The Graph structure of two-player games

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Introduction:

- 1. Fundamental question in Game theory: Representing preference
 - Von Neumann and Morgenstern's axiomatisation of utility
 - John Nash equilibrium concept
- 2. Non applicability of utility values to Ordinal Games
- 3. Non-convergence of Nash equilibria
- 4. Response Graphs

Response Graphs:

Graph G = (N,A), $N = \{n1,...,nk \text{ (finite set of nodes)}\}$, $A \subseteq N \times N$

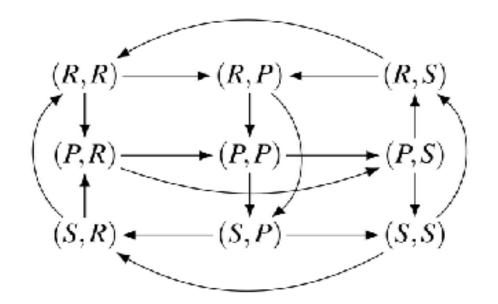
- 1. If arc $(x,y) \in A$ then $x \rightarrow y$ \Rightarrow Preference goes from x to y
- If arc (x,y) ∈ A and arc (y,x) ∈ A then x y
 Preference alloted equally to x and y.

The response graph of the game is the graph whose node set is $S1 \times S2$, with an arc (s1,s2) (t1,t2) if the profiles (s1,s2) and (t1,t2) are i-comparable and $ui(t1,t2) \ge ui(s1,s2)$.

Example:

Rock
Player 1 Paper
Scissors

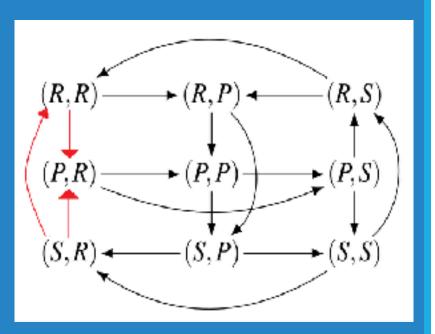
	Player 2	
Rock	Paper	Scissors
0,0	-1,1	1, -1
1,-1	0,0	-1,1
-1,1	1,-1	0,0



Response Graphs:

- 1. Utility payoffs serve to instantiate preference orders, so cut the intermediary.
- 2. No need for labelling by profiles
- 3. 'Sink strongly connected components' in 'sink chain components'

Sink Components



- 1. Solution concept generalizing pure Nash equilibria
- 2. Contained in Sink chain component
 - → A topological concept emerged from Fundamental theory of Dynamical Systems.
- 3. Represent 'long-run' outcome of dynamic processes (learning or evolution of a game).
- 4. Dynamic and Predictive solution concept for games unlike Nash equilibria

Weighted Response Graphs:

Generalization of Response graphs for biased games.

• Arc weights equal payoff differences for associated player.

Decomposition to strategic equivalence.

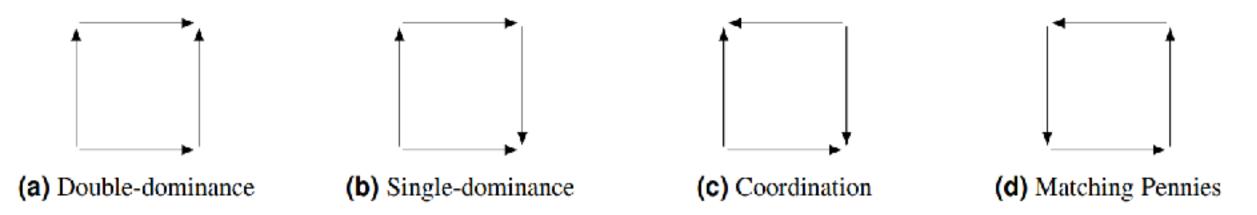
- Depends on the payoff for mixed profiles
- (S,P) preferred over (P,R)

Modulo Isomorphism: preference equivalence.

Preference Games:

- 1. Preference single dominance games
- 2. Preference double dominance games
- 3. Preference-zero-sum games
- 4. Preference-potential games

Preference Games:



The four non-isomorphic response graphs of generic 2×2 games.

