

Process modeling in Petri Nets

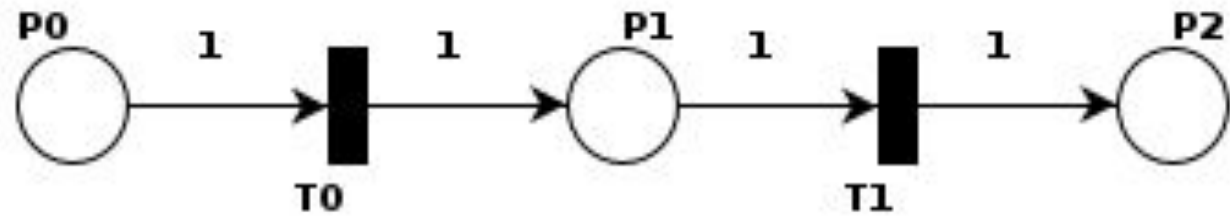
Basics

Rushikesh K Joshi

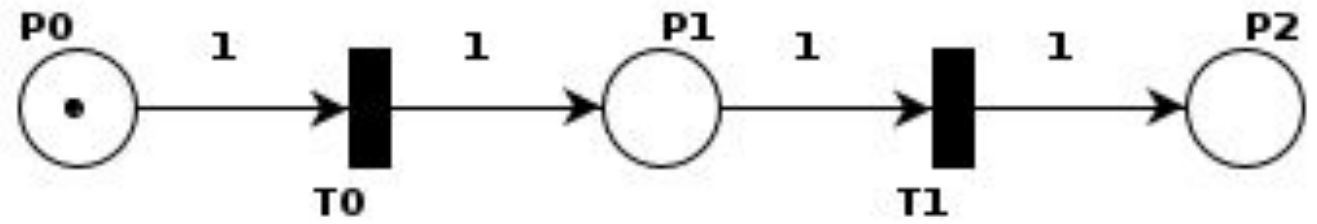
Computer Science & Engineering

IIT Bombay

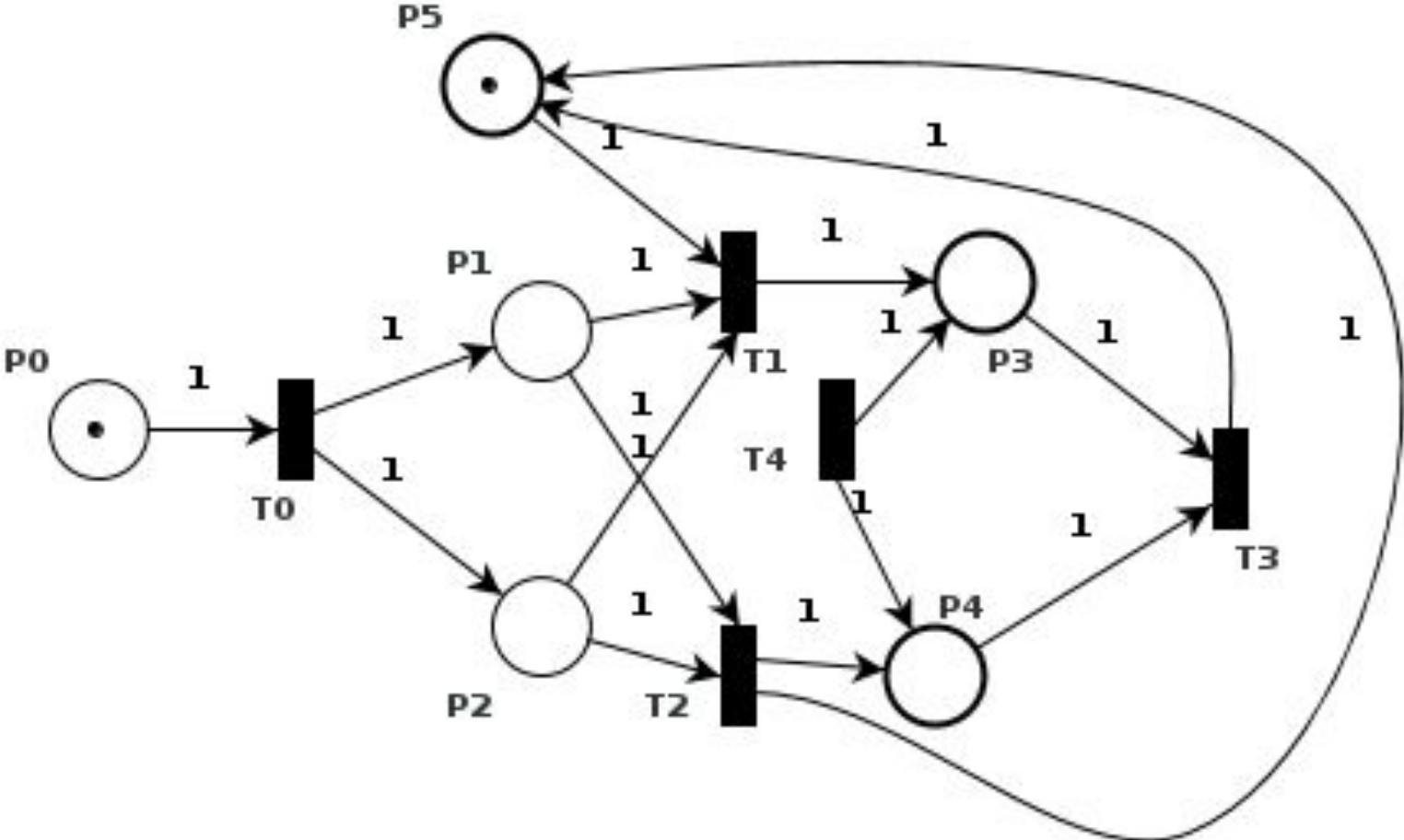
Example Petri Nets



Petri Net with a token



A Complex Petri Net



Circles are called **Places**

Rectangles are called **Transitions**

Tokens can be inserted in a set of places as **initial condition (A Marking)**

Then as per the **rules** of the petri-net language, the tokens move through the Petri Net (Markings change)

This is “**execution**” of our process, that is, state changes.

At a given point of time, **multiple places may have tokens**,

The set of places which have tokens in them at a given point of time is called **Marking**

If pre-conditions of transitions are fulfilled (all its input places have tokens), the transitions are called **“enabled”**

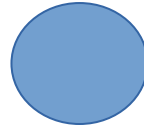
One of the triggered transitions can **“fire/trigger”**

Once a transition fires, the **tokens** in the pre-places are **moved out** into output places of that transition

Basic Petri-net Elements

- Places

- Represent states



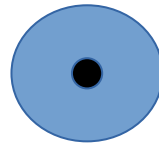
- Transitions

- Represent actions



- Tokens

- Part of current state



- Arrows

- Flow elements

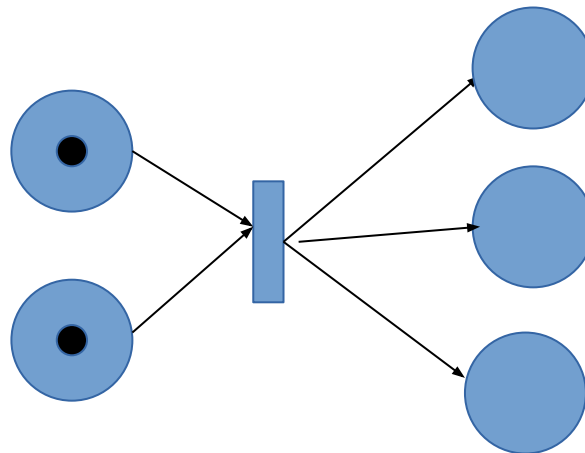


Elementary Petri-Nets

- at most one token per place

Transition Enabling Rule

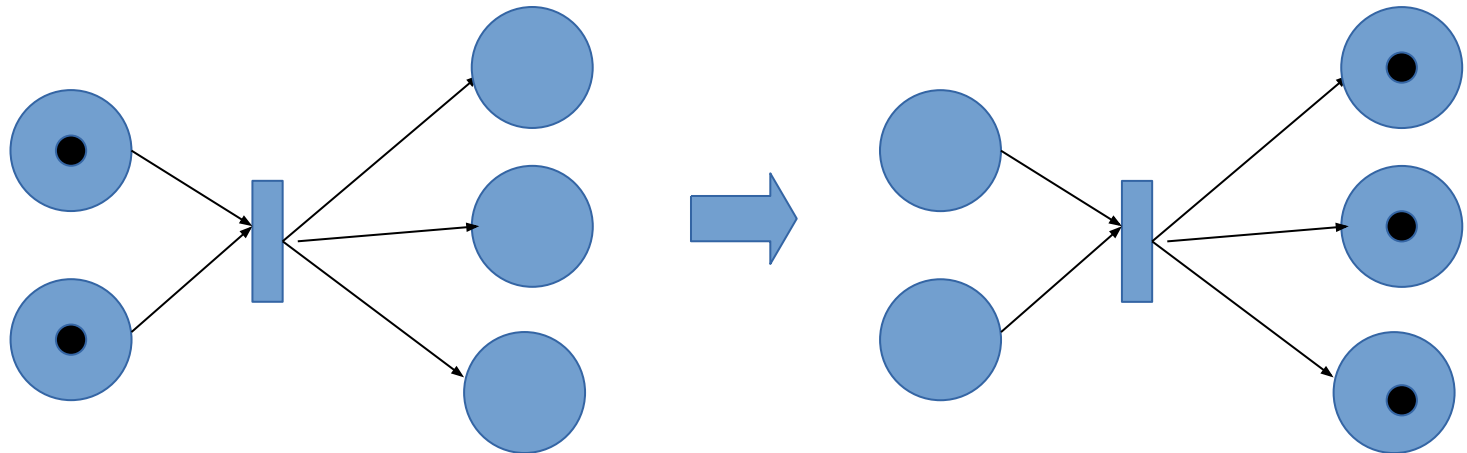
- each of its preplaces (preconditions) has a token
- after it is fired, each of its post places must be able to hold one new token



