

Bridging the gap between Course Work and Real Life Problems

Under the guidance of Prof. Milind Sohoni

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IIT Dharwad

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Today...

- What is the current course structure?
 - How to apply the concepts taught in classroom.
 - Connection with reality.
 - How can we enrich our curriculum?
 - Why engagement with society is important?
- Case Studies
 - Analysing data. The census. The railway system. The bus-depot.
- The Dharwad City Bus Depot
 - Various steps and its connection with our curricula.
 - Intermediate and final output.

What is the current course structure

and what is its connection to everyday problems?

- Asymptotic Notation
- Sorting and Searching
- Divide and Conquer
- Greedy Algorithms
- Graph Theory
- Dynamic Programming
- NP-completeness
- and much more.....

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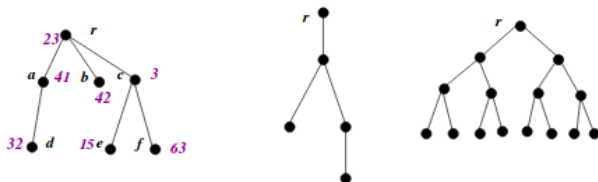
How to adapt them to everyday life?

Real Life Problems are manifold complex than standard textbook problems!

What is being taught?

A rooted tree

Rooted Tree: $T(V, E)$, r is a rooted tree iff (i) T is a tree and $r \in V$ is a vertex (called the **root**, and (ii) there is a function call $level : V \rightarrow \mathbb{Z}$ such that (a) $level(r) = 0$, (b) $level(v) \geq 0$, and (c) each vertex v of level $d > 0$ is connected to exactly one vertex w of level $d - 1$. Then w is called the **parent** of v and v the **child** of w .
A vertex with no children is called a **leaf node**.

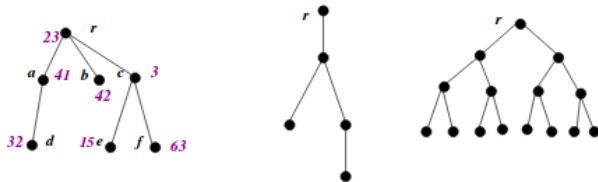


A tree with n vertices has exactly $n - 1$ edges.

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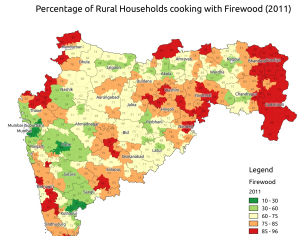
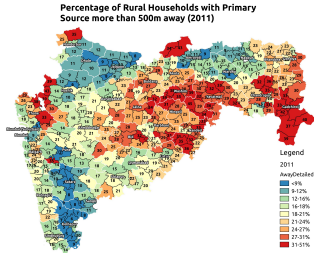
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Too much abstraction!

And what is the reality ...

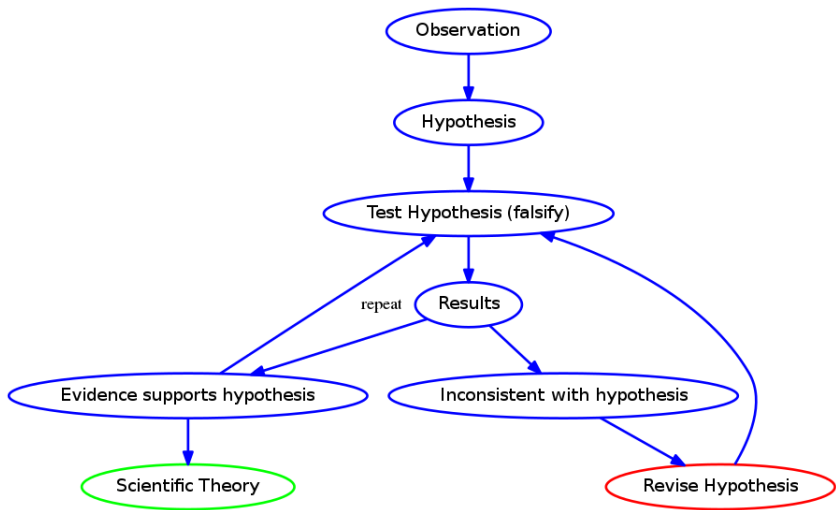


And what is the reality ...



