

* interrupt handling (with VMs & VMMs)

- interrupt traps

external / I/O interrupt → phys. signaling to the CPU
 exception
 explicit interrupt

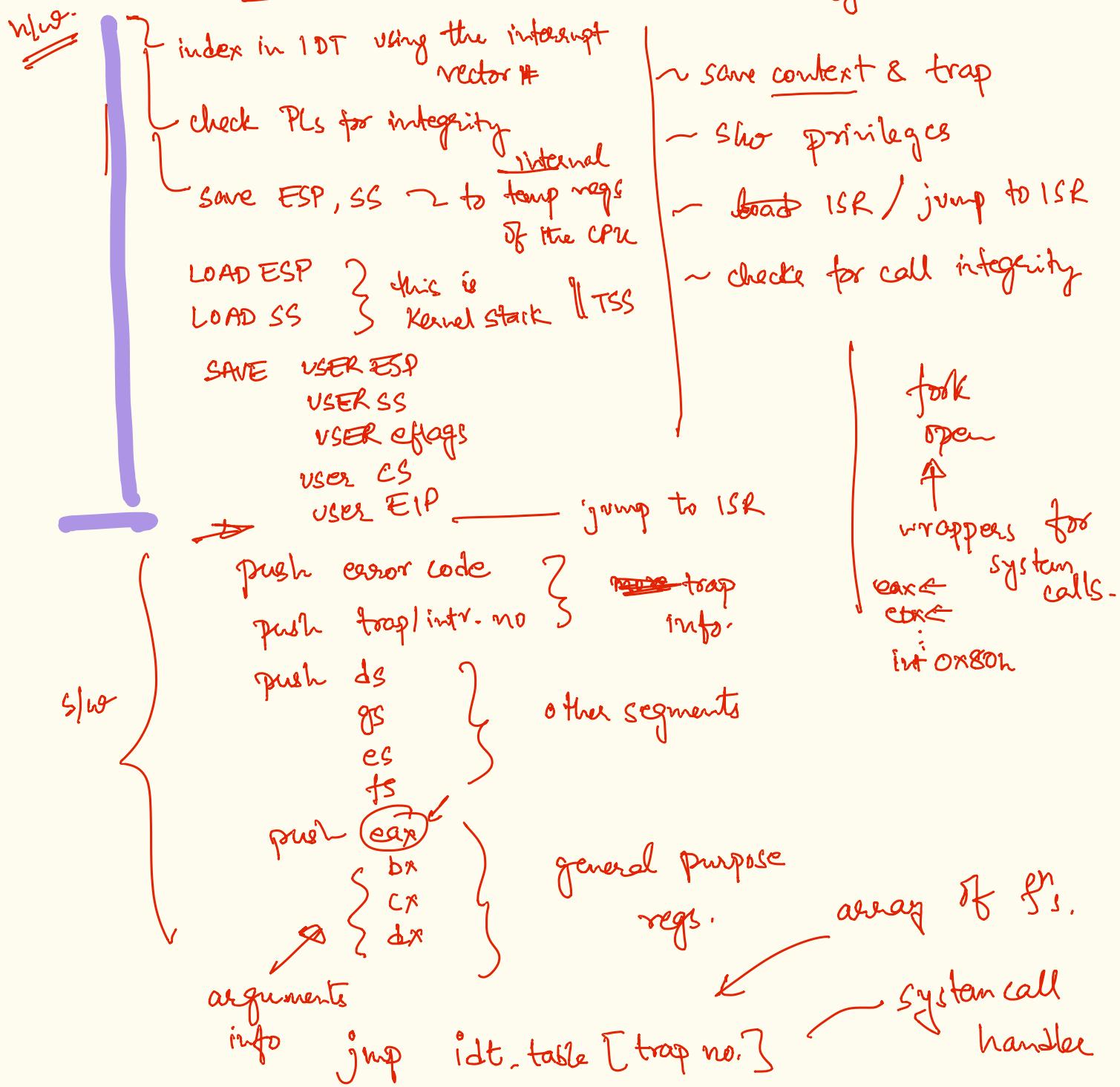
S/W →
 PG·filt., null deref.