The above transition in the picture is enabled

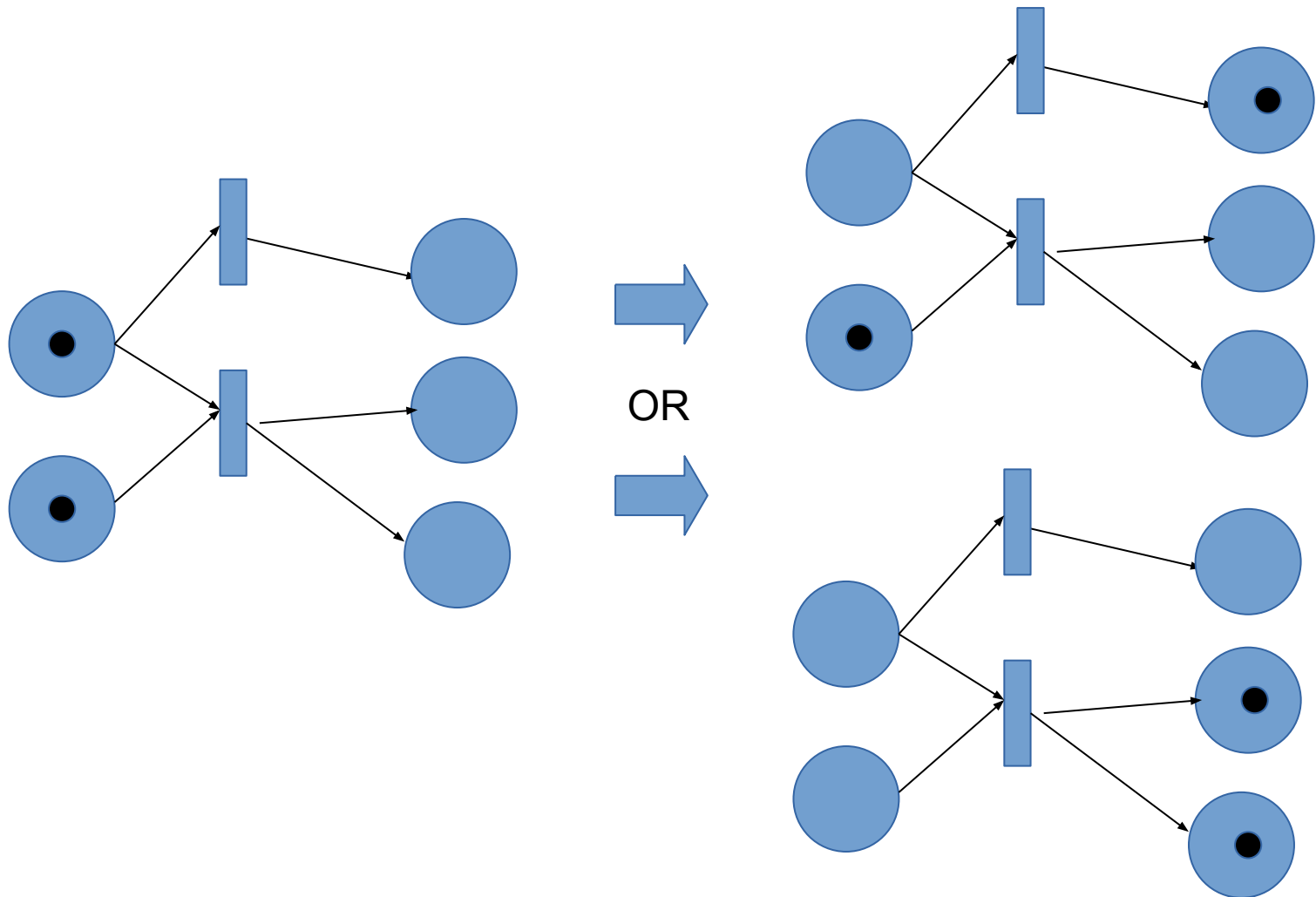
Transition Firing Rule and Example

- Enabled transition **fires**, and it moves the token(s) downstream Into all post places
 - One token is removed from every preplace
 - One token is deposited in every output place

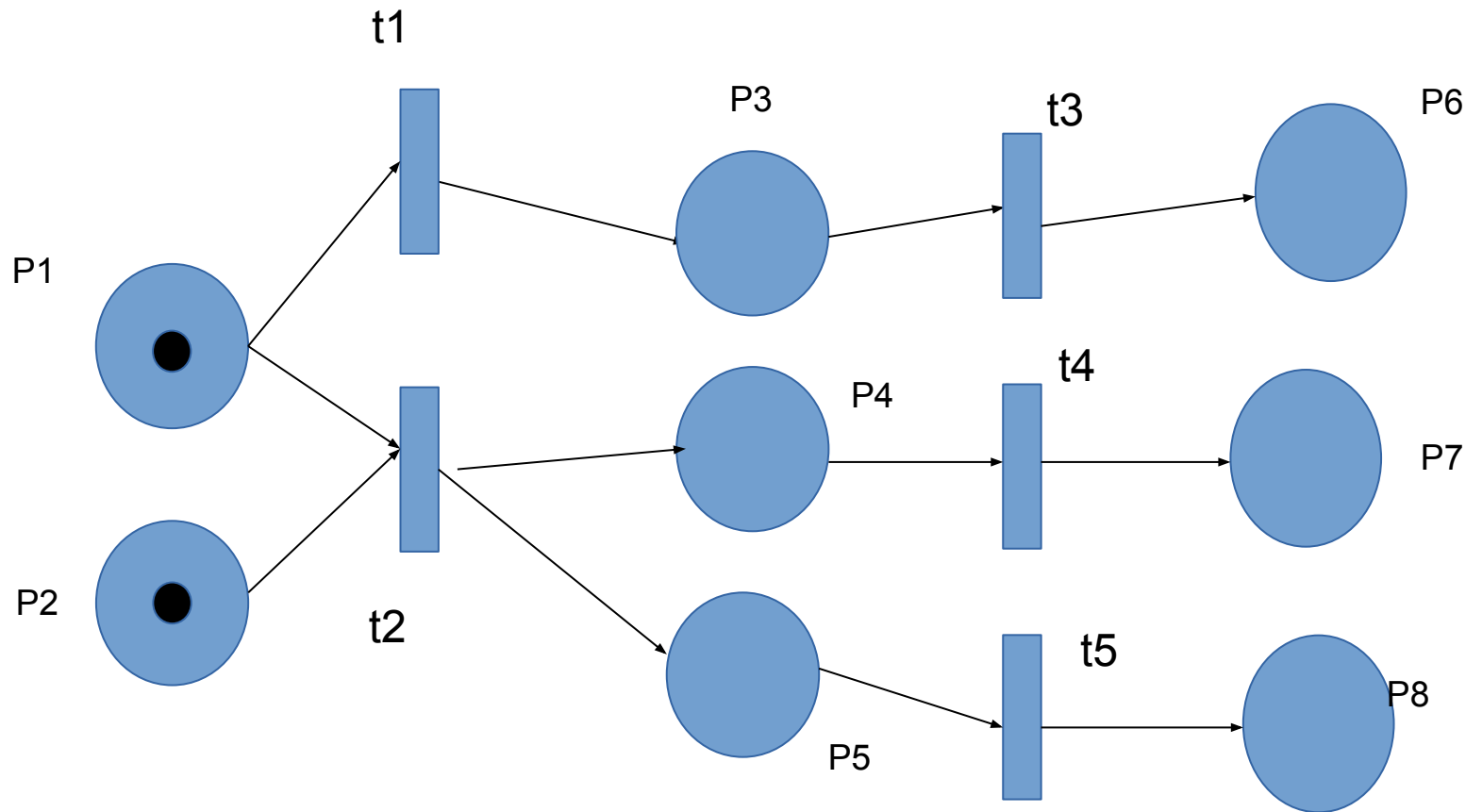


Transition Enabling Rule- multiple transitions may get enabled due to a single marking

- A marking may enable one or more transitions
- Any one of them can be fired



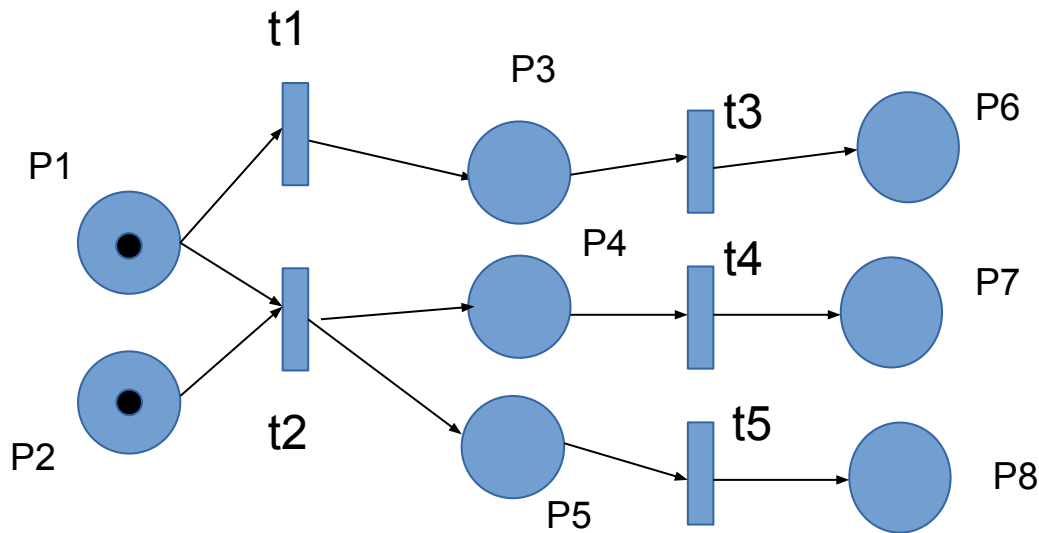
We label the places and transitions
So that we can analyze the state changes



Current Marking = {P1, P2}

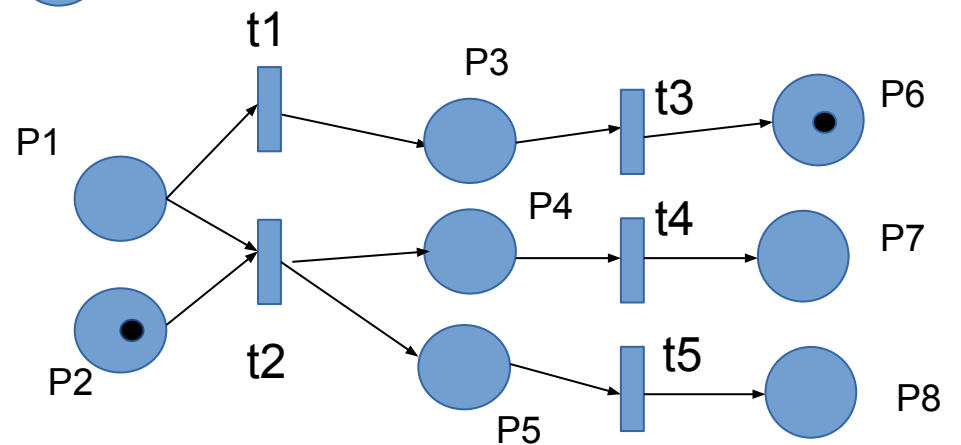
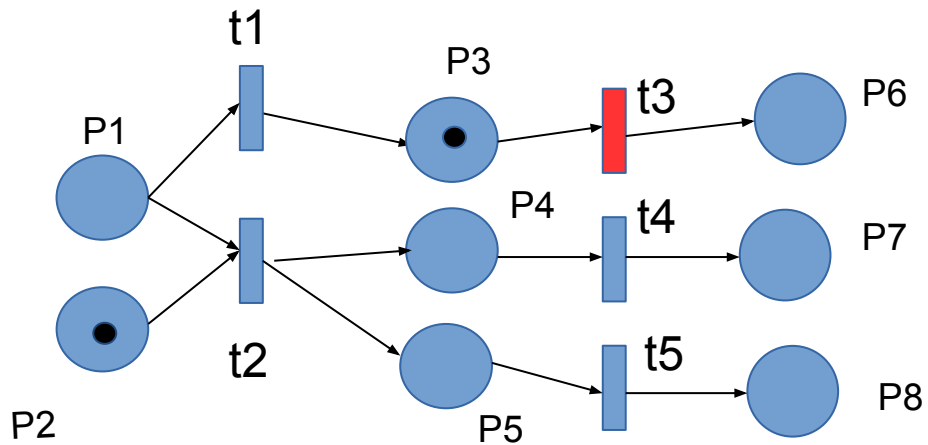
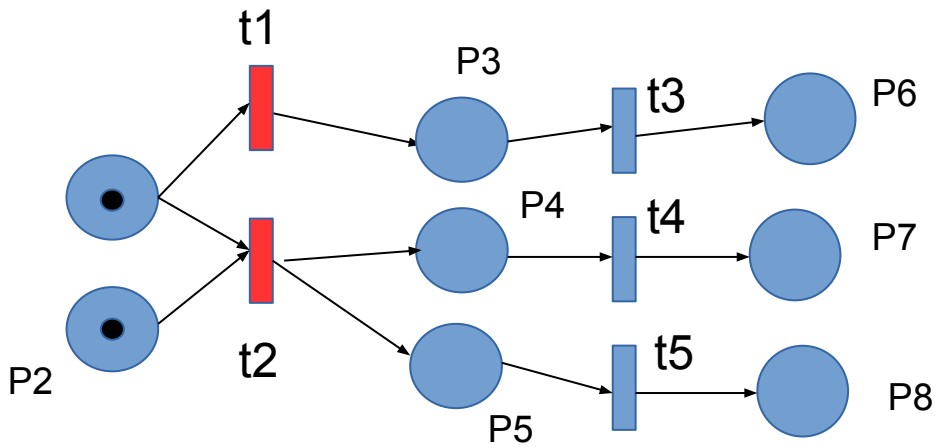
Set of enabled transitions = {t1, t2}