Data from Census 2011

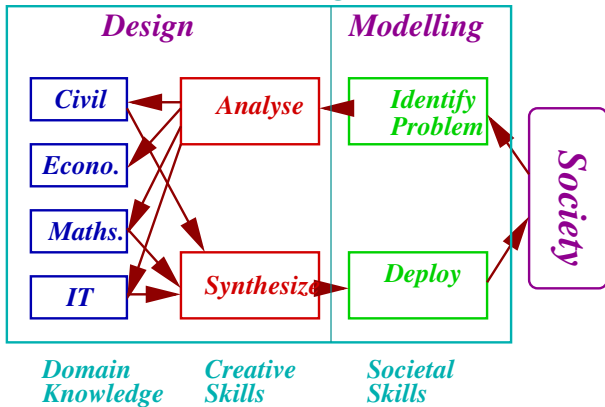
Science



Scientific Method Flow Chart

Engineering

The True Engineer



Engineering v/s Science

- Scientist **describes** society, engineer wants to **change** it. So how should the change be?
- Obviously which generates **value**, which is why we engineers need to **understand** society and how it operates.

Ref: “Engineering Teaching and Research and its Impact on India”, CURRENT SCIENCE, VOL. 102, NO. 11, 10 JUNE 2012 (<https://www.cse.iitb.ac.in/~sohoni/RD.pdf>)

Problem Statement

- How do we analyse a real-life situation?
- How to use our course-material to **model** real-life situations.

Solutions?

Students and Case Studies

- Untapped resources of our country
- Can work in teams with the government
- Benefit of society
- Hands-on experience for students

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Several development issues require good engineering methodology!

The South Western Railway Timetable

Table No.

14



Bengaluru City - Hubballi - Miraj

BG

TRAIN NAME	Kochi- velli Hubballi Express	Yes- vantpur Bikaner Bi Weekly Express	Yes- vantpur Karwar Express via Arsikere	Beng- galuru - Hub- balli Jan- Shatabdi Express	Mysuru - Talgup- pa Town Intercity Express	Puduch- ery - Dadar Express	Tiru- nelveli - Dadar Express	Mysuru -Dadar Sharavati Express	Yesvant- pur Shiva- mogga Express	Yes- vantpur Barnier AC weekly Express	Mysuru- Varanasi Express	Baiyap- ana- halli Tata Ex- press	Hubballi Loka- manyu Tilak (T) Express
Train Number	12778	16587	16515	12079	16206	11006	11022	11036	16579	14805	16229	18112*	17317
Class of accommodation	2A, 3A SL, II	2A, 3A SL, II	CC, 2S, II	CC, 2S	CC, 2S, II	2A, 3A, SL, II	2A, 3A, SL, II	2A, 3A, SL, II	CC, 2S	1A, 2A, 3A	2A, 3A, SL, II	2A, 3A, SL, II	2A, 3A, SL, II
Days of operation from originating station	Thu	Su, F	M, W, F	Daily	Daily	Su, Tu, W	M, Th, F	Su	Sa, Su	M	Tu, Th	Su	Daily
From Table No.	6A				79	1	6A	79			15A		
Bengaluru City Jn.	Km 0	a d		06.00							09.55 10.00		
Malleswaram	3	d											
Yesvantpur Jn.	6	a d	04.30 04.40	05.00 07.00	06.09 06.10	06.00 06.30	04.15 06.30	Mysuru Dept. 06.05 via Has- san	09.00 10.03 10.05	10.30 11.20	10.20 10.22 11.18 11.20	10.30	
Tumakuru	70	a d	05.35	06.03 06.05	07.03 07.05	Mysuru Dep. 05.50	07.30 07.32						
Ammasandra	114	d											
Tiptur	141	d		09.22		08.42	08.42		11.30				
Arsikere Jn.	166	a d	07.15 07.20	07.45 07.50	09.55 10.05	08.30 08.35	08.43 09.10 09.15	09.10 09.15	09.10 12.05 12.10	13.10 13.15	13.10 13.15	13.10 13.15	
Kadur	206	d					09.30	09.50	09.50	12.55	13.50	13.50	
Birur Jn.	211	a d	08.08 08.10		Manga- lore Jn. Arr.	09.13 09.15	10.03 10.05	10.03 10.05	10.03 10.05	13.08 13.10		13.58 14.00	

