

spinlocks,
diable interrupts
diable pre-emption

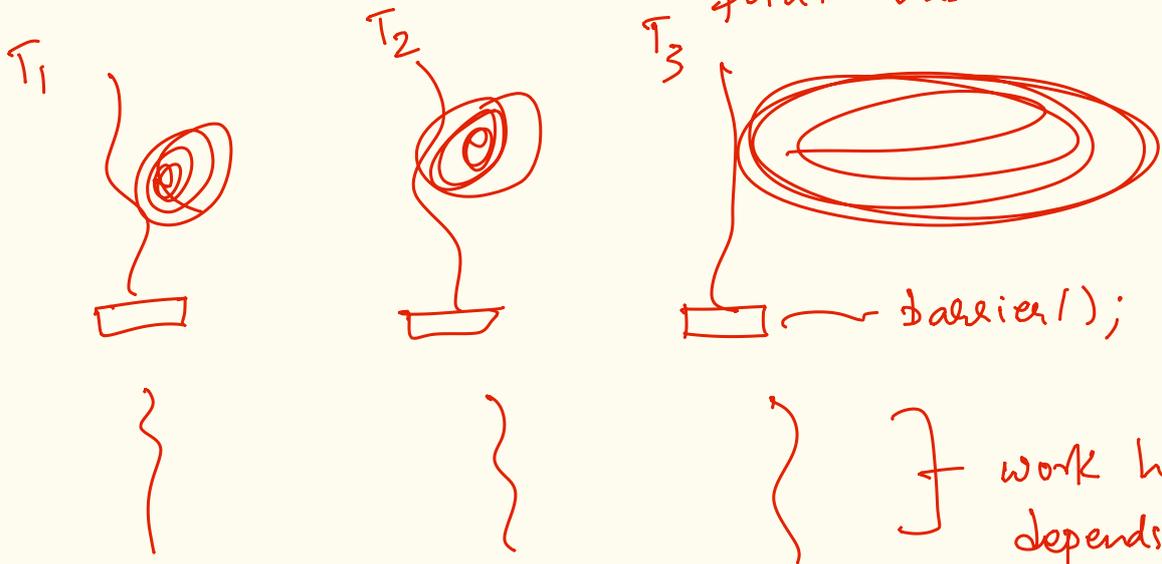
lock;
spinlock
mutex

condition variables
semaphores

barrier

primitive to hold all threads
till a condition/execution

point has reached.



⊛ implement a barrier.

init ————— how many instances
 barrier ————— should reach the barrier!
 ————— wait for barrier condition.

⊛ barrier using semaphores.

init
down
up

Semaphore S_1, S_2

$S_1.init(K);$

$S_2.init(0);$

barrier() {

down(S_1);

if ($S_1 == 0$) $S_2.init(K);$

down(S_2); ← sleep/block till

all threads

reach barrier;

a reset of semaphore will also check for wakeup condⁿ.

}

H.W

correctness of nested barriers.

barrier(B);



barrier(B);

barrier(B);



barrier(B);

barrier(B);



barrier(B);

1. verify that above implementation has/does not have a problem.

2. if problem, what solution?

if no problem, really?

assume.

$S_2.init(K)$

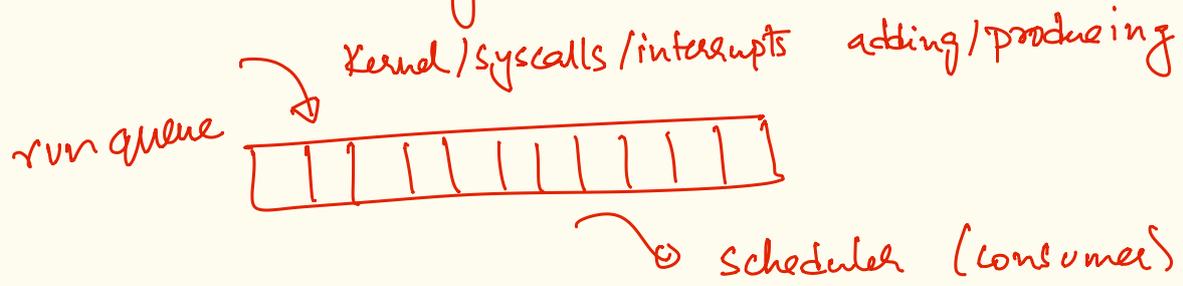
$S_1.init(K)$

as part

of the if ($S_1 == 0$)

condition.

producer-consumer synchronization problem.



~ producers add to queue till size of queue

~ consumers consume/dequeue from queue if queue not empty. } atomic synchronization

implement.: MAX. is size of queue.

using semaphores

a solution of the prod-con synchronization problem:

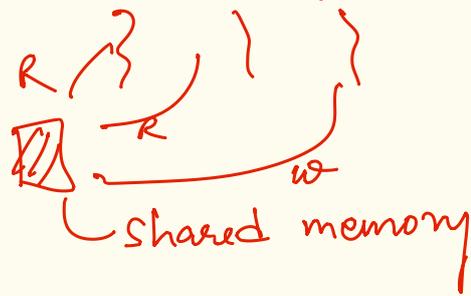
Semaphores $S_p(MAX); S_c(0);$

```
producer() {  
    down(Sp);  
    spinlock(L);  
    K = (index++ % MAX);  
    spinunlock(L);  
    addqueue(K);  
    up(Sc);  
}
```

```
consumer() {  
    down(Sc);  
    spinlock(L);  
    index = (index - 1) % MAX;  
    spinunlock(L);  
    dequeue(index);  
    up(Sp);  
}
```

or use a pindex & cindex

reader-writer synchronization problem.



~ if reader; allow other readers
; do not allow other writers

if writer; do not allow readers
or other writers.

H.W2 implement RW lock using condition variables.

Reader Lock

Writer Lock

Reader Unlock

Writer Unlock