Firing Sequence (Trace) is sequence of transition firings
We can have many traces in a given marked net, since a marking can enable more than one transitions

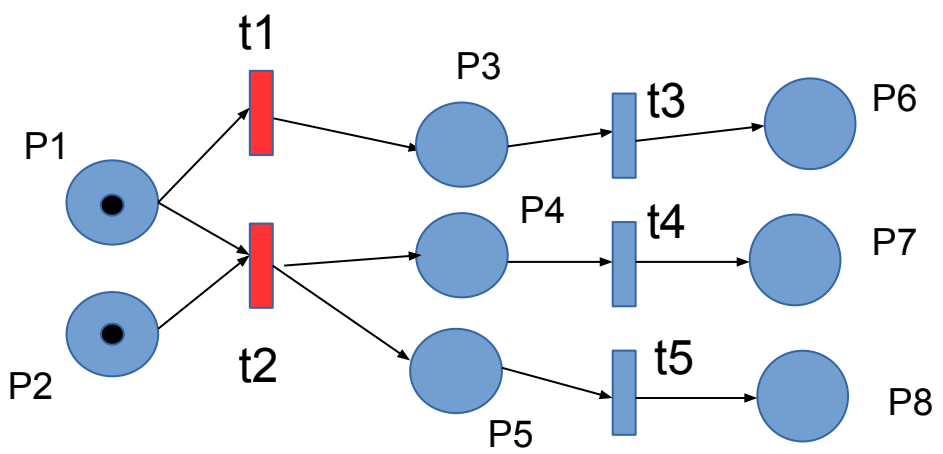


- Here, the following firing sequences (traces) are possible. Note that only one transition fires at a time
 - $t1 \rightarrow t3$ (trace 1)
 - $t2 \rightarrow t4 \rightarrow t5$ (trace 2)
 - $t2 \rightarrow t5 \rightarrow t4$ (trace 3)

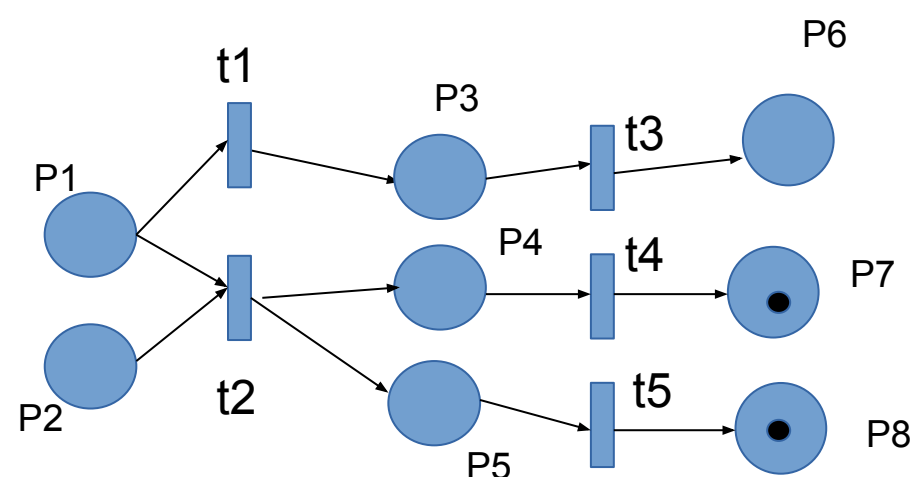
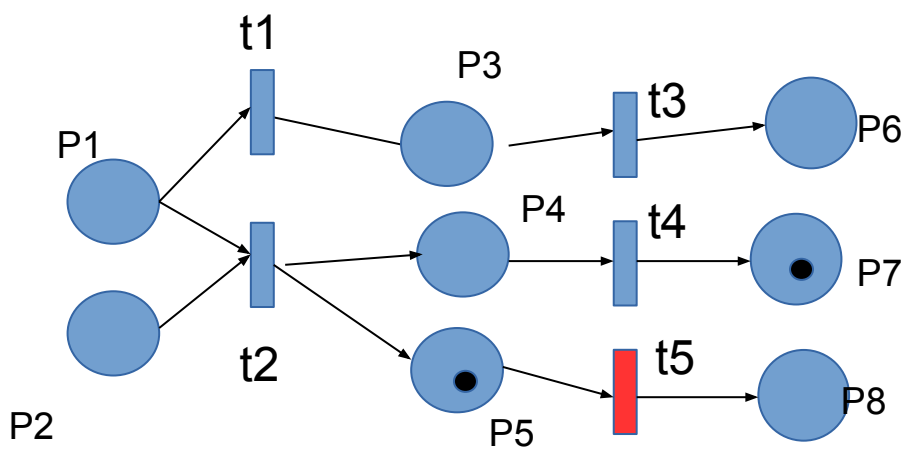
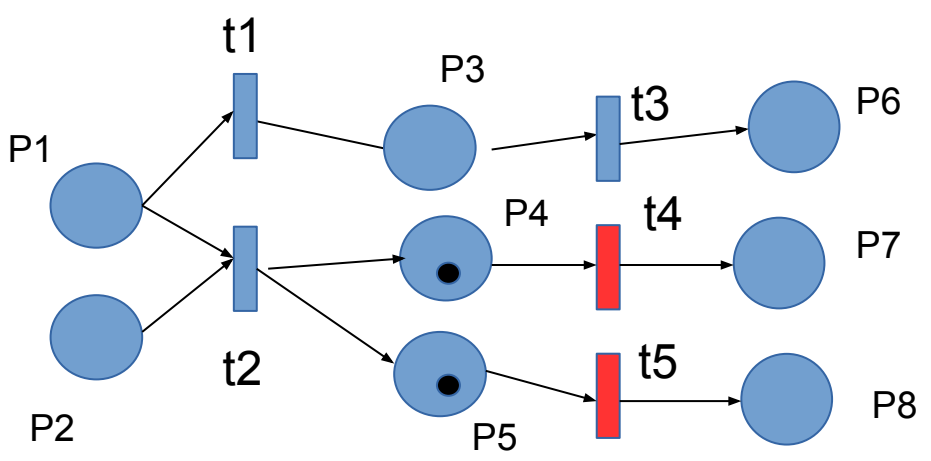
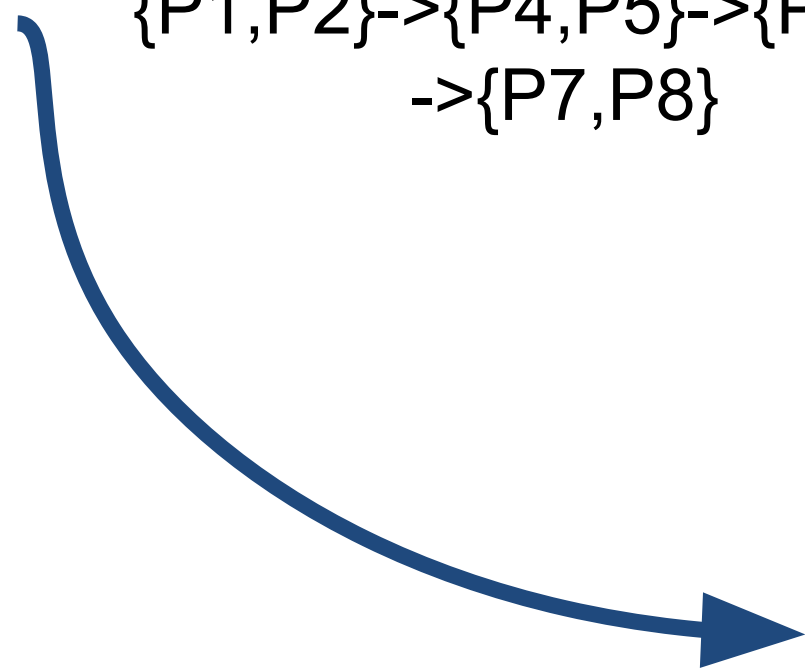
Trace 1 : t1 -> t3
 Through states (markings):
 $\{P1, P2\} \rightarrow \{P3, P2\} \rightarrow \{P6, P2\}$