<https://www.cse.iitb.ac.in/cs213d/SWRTimetable.pdf>

The South Western Railway Timetable

Table No.
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Bengaluru City - Hubballi - Miraj

BG

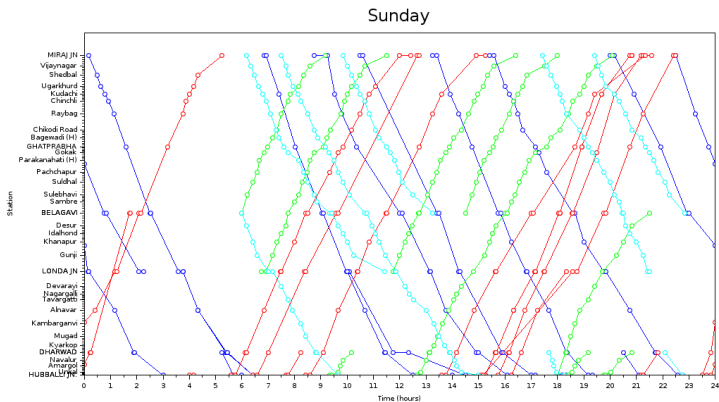
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Tumakuru	70	a d	06.03 06.05	08.03 08.05	07.03 07.05	Mysuru Dep. 05.50	07.30 07.32	07.30 07.32		10.03 10.05	11.20	11.18 11.20	11.20
Ammasandra	114	d											
Tiptur	141	d			09.22		08.42	08.42		11.30			
Anikere Jn.	166	a d	07.15 07.20	07.45 07.50	08.30 10.05	08.43 08.45	09.10 09.15	09.10 09.15	09.10 09.15	13.10 13.15	13.10 13.15	13.10 13.15	13.15
Kadur	206	d					09.50	09.50	09.50	12.55	13.50	13.50	
Birur Jn.	211	a d	08.08 08.10				09.13 09.15	09.33 09.35	10.03 10.05	10.03 10.05	13.08 13.10		13.58 14.00

- Extremely rich data-set.
- Available in pdf format online.
- Cannot run basic computation such as number of trains operational at a given time instant!

From PDF to Excel...

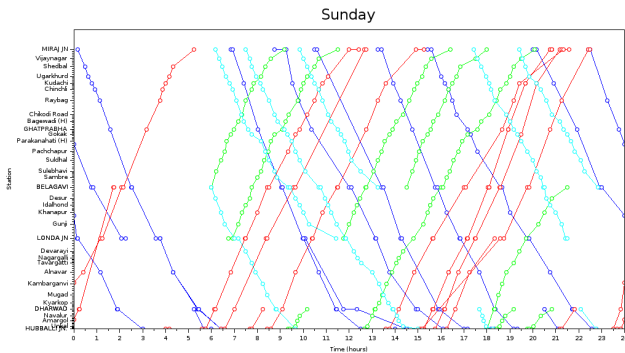
17317	12725	12629	22658	17309	17316	16535	12781	16589	16532	16534	17311	16506	16508
Daily	Daily	Tu, Th	W, Sa	Su, Tu	Tu	Daily	F	Daily	F	Su	F	Sa	M, W
15.45	21.1	21.4	21.4	23.4	23.4	4	4.3	5.35	6.25	6.25	6.35	6.25	6.25
16.08	21.15	21.5	21.5	23.5	23.5	4.1	4.4	5.45	6.35	6.35	6.45	6.35	6.35
16.1	21.5	22.18	22.18	0.13	0.13		5.03	6.08	7	7	7.12		
16.1				0.15	0.15		5.05	6.1				7	7
16.45								6.5					
17.29								7.28	8.23	8.23	8.38	8.23	8.23
17.3								7.3	8.25	8.25	8.4	8.25	8.25
								8					
18.33		0.3	0.3	1.43	1.43		7.15	8.25	9.35	9.35		9.35	9.35
18.35		0.32	0.32	1.45	1.45		7.2	8.3	9.4	9.4		9.4	9.4
								9.2					
19.3								9.4					
								9.5					10.5
								10.1					
								10.3					
								10.4					
20.1								10.5					
								11.05					
21.1		3.15	3.15					12	12.4	12.4		12.4	12.4
21.15		3.2	3.2					12.25	12.45	12.45		12.45	12.45