* Native systems int processing

→ int 0x80

→ user space

int
 system call

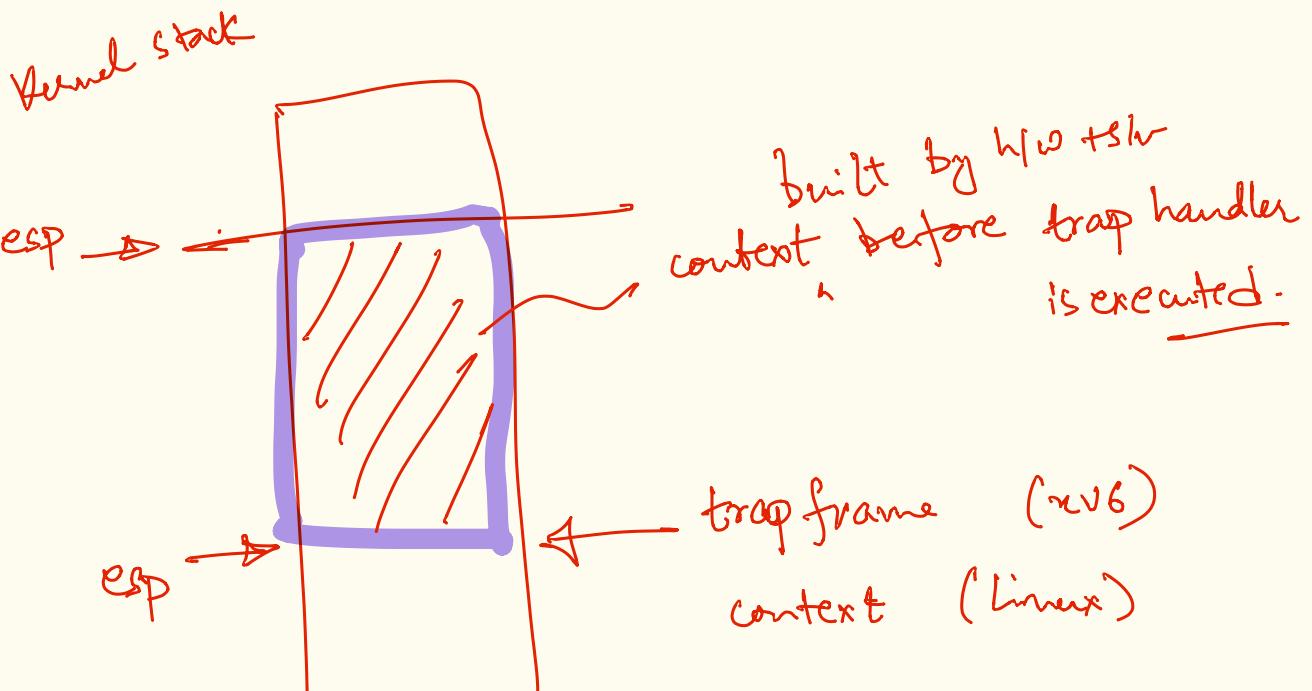


```

system-call-handler( ) {
    fn = sys-call-table [eax];
    call fn;
    iret;
}

```

Sys call no.



return path in the ISR

- pop al

pop ds
gs

es

fs

adjust ESP ~ to move over error code
↓ trap no.

iret

 } ~ hw-assisted instruction.

 rolls back ESP B Kernel

 } loads user reg's.

 } jumps to CS:ESP w/ PL=3

(2) what happens to system calls from VMs?

(i) full virtualized VMs. ~ binary translation / HVM

- on interrupt

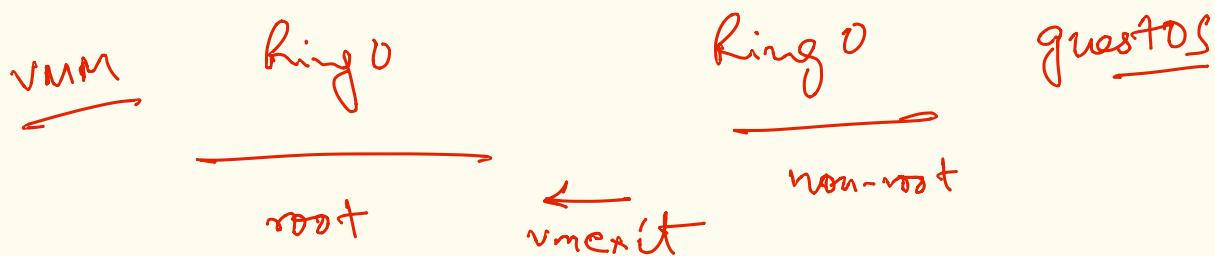
{ control in VMM
context on VMM stack
ISR uses context to determine that this int. should be handled by guest.
emulate an interrupt

→ control to VMM
VMM needs to pass interrupt to guest
guest executes handlee
~ int in guest trap to VMM

{ copying context from VMM to guest Kernel stack? ~ update to ESP & SS
show to guest at CS: EIP if guest ISR.
→ VMM returns to guest user space
trap to VMM

⊕ Similar in PV, but via hypercalls for critical stati.

interrupts with HW assisted VMs.



~ depends on exit conditions & IDT setup.

Assignment #3

~ build a simple VM with KVM
use

~ qemu + KVM

(i) User process, emulator

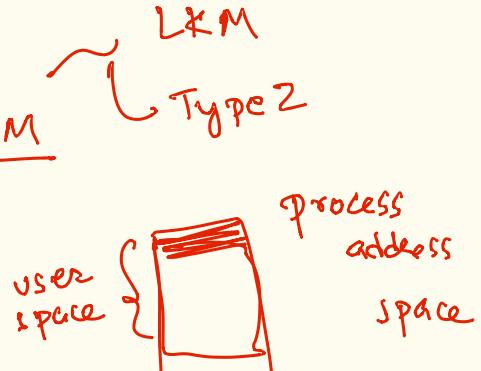
(ii) one soln. will look like qemu $gpa \rightarrow gpa$ mpa
device file $hva \rightarrow$

(iii) /dev/KVM

$fd = open(/dev/KVM);$

$ioctl(fd, \underline{params})$

action arguments



ioctl ~ generic
system call.