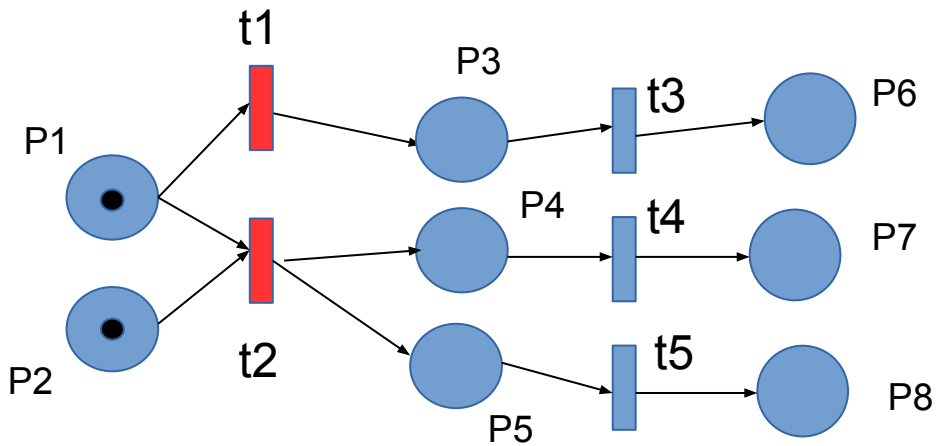
Enabled transitions are shown in red



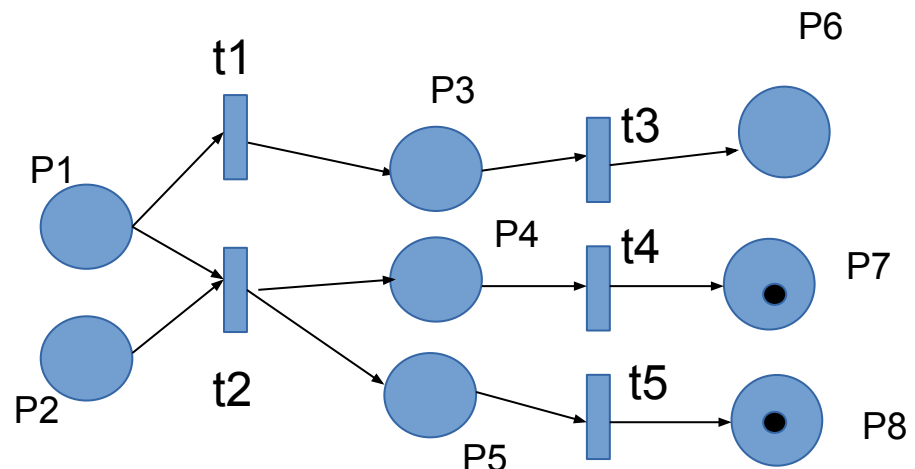
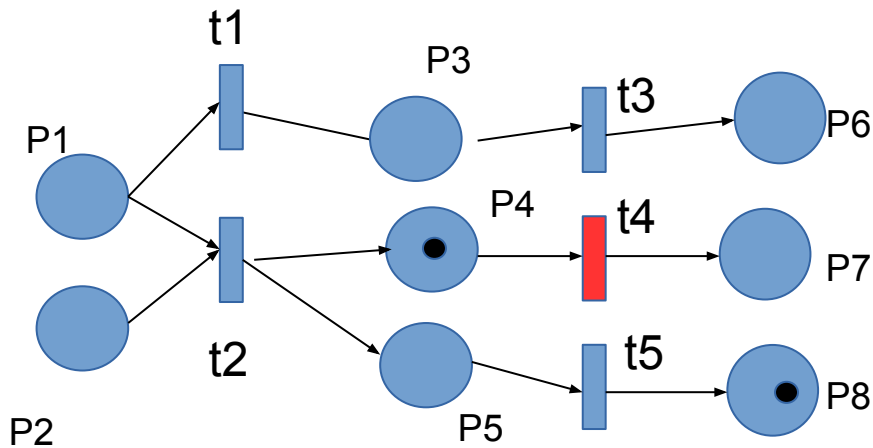
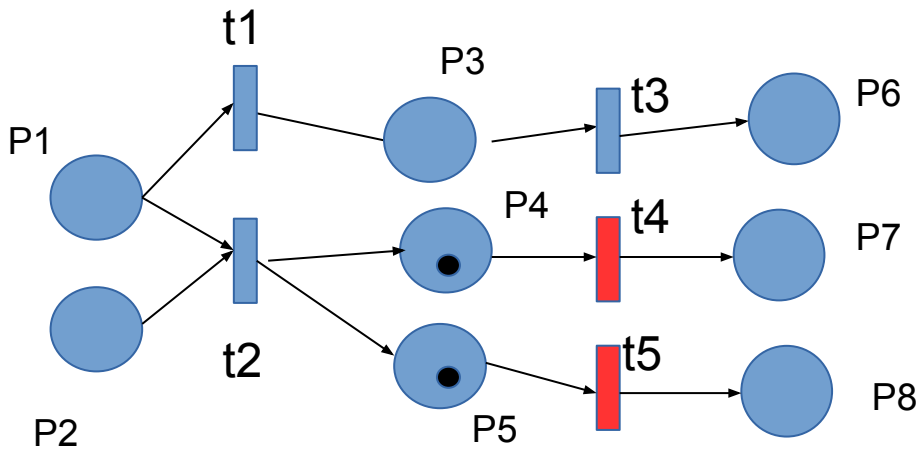
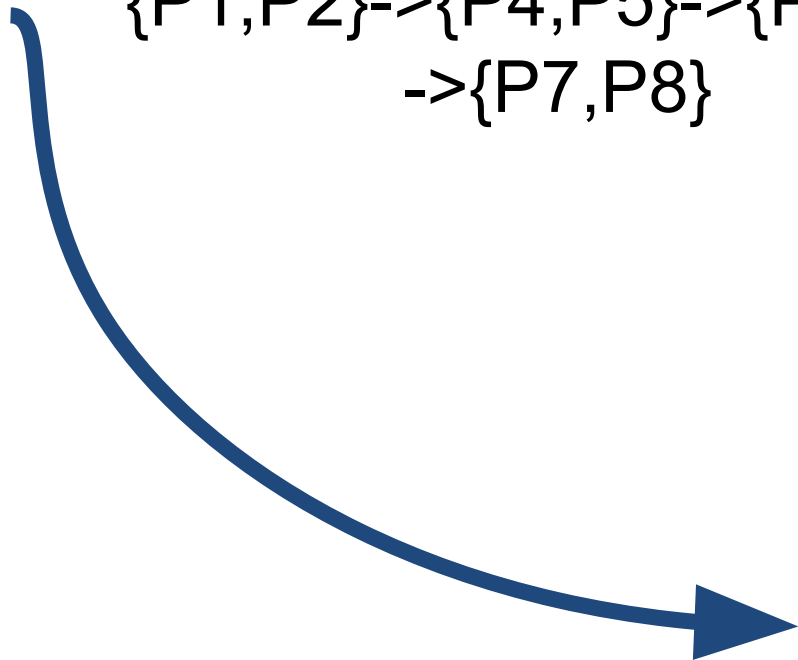
Trace 2 : t2 -> t4 -> t5
 Through states (markings):
 $\{P1, P2\} \rightarrow \{P4, P5\} \rightarrow \{P7, P5\} \rightarrow \{P7, P8\}$



Enabled transitions are shown in red



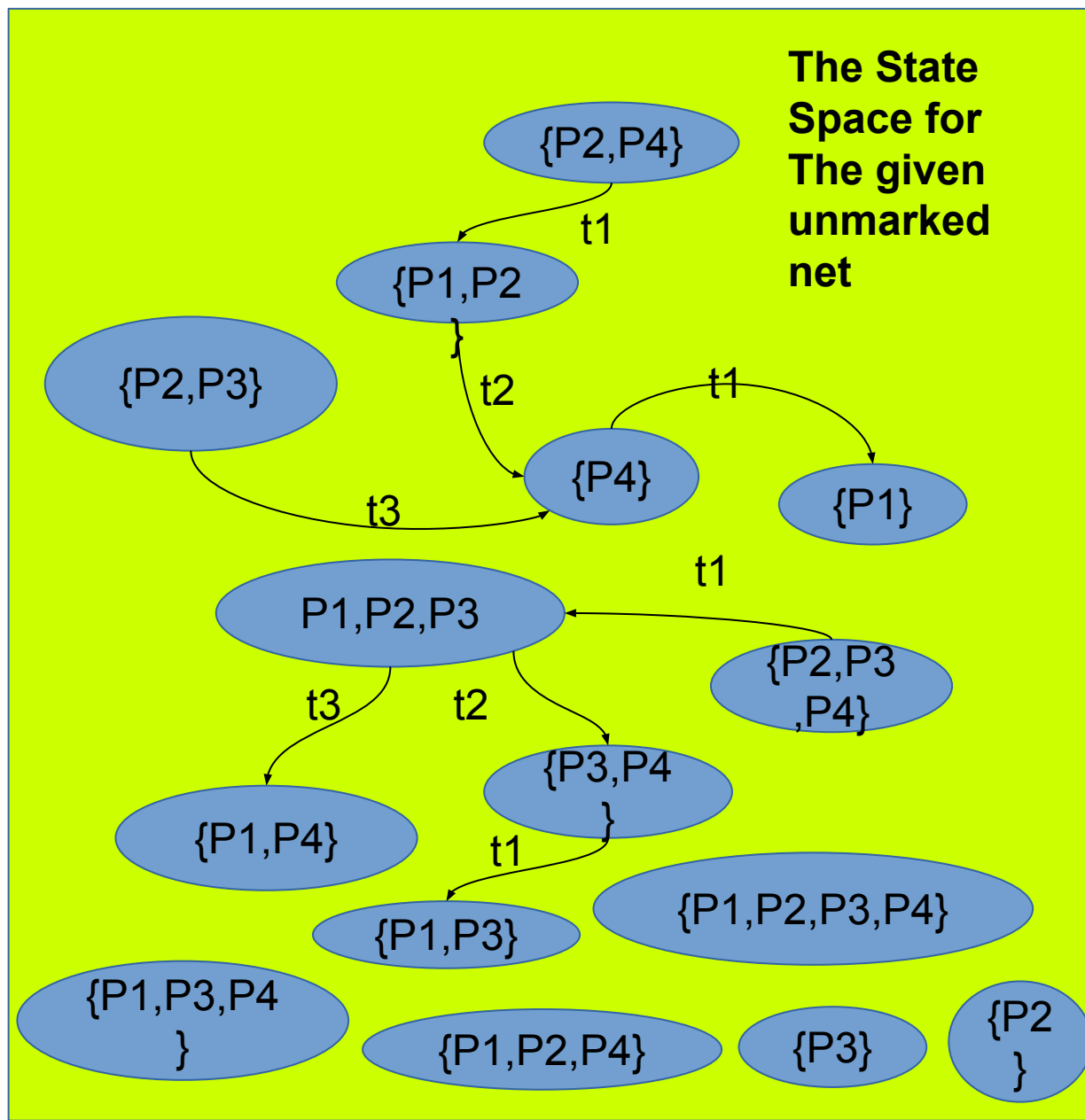
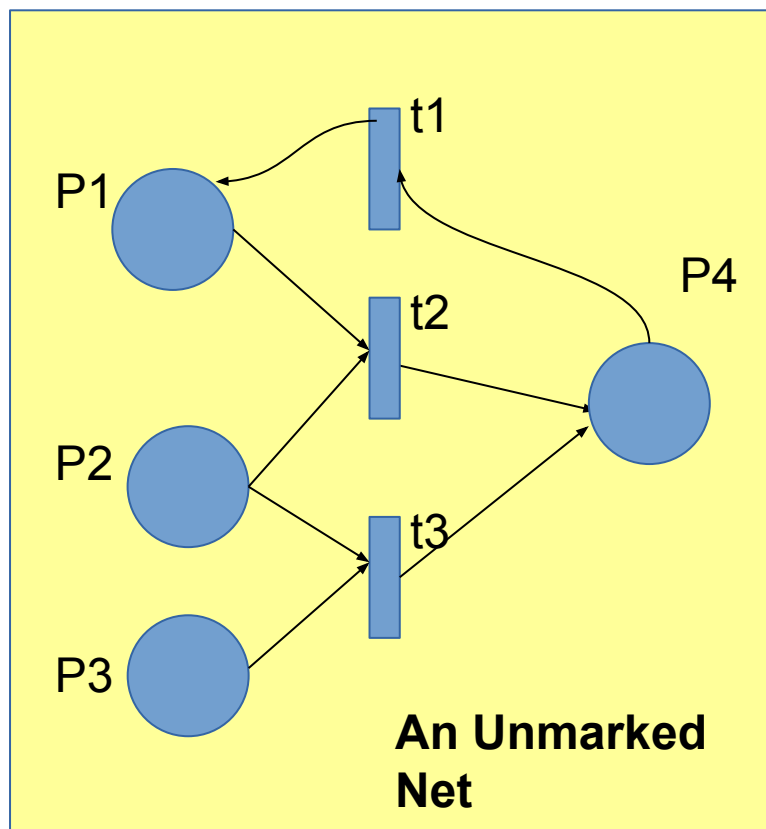
Trace 3 : t2 -> t5 -> t4
 Through states (markings):
 $\{P1, P2\} \rightarrow \{P4, P5\} \rightarrow \{P4, P8\}$
 $\rightarrow \{P7, P8\}$