...To Graphical Visualization



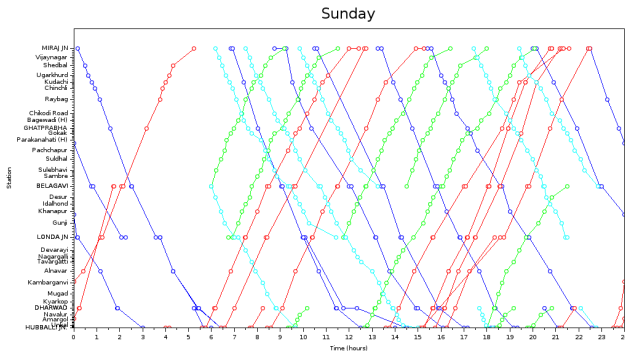
Software Used: Scilab

...To Graphical Visualization



- Space-time representation of timetable
- Can be used to analyze **bottlenecks in schedule**, **minimization of delays**, possible collision domains.
- **Time Table Optimization!**

...To Graphical Visualization



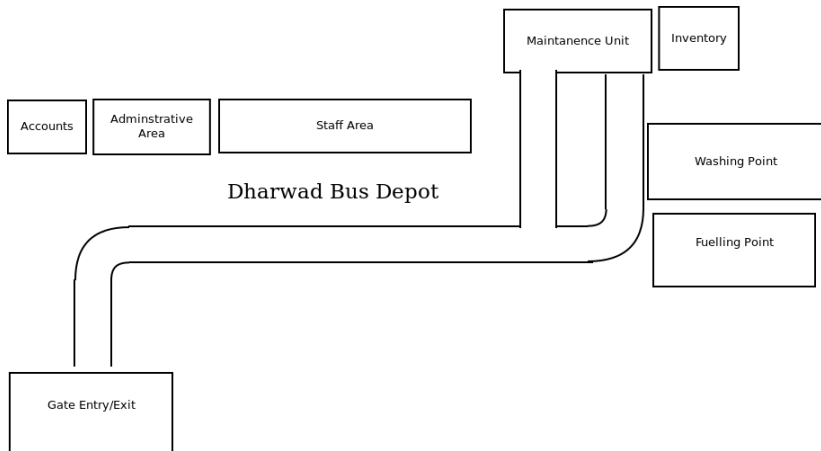
The above example was based on **secondary data** and its analysis. That is important too, since it **trains** you in understanding how to **represent** and what would be of interest to people and the limitations of the implementation agency, in this case, a single track. Also understand what could be changed, e.g., signal spacing and loops, but that needs primary work. **The next example is different.**

Dharwad bus Depot

The Dharwad Bus Depot

- One of the eight division headquarters under North Western Karnataka Road Transport Corporation.
- operates 208 schedules with strength of 228 vehicles.
- covers 55k kms and is utilising the services of 1000 regular employees, which includes Officers, Supervisory and Administrative staff, Mechanical staff and Drivers and Conductors.
- Plethora of natural data generated on a daily basis.

