

## Self-assessment Quiz

Total Marks: 80

**Note.** You are supposed to solve these problems without any discussions with anyone. Try to write your answers succinctly.

**Que 1.** [20] A graph is called bipartite if its vertex set can be partitioned into two parts, say  $V_1$  and  $V_2$ , such that all the edges connect a vertex from  $V_1$  to a vertex from  $V_2$ . Let us say we are given a graph by its adjacency matrix. Can you design an algorithm which can test whether the graph is bipartite? Formally argue that your algorithm is correct.

*Hint:* Start by assuming that the graph is indeed bipartite with  $V_1, V_2$  as the two parts. Pick an arbitrary vertex  $v$ , which we can assume belongs to  $V_1$  (why?). Now, any neighbor of  $v$  should belong to  $V_2$ . Similarly, for any vertex  $u \in V_2$ , any neighbor of  $u$  must belong to  $V_1$ . This way you can try to mark all the vertices as either belonging to  $V_1$  or  $V_2$ . Now, how would you detect if the graph was not actually bipartite?

You can try to use a depth-first-search or breadth-first-search for exploring the graph.

**Que 2.** [10] Mark true or false.

- $(2n + 1)^2 = O(n^2)$ .
- $f(n) = O(g(n)) \implies 2^{f(n)} = O(2^{g(n)})$ .

Prove that if  $f(n) = O(g(n/2))$  then there is a polynomial function  $g(n)$  such that  $f(n) = O(g(n))$ .

**Que 3.** [10+10] Part(a). Let  $G$  be a directed graph with  $N + 1$  vertices, for a large number  $N$ . Imagine the vertices to be points on the number line. That is, the vertices are indexed  $0, 1, 2, 3, \dots, N$ . For each  $i \geq 2$ , we have two edges going out of the vertex  $i$  – one going into  $i - 1$  and the other into  $i - 2$ . Moreover, there is one edge from 1 to 0. Show that the number of distinct paths from vertex  $N$  to vertex 0 is at least  $2^{\lfloor N/2 \rfloor}$ .

**Hint:** Observe that there are two paths from vertex  $i$  to vertex  $i - 2$ . It suffices to count only a subset of all paths.

Part (b). Consider the following recursive procedure to compute the fibonacci series.

```
function Fibonacci(n):
  if n = 0 or n = 1: return 1;
  else: return Fibonacci(n-1)+Fibonacci(n-2);
```

Using part (a) show that this procedure takes  $\Omega(2^{\lfloor N/2 \rfloor})$  time to compute Fibonacci( $N$ ).

**Que 4.** [10] Prove the following. In a party with 100 people, there must at least two people who have the same number of friends in the party.

**Que 5.** [10] Consider the following algorithm.

```
input: a positive number  $n$ ;  
 $i \leftarrow 2$ ;  
while(  $i \leq n$ ):  
     $i \leftarrow 2 \times i$ ;  
end while
```

The number of iterations the algorithm makes is  $O(?)$ .

**Que 6.** [10] Let a class have 25 students. Show that the probability that at least two students in the class will have the same birthday is at least  $1/2$ . (Ignore February 29 and assume that each student got a birthday uniformly randomly and independently from the 365 possible dates.)

**Hint:** Think about the probability that all of them have different birthdays.