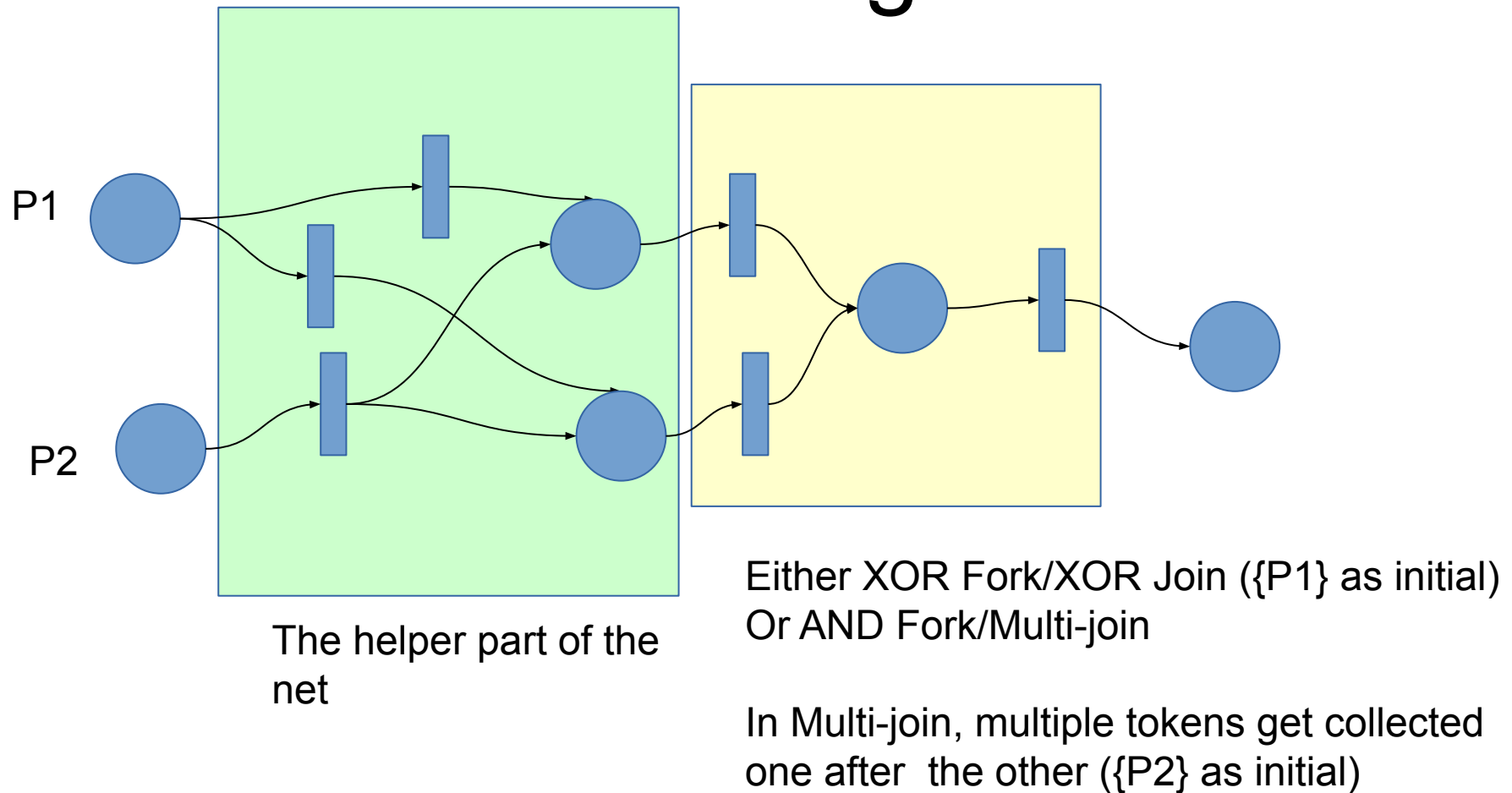
Enabled transitions are shown in red

State Space: State transition diagram

set of all possible markings, and the transitions through them

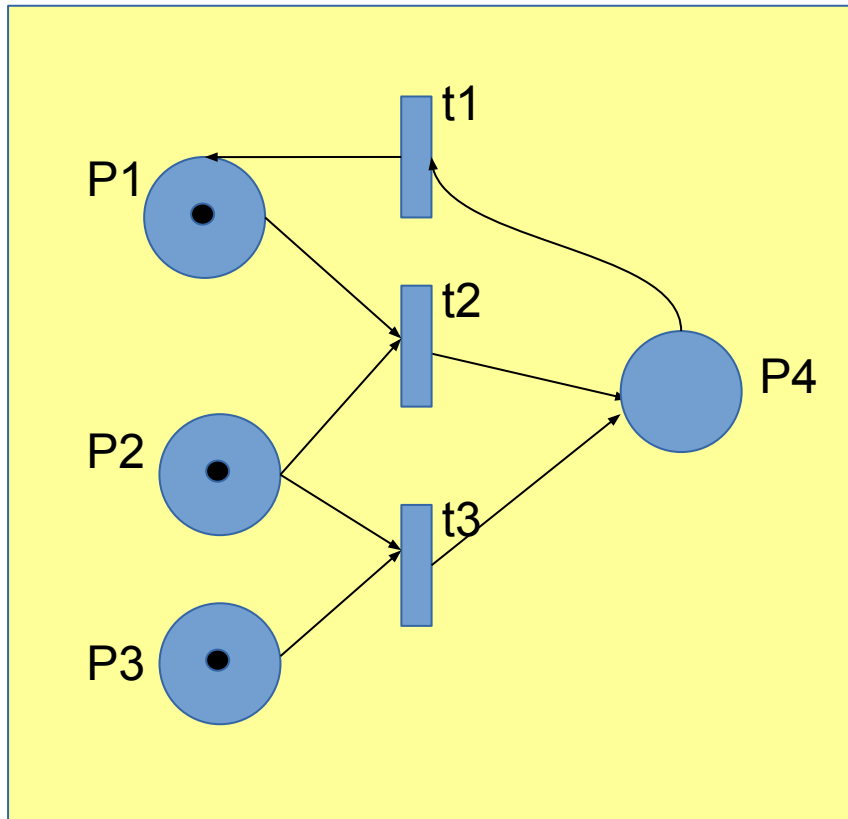


The same Petri-net for different purposes with different initial markings

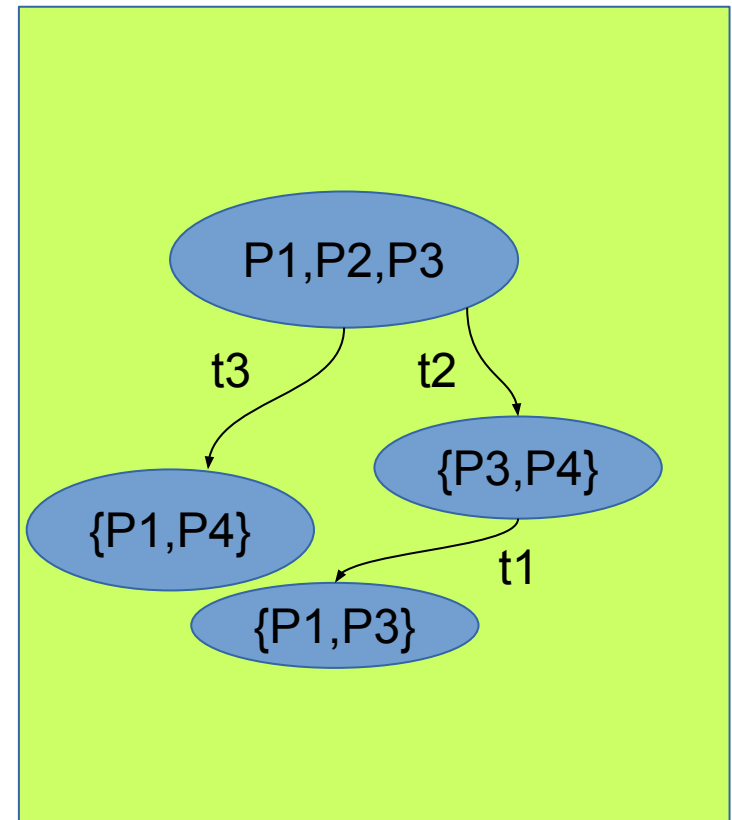


Reachability Graph

= state space reachable from a given initial marking

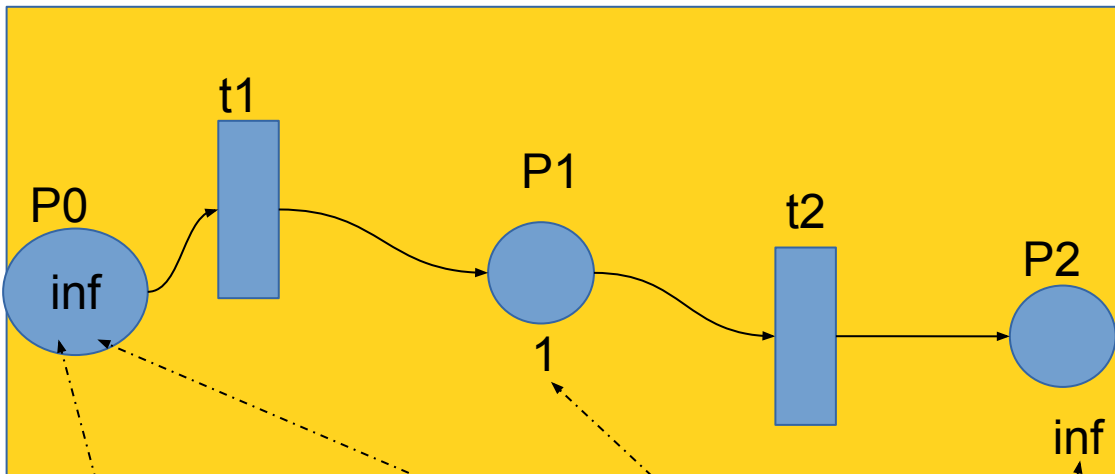
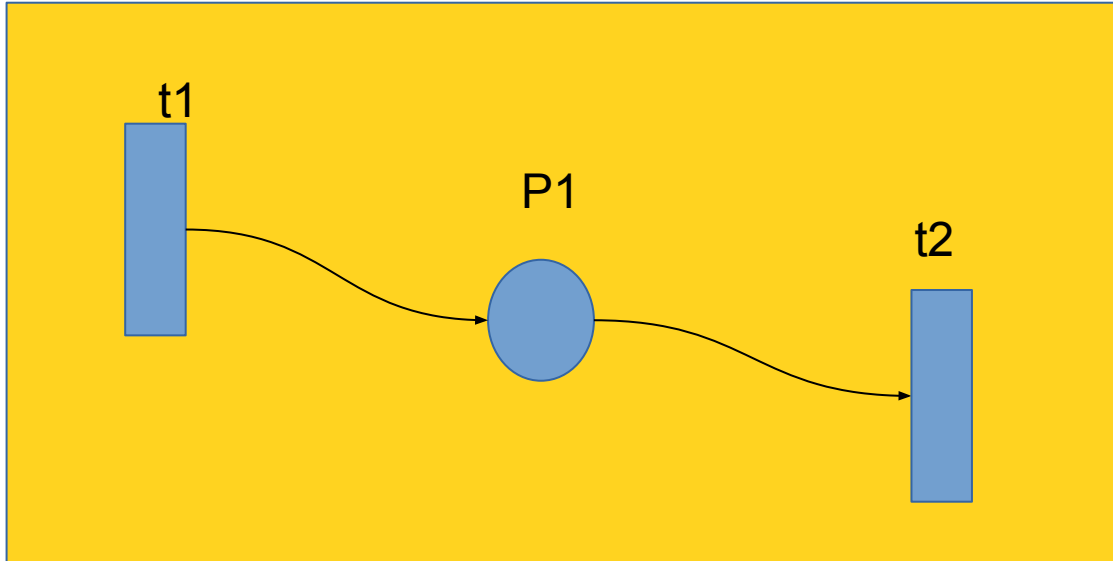