Dharwad Bus Depot: Basic Physical Units



Basic Data-sets

NWKMOWSTAT: 20-Mar-17

NORTH WEST KARNATAK ROAD TRANSPORT CORPORATION
DHARWAD DEPOT DHARWAD

DAILY OPERATIONAL STATISTICS (ST-1) Dt 19-Mar-17

PARTICULARS	CITY	SUBURBAN	MOFFUSIAL	TOTAL	REVENUE
SCH K.M.	85	51	72	208	
CANC K.M.	14327	13872	27230	55429	327180 CITY
EXTRA K.M.	26	0	112	138	253756 SUBURBAN
EFFY K.M.	14301	13872	27118	55291	561929 MOFFUSIAL
TR REVENUE	327180	253756	544609	1125545	1142865 TOTAL
E.P.K.M.	22.88	18.29	20.08	20.36	0 S.C. PASS
Pass Rev	0	0	0	0	1142865 GR TOTAL

REVENUE SUMMARY	VEHICLE SUMMARY	E.P.K.M.
TRAFFIC REVENUE	TOTAL VEHICLE	22.88 CITY
PUBLIC PASS AMT	OFF ROAD	18.29 SUBURBAN
CITY/SUB PASS AMT	VEH AT D.W.S.	20.72 MOFFUSIAL
S.C. PASS	FOR DOC	
W. B. AMOUNT	FOR REPATES	
E.L.T.	V.W./CC/EXT/PAIR	
	LOAN	
GROSS REVENUE	BREAK DOWN	
E.P.K.M.	ACCIDENT	
	NOT ARRIVED	
	ON ROAD	
	SPARE VEHICLE	

ROUTE	EXTRA	C.C.	FAIR	CAUSEWISE CANCELLATION	TRIPS
KILOMETER				SHORTAGE DR	
REVENUE				SHORTAGE CHD	48
E.P.K.M.				SHORTAGE CREW	
				SHORTAGE VEH	
	KM	Rev	FROM	VEH. REPAIR	54
	0	0	#DIV/0!	TYRE PUNCTURE	24
TRIP DETAILS	CITY	SUBURBAN	MOFFUSIAL	BREAK DOWN	
TO BE OPERATED	1232	448	263	STRKE	
CANCELLED	2	1	0	ROAD BLOCK	
OPERATED	1230	445	263	D.A.	
				V.W./CC/EXT	
LATE DEPARTURE				VIATION	
BY 15 MINUTES	DR CHANGE	64	66.00	E.T.M	
BY 30 MINUTES	CR CHANGE	56	71.00	OTHERS	12
BY 45 MINUTE	VEH CHANGE	2	\$9.00	GRAND TOTAL	138
1 HOUR & ABOVE	REGULARITY	2			0.25 %

T.T. A.T.I. CHARGEEN

A.T.S. A.W.S.

DEPOT-MANAGER
N.W.K.R.T.C.
DHARWAD DEPOT

NIPK (DAILY REPORT)

20-Mar-17

NORTH WEST KARNATAKA ROAD TRANSPORT CORPORATION DHARWAD DEPT. DHARWAD

DAILY OPERATIONAL STATISTICS (BT - T) 19-Mar-17

PARTICULARS		CITY	SUBURBAN	HOBBALIA	TOTAL		
SCHEDULES		85	20	12	117		
CONV. K.M.	1457	1380	0	2736	3143		
TR. K.M.	20	20	0	40	40		
DEPT. K.M.	0	0	0	0	0		
TR. REVENUE	15161	1187	0	3320	16368		
D.P.M.	27240	25736	0	52976	52976		
Pass Rev		27.65	18.24	20.31	26.16		
Pass Rev %		100	65.99	73.45	100		
TRAFFIC SUMMARY		VEHICLE SUMMARY					
PUBLIC PASSENGER		TOTAL VEHICLES		228		E.P.K.M	
TRAF. REV. PASSENGER		DEPT. PASSENGER		228		22.85 CITY	
CITY PASSENGER		HOB. PAS. P.		2		20.52 HOBBALIA	
S.B.C. PASSENGER		HOB. PAS. P.		2			
S.B.C. UNIT		HOB. PASSENGER		2			
		HOB. PASSENGER		2			
		HOB. PASSENGER		2			
		HOB. PASSENGER		2			
		HOB. PASSENGER		2			
		HOB. PASSENGER		2			
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		HOB. PASSENGER		2			
		HOB. PASSENGER		2			
		HOB. PASSENGER		2			