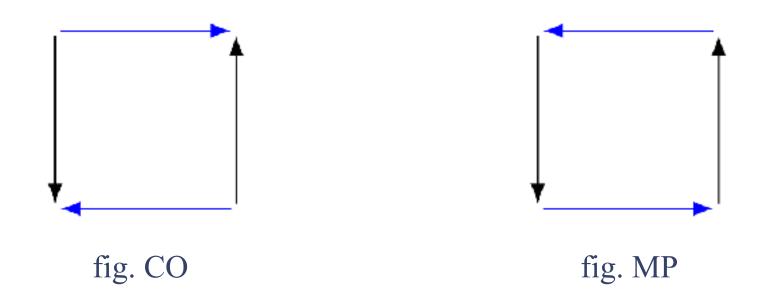
Two-Player zero-sum and Potential Duality:

$$(c,-c) \leftarrow c-d \qquad (d,-d)$$

$$(-c)-(-a) \qquad \qquad \downarrow (-b)-(-d)$$

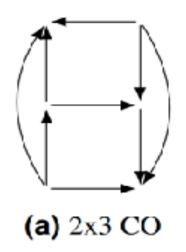
$$(a,-a) \longrightarrow (b,-b)$$

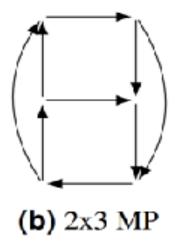
Two-Player zero-sum and Potential Duality:

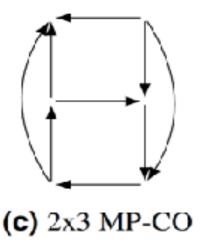


The CO graph on the left when reflected for second player we get graph for matching pennies

Applications: 2x3 Games

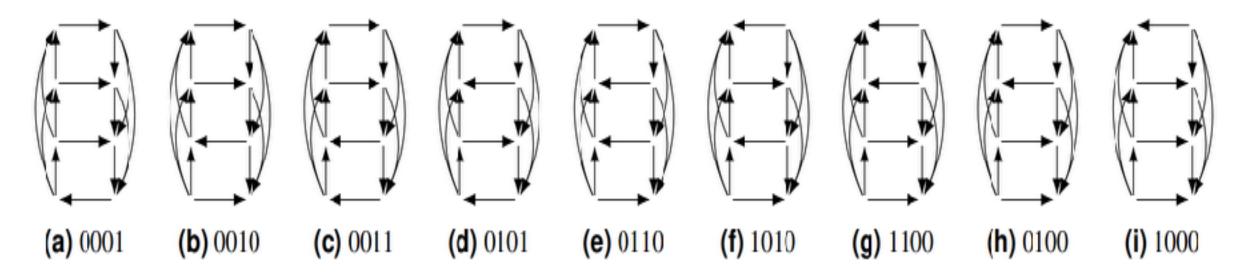






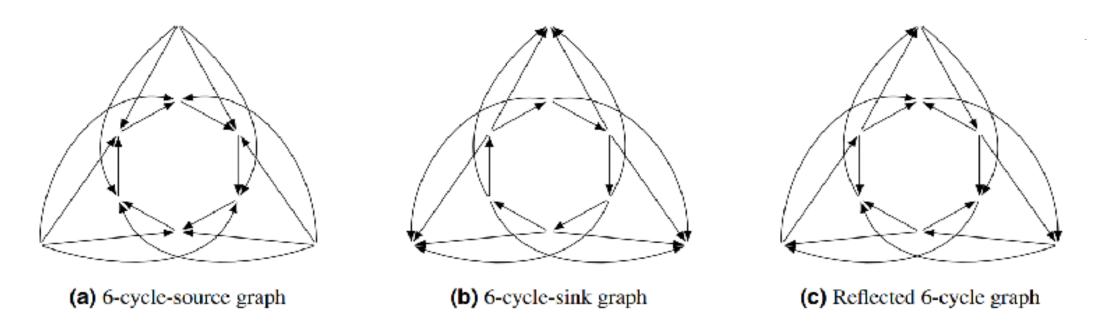
- Three fundamental non-isomorphic response graphs for 2x3 games
 - a) Preference-potential
 - b) Preference-zero-sum
 - c) Neither of them (unique minimal example)

Applications: 2x4 Games



- Binary encoding represents direction of arc for player 2
 - a) Preference-potential \rightarrow g) and i).
 - b) Preference-zero-sum \rightarrow a) and c).
 - c) Neither of them \rightarrow rest of them

Applications: 3x3 Games



- a) and its reversal b), do not contain MP and yet are not preference-potential.
- Reflection of player preferences in a) and b) gives c). It does not contain CO yet it is not preference-zero-sum.

Conclusion:

- 1. We defined a model which captures underlying notion of strategic preference without concerning about cardinal values for payoff.
- 2. This model can be used when access to or knowledge of cardinal payoffs is implausible.
- 3. We showed that two-player potential and zero-sum games have analogous properties of sink component.

References:

Biggar, Oliver, and Iman Shames. "The graph structure of two-player games." *Scientific Reports* 13, no. 1 (2023): 1833. (Main Paper)

Link: [2209.10182v2] The graph structure of two-player games (arxiv.org)

Thank You!