A Net with initial marking as $\{P1, P2, P3\}$



Reachability Graph for the given net with given initial marking $P1, P2, P3$

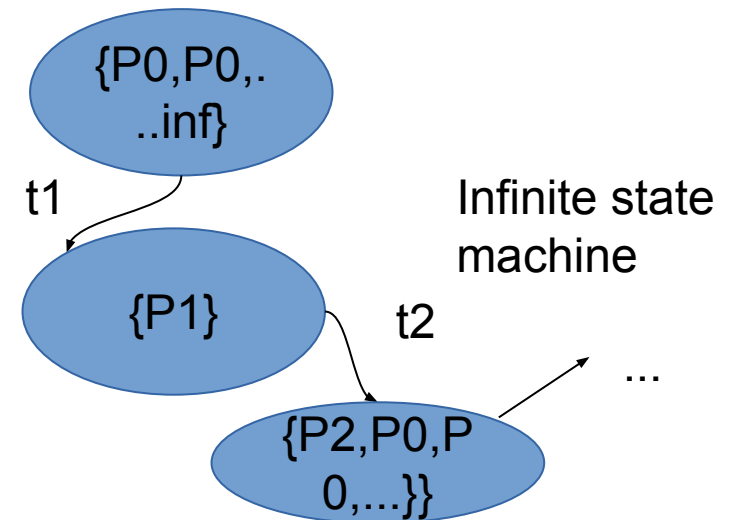
Infinately enabled net



P0 has inf tokens

The number written inside represents those many (initial) tokens in the given state of the machine.

Capacity of P_1 is 1,
 capacity of P_0 ,
 Capacity of P_2 is inf
 Inf is default capacity in classical nets



A Problem: Rules of one masters program

- It's 2 semester program
- At most 10 courses
- Minimum 8 courses
- Per semester max 5 courses
- One R&D project can be taken as one course
- One masters project can be taken in place of 3 courses.
- MTP 1 in sem 1: counted as 3 courses, MTP 2 in sem 2: counted as 2 courses
- 1 seminar in 1st sem is must
- Backlogs of sem1 can be taken in sem2
- If any backlogs remain, one more semester is granted

Simplified Version of the Problem

- 5 courses
- 2 semester
- 2 courses per semester: total exactly 4 courses
- A course cannot be taken twice

- Make your net to represent this system
- Try to reduce the no. of transitions
- Model with: One transition representing one course

Vending machine Problem

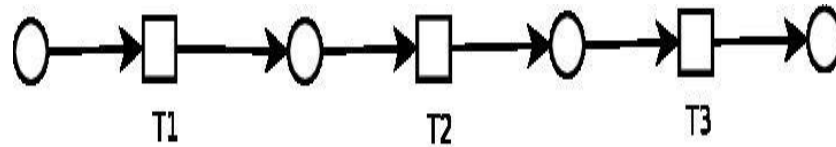
- Insert Rs. 25: currency Rs. 5, Rs. 10 coins accepted e.g. 5+5+5+5+5, 10+5+10 etc. (design as small a machine as possible)
- Choose the drink
- Pick it up
- The machine is ready for the next task
- Time out of 1 minute, the machine returns all the coins inserted and resets to initial state
- Invalid coin is rejected, and all the coins are returned, the machine resets with a spoken message of invalid coin.

Workflow Nets

Workflow Nets

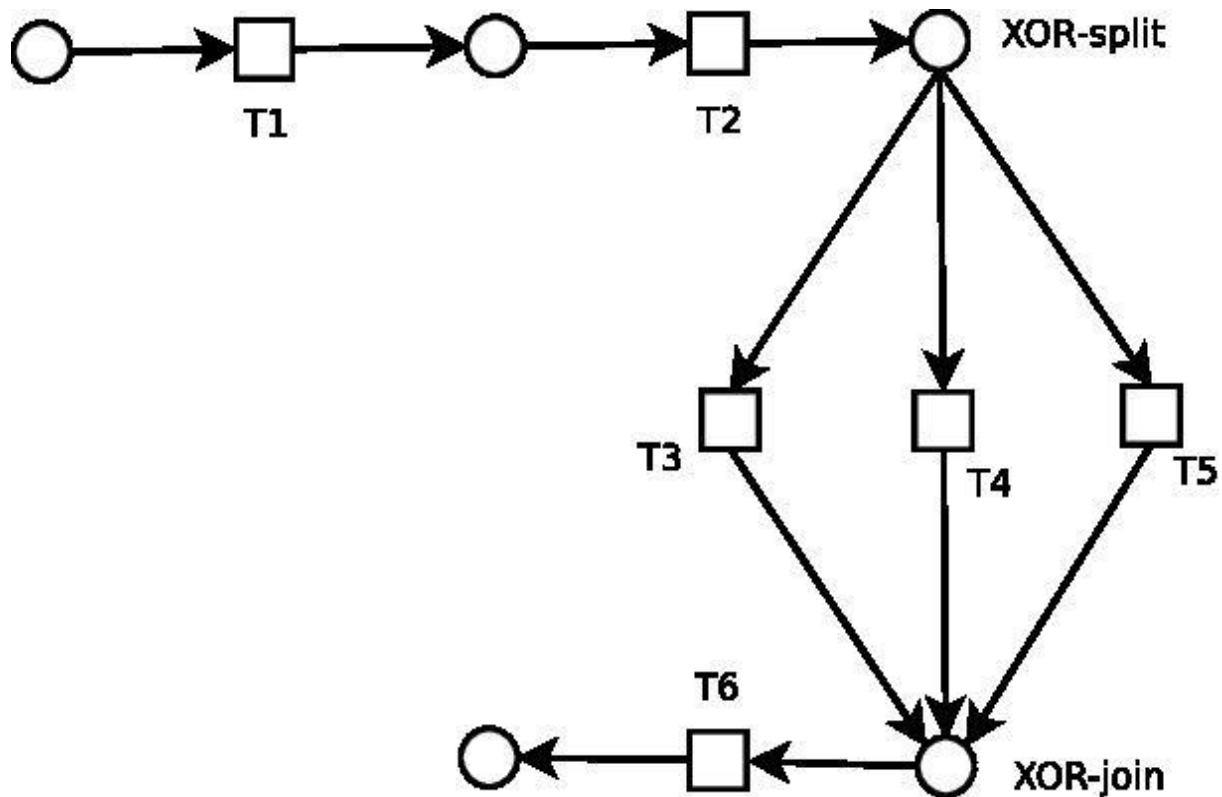
- Unique source place
- Unique sink place
- Connected
- Unique initial marking, unique terminal marking
- Well-formed – every transition is reachable, every marking is reachable, every marking terminates