- Daily Operation Statistics
- Summary of revenue collection, EPKM (earnings per km), vehicle allocation, cancellations/late departures.

Basic Data-sets

[illegible]

[illegible]

- **Log sheet**
- To be filled in by the conductor after every trip.
- Contains stopwise distribution of ridership.

Basic Data-sets

CC	F M	1 -	07:05	7:20	-	0:00	89	
DCS-3B	19	CITY	CBT	SOMESHWAR TMP	6	14:15	14:35	YAMMIKERE, R.MATH
	20	"	SOMESHWAR TMP	CBT	6	14:40	14:00	YAMMIKERE, R.MATH
	21	"	CBT	TEJASHWI NAGAR	5	15:05	15:20	YAMMIKERE, R.MATH
	22	"	TEJASHWI NAGAR	CBT	5	15:35	15:50	YAMMIKERE, R.MATH
	23	"	CBT	TEJASHWI NAGAR	5	15:55	16:10	YAMMIKERE, R.MATH
	24	"	TEJASHWI NAGAR	CBT	5	16:20	16:35	YAMMIKERE, R.MATH
	25	"	CBT	TEJASHWI NAGAR	5	16:40	16:55	YAMMIKERE, R.MATH
	26	"	TEJASHWI NAGAR	CBT	5	17:10	17:25	YAMMIKERE, R.MATH
	27	"	CBT	SARASWATIPUR	4	17:30	17:45	YAMMIKERE
	28	"	SARASWATHIPUR	CBT	4	17:50	18:05	YAMMIKERE
	29	"	CBT	TEJASHWI NAGAR	5	18:20	18:35	YAMMIKERE, R.MATH
	30	"	TEJASHWI NAGAR	CBT	5	18:40	18:55	REST
	31	"	CBT	TEJASHWI NAGAR	5	19:10	19:25	YAMMIKERE, R.MATH
	32	"	TEJASHWI NAGAR	CBT	5	19:40	19:55	YAMMIKERE, R.MATH
	33	"	CBT	TEJASHWI NAGAR	5	20:10	20:25	YAMMIKERE, R.MATH
	34	"	TEJASHWI NAGAR	CBT	5	20:30	20:45	REST
	35	"	CBT	S.R.NAGAR	5	21:15	21:30	YAMMIKERE, R.MATH
	36	"	S.R.NAGAR	CBT	5	21:35	21:50	YAMMIKERE, R.MATH
N/O CBT		1 -	07:35	6:00	-	0:00	90	179

Form IV(The city bus schedule)

Form IV: Our first Data-set

The Schedule

- SCH.NO.: The Schedule Number (Typical schedule is 8 hours long.)
- TRIP: trip number
 - FROM and TO: The first and last stops.
 - KM: distance covered
 - DEP and ARR : departure and arrival timings
 - VIA: place en route

Very linear representation of data.

- No information about sub-stops.
- Does not describe the spatial distribution of route-data.
- Insufficient for proper analysis.

Questions

It is natural to generate many domain specific questions of value:

- **Passenger**: What is the route to travel from Ganesh Temple to Central School, starting at 7:15 AM, in the minimum possible time?
- **Depot Manager**: How to maximize the profit subject to budget constraints?
- **Planning Committee**: Which are the more profitable areas to add a new bus route?
- **Researcher** : How much percentage of the rural population has access to a bus within 500m?

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We will answer the first question through this study and do preliminary work to answer the other questions.

Work done at IIT Dharwad

The 39 students of the course CS213 (Jan. 2017-May. 2017), at IIT Dharwad map Dharwad city bus routes.

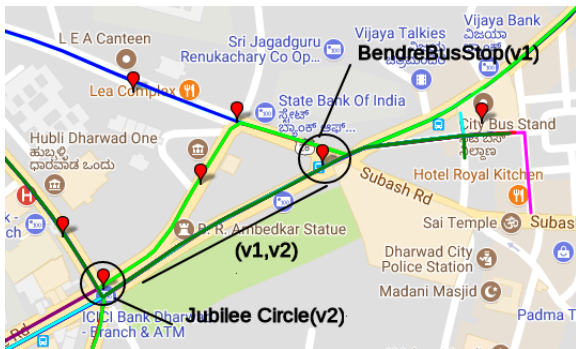
- They first analyze the 8000-line city bus schedule.
- prepare a summary and allocate routes
- travel them and generate kml files
- and finally compile them together.

Summary and Route Allocation

S. No.	Sch. No.	From	To	Km	Via	Name	Roll No.	Status	Route No.
1	DCS-45	CBT	ADARSHA NAGAR	7	KCD, DSK, SDN, ADS	Deeksha, Adhok	32, 33	Done	1
2	DCS 9, DCS-44A, DCS-2A	CBT	AGRI COLLEGE	6	NBS	Shruti, Ketan	02, 31	Done	5
3	DCS 37A, DCS 28A	CBT	ANJANEY NAGAR	7	KCD, DSKP	Yeshwanth, Niteesh, Nikshep	08, 21, 23	Done	37
4	DCS-50, DCS-11A	CBT	BANASHREE NAGAR	7	KCD, DSK, SDN, ADS	Riya, Shreyas, Prabhakar	07, 10, 9	Done	4
5	DCS 34A	CBT	CHAITANYA NAGAR	5		Yeshwanth, Niteesh, Nikshep	08, 21, 23	Done	34
6	DCS-46	CBT	DHARAWAD H.C	15	NBS, MMT	Ashish, Pranay	08, 30	Done	18
7	DCS-36A, DCS 56A, DCS 57A, D	CBT	DRS	3	YMK	Shruti, Ketan	02, 31	Done	11
8	DCS 62A	AGRI COLLEGE	DRS	9				Too early Morning	-
9	DCS-26A	CBT	GAMANAGATTI	14	TOLL NAKA,	Riya, Shreyas, Prabhakar	07, 10, 9	Done	9
10	DCS-61, DCS-23B	CBT	HANUMANT NAGAR	7		Harshal, Ritwan	3, 15	Done	13A
11	DCS 7	CBT	Itigatti	12	YARIKOPPA	Shruti, Ketan	02, 31	Done	7
12	DCS-24A, DCS-29	CBT	Jogellappur	10		Vishnu, Manoj	14, 19	Done	18
13	DCS-5A, DCS-23A, DCS-61	CBT	KALYAN NAGAR	7	Barakotri	charu, prashanth	1, 38	Done	23
14	DCS 53	CBT	KAYAKANAGAR	11	NBS	Ishan, Akash	13, 16	Discontinued	-
15	DCS-14A, DCS-38, DCS-2A, DCS-9	CBT	MADHIAL	3		Ashish, Pranay	08, 30	Done	14
16	DCS-3A, DCS 60A	CBT	MAHIBUB NAGAR	3		Ashish, Pranay	08, 30	Done	60
17	DCS-12	CBT	MALAPUR	3	SIVAJI CIRCLE	Tirumal, Bhanu	23, 24	Done	34
18	DCS 1A, DCS-32	CBT	NAVALUR	7	TOLL NAKA	Yeshwanth, Niteesh, Nikshep	08, 21, 23	Done	1
19	DCS-22, DCS-44A, DCS-65	CBT	NBS	2.5	M.HSP	Ishan, Akash	13, