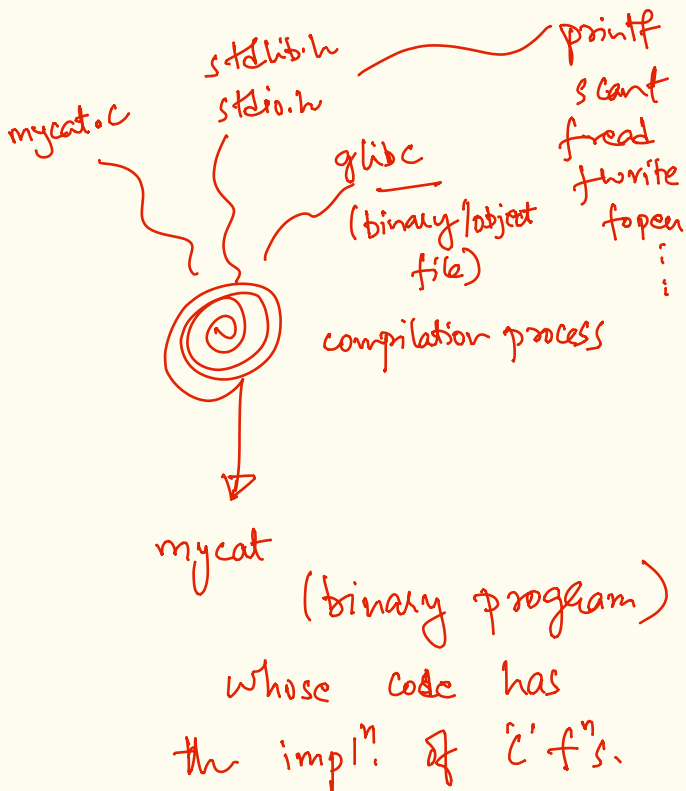
sleep

getpid

getppid

pipe

dup

~~printf~~

① Lab 2

② mycat.c
mycat2.cmycat3.cwhich executes as
a child process
of mycat2mycat2 waits till
mycat3 is done.mycat4.c (extend mycat3)
which writes output
to a file instead
of to terminal.

③ Write a program (P1)

- that overwrites itself.
w/ program (P2).

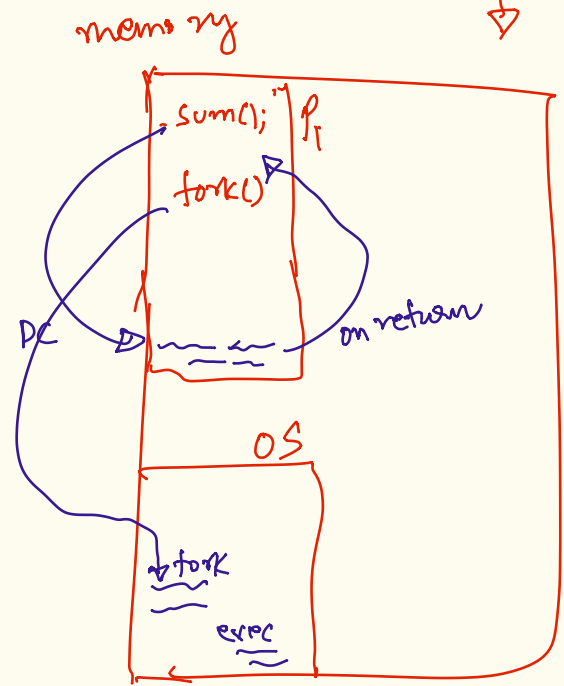
- Does anything after
exec execute within
the calling program?

(*)	API	vs	ABI
	application programming interface		appl ⁿ . binary interface
	- source code level interaction		~ runtime-used interface ~ interaction between compiled/ binary programs

P1:

```
int a=23;
int b=57;
sum(a,b);
fork();
:
:
```

```
int sum(int a, int b)
return (a+b);
```



Qs for the system call mechanism design

~ how to invoke a system call?

~ mechanism to switch privilege levels?

~ mechanism to save & restore context

CPU state

return addr.

return value

~ -- to handle/pass arguments

~ -- to specify which system call

& use to identify locⁿ in OS program.

x86 instruction

↓ jmp <memory addr>

PC is set to mem. addr.

how to invoke a system call?

explicit interrupt

interrupts

hardware → external IO

software { implicit { div by zero
null ptr access

explicit { system call

x86 ISA

(execute an instruction explicitly)

int 0x80 → interrupt vector
instruction to generate an interrupt

- slow CPU to highest PL

- saves context of running process { temporary buffer
(all regs. of CPU) } kernel stack

- switches to kernel stack

- jumps to interrupt handler function

iret → return from interrupt | instruction of the ISA

- restores context of CPU regs.

~ switches to user stack

~ changes PL to user-mode

~ jumps to user-space return address

interrupt handler / isr — interrupt service routine

IDT location is stored in a CPU register

IDTR ~ register storing address of IDT.