Workflow Patterns: SEQUENCE

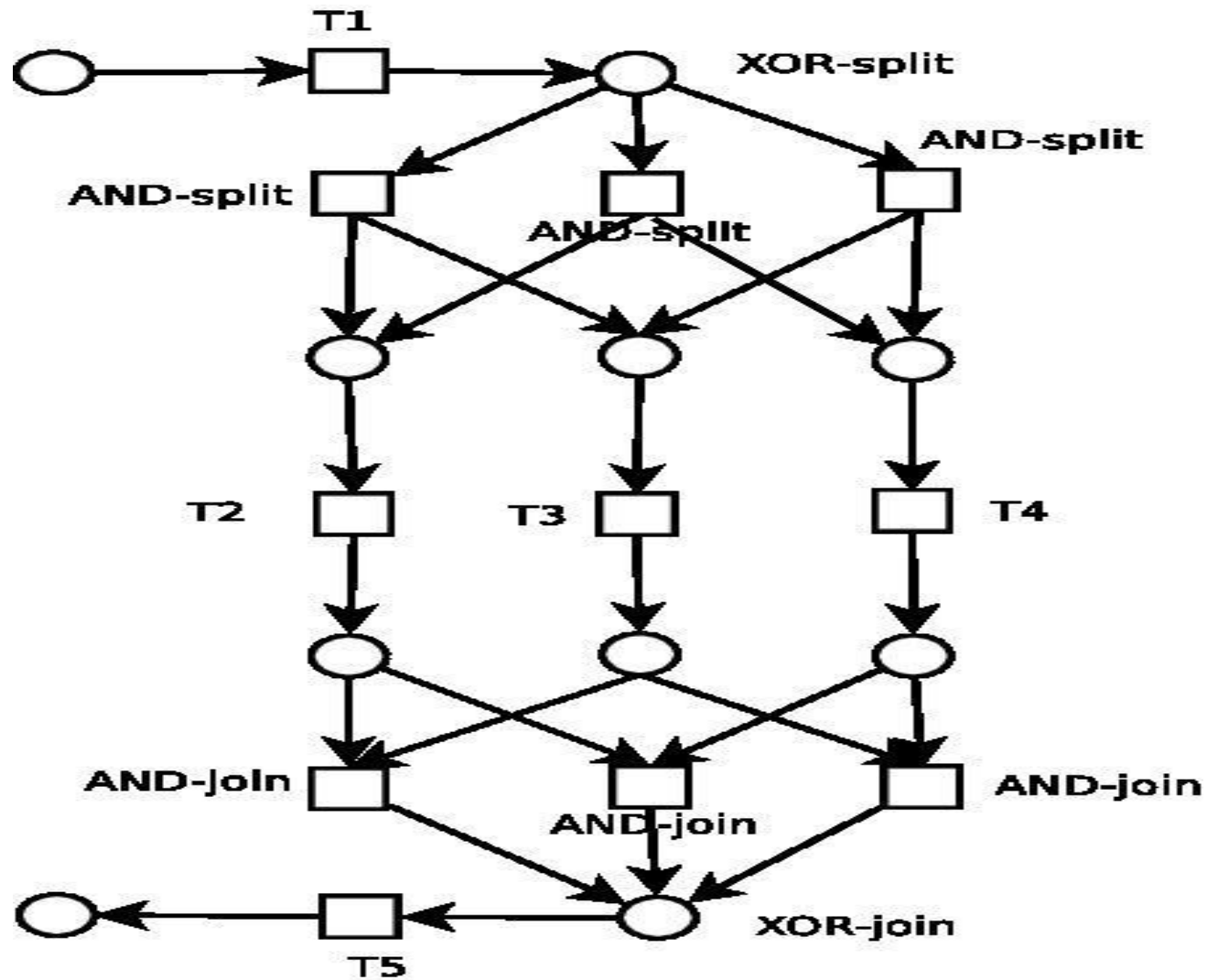


Sequence

Workflow Patterns: XOR

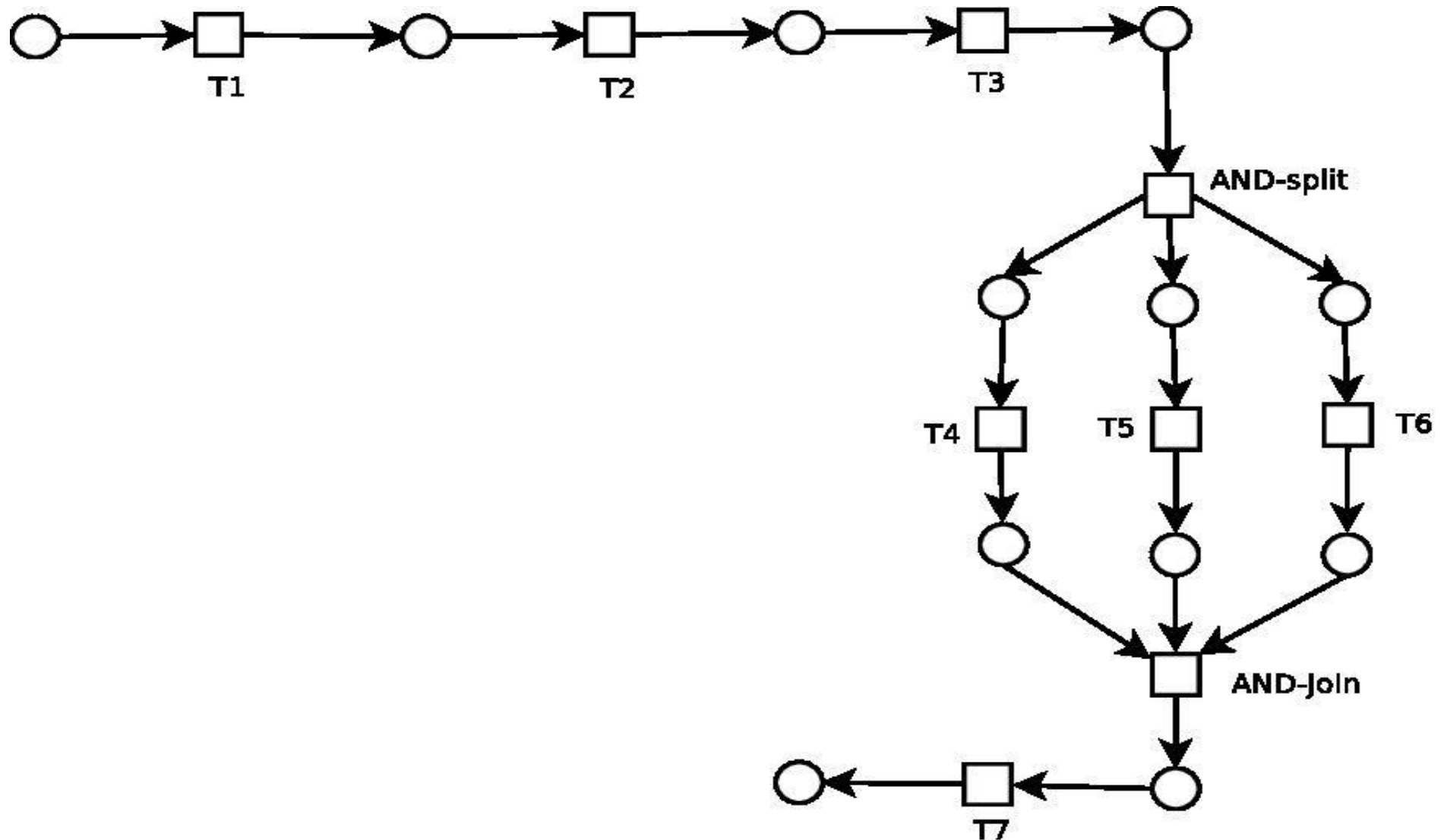


Workflow Patterns: 2/3 Choice



2/3 split-join

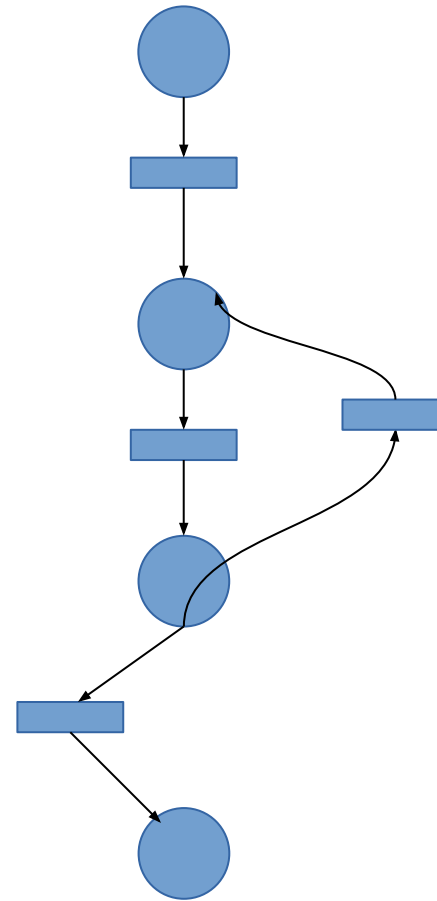
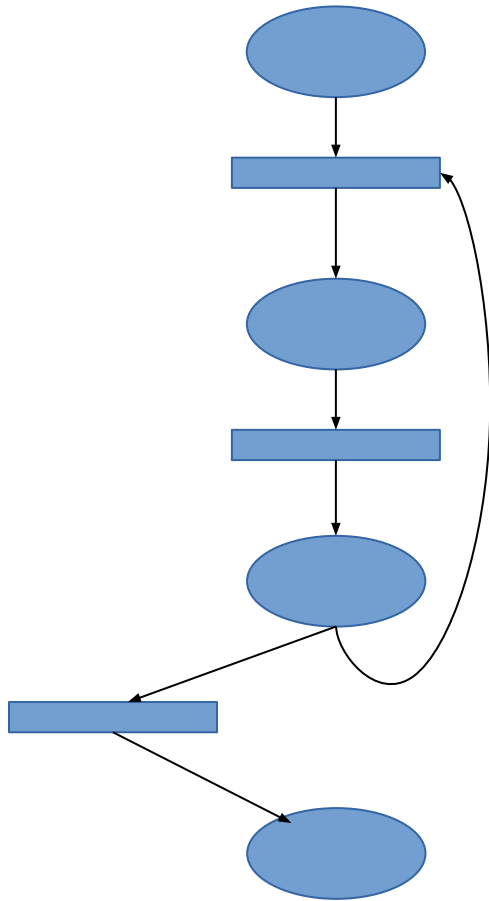
Workflow Patterns: AND



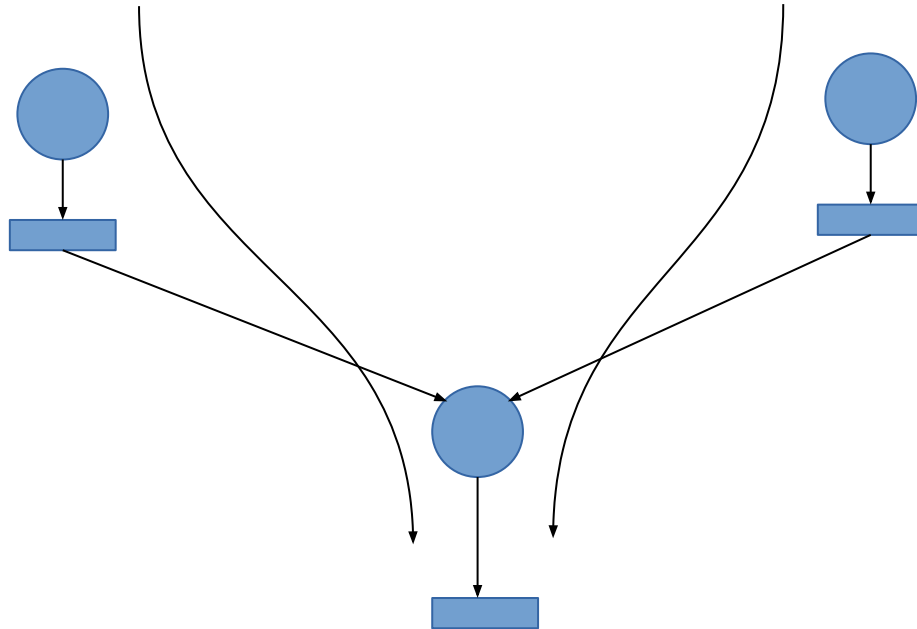
Workflow Patterns: Iteration

which one is correct? Which one is incorrect?

Why?



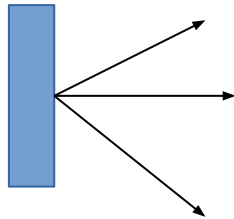
Workflow Patterns: Multi-merge



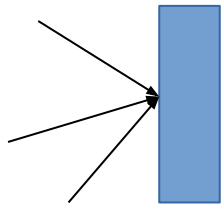
Tokens can arrive both ways, and they are all sent down

Split/Join

- Parallel split

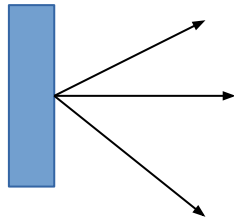


- Synchronization (parallel merge/AND join)

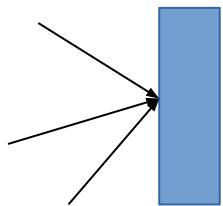


Split/Join

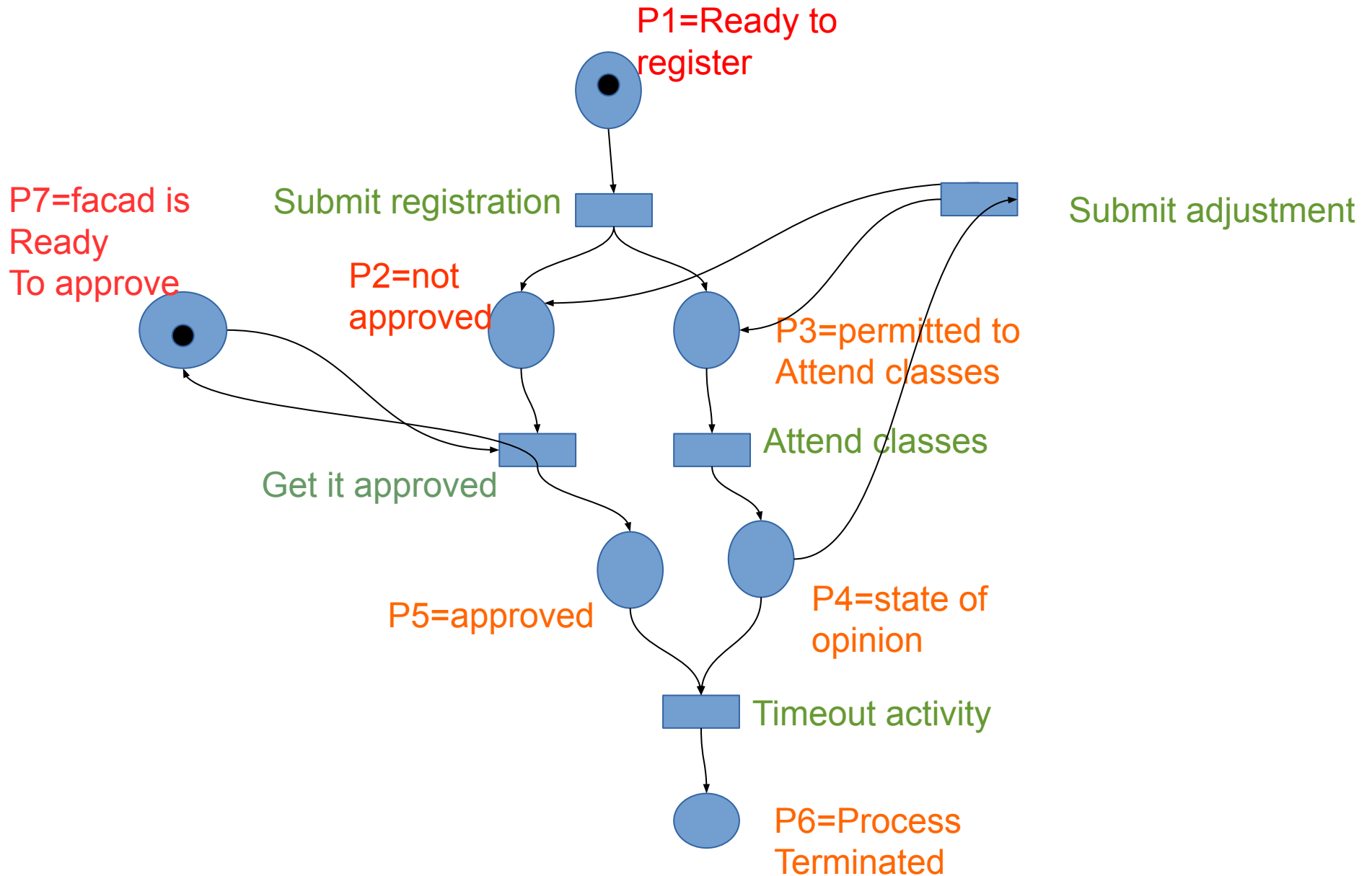
- Parallel split



- Synchronization (parallel merge/AND join)



Roles as tokens in places



Types of activities

- . Automatic Activity

- Computer can execute it fully

- . (when enabled, it is executed automatically such as by an algorithm, script task etc.)

- . User Activity

- A human being executes it

- . (though enable, it is done manually)

- . Message Activity

- An external message triggers the task instance

- . (though enabled, it requires a message to trigger it)

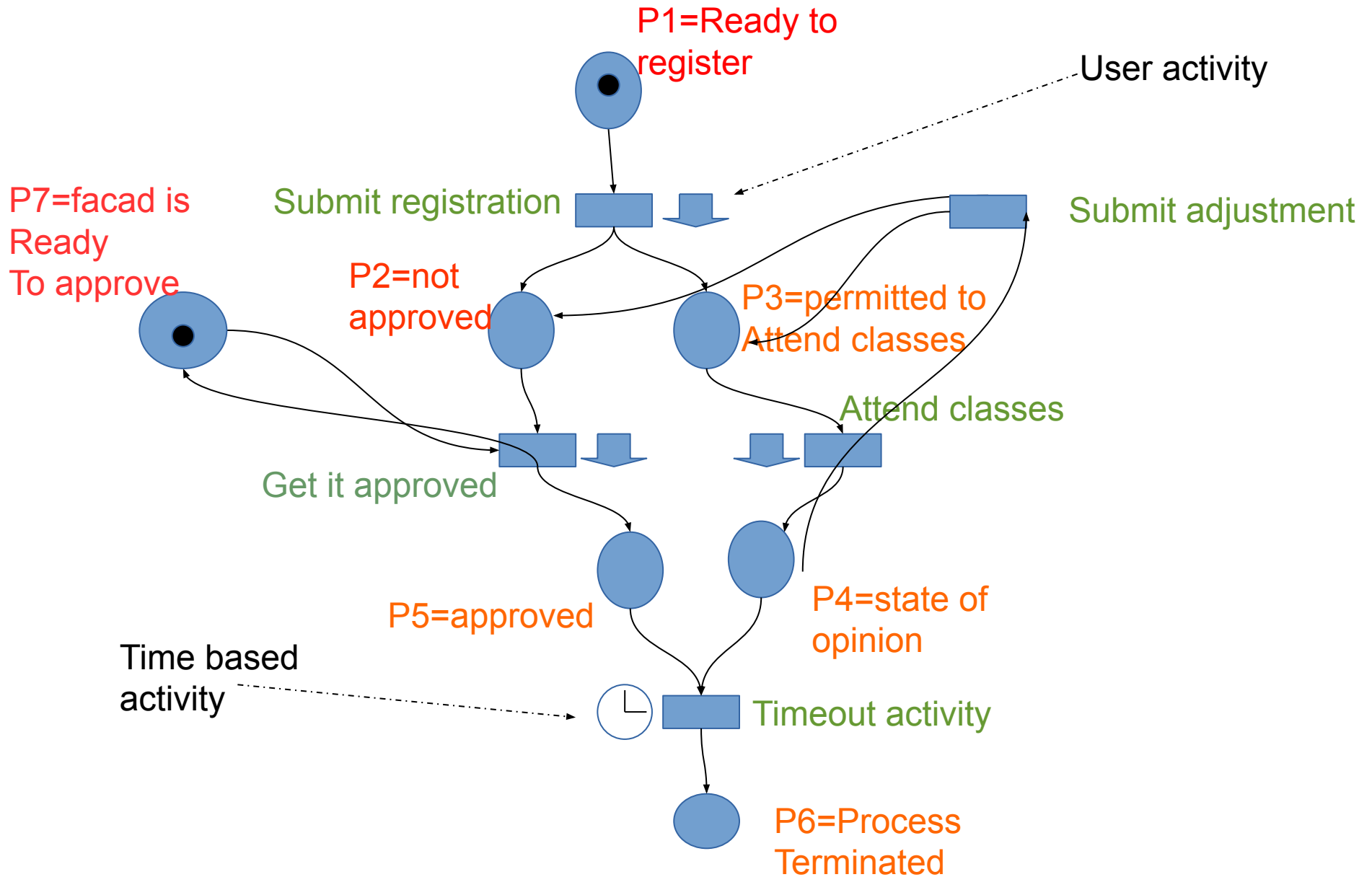
- . Time Triggered Activity

- Task needs to be triggered at a particular time, or after a certain period of timeout

- . (though enabled, time has to trigger it)

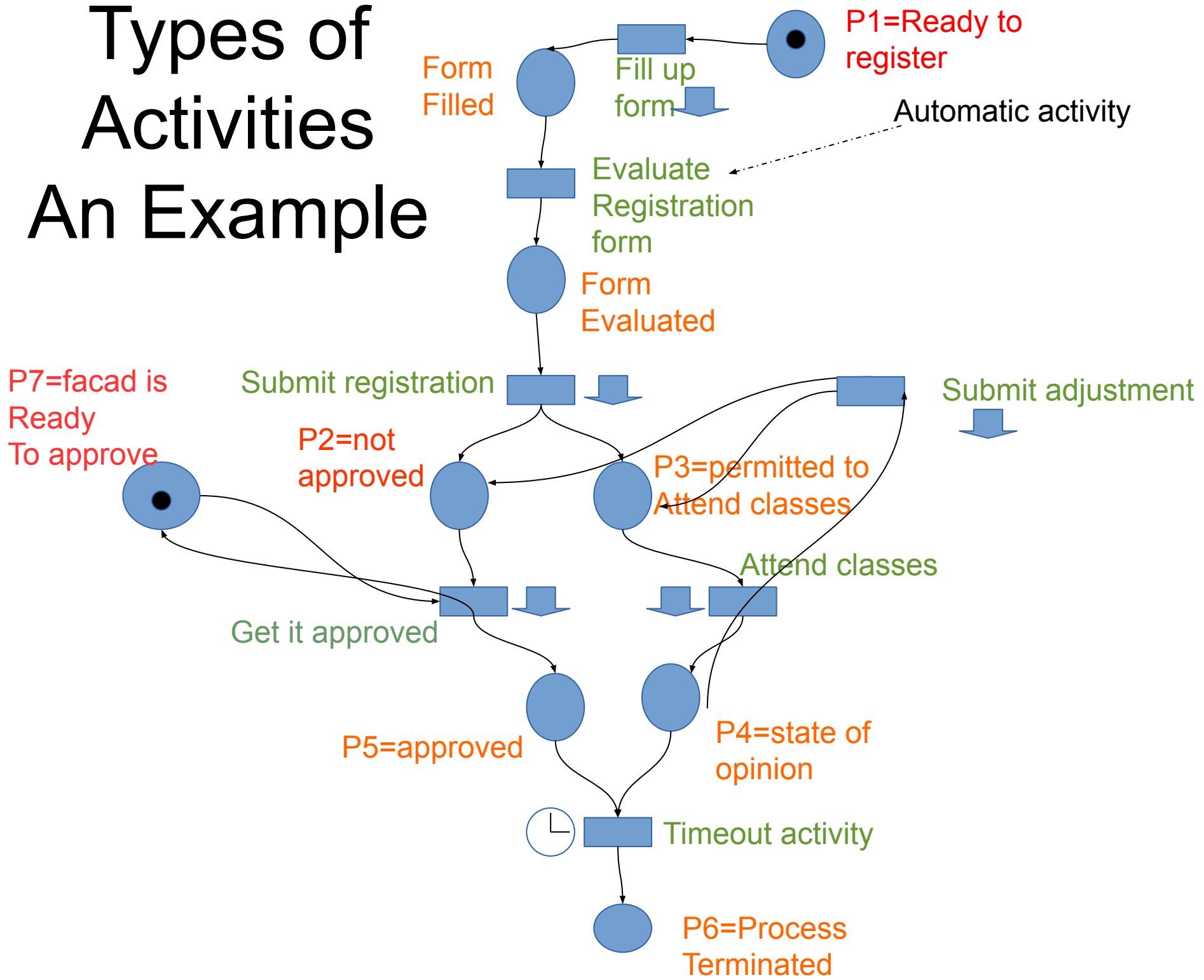
Types of Activities

An Example



Types of Activities

An Example



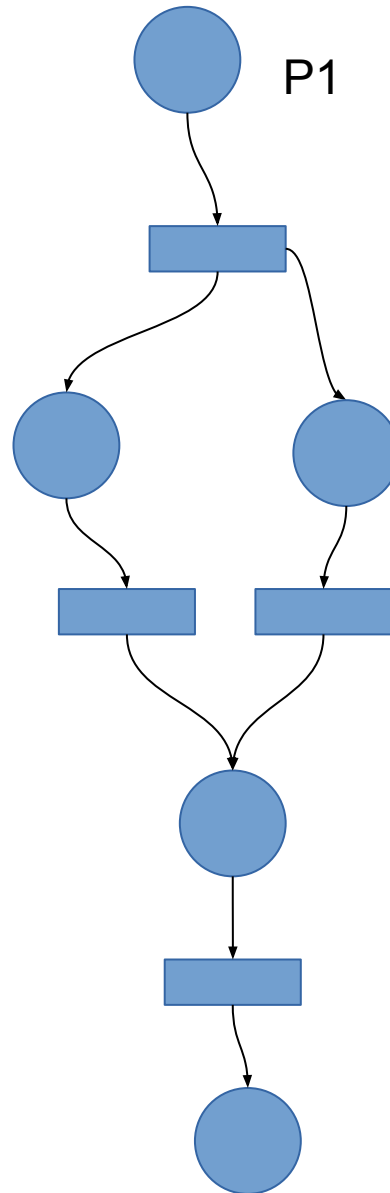
Problem

- Re-engineer the above net to reflect the process that we are actually following
- Make sure all traces that we actually take are in
- Also make sure the traces that we donot take are not in
- i.e. the aim is to get the exact model that is necessary and sufficient to represent the present process.

Classical Petri Nets

- A place can contain 0 or more tokens (more than 1)
 - The state then needs to mention the count of tokens held in places
 - e.g. {1 p1, 2 p3, 4 p4}
 - This state is different from {2 P1, 2 P3, 4 p4}

Places with multiple tokens



Build a state space

Initial marking: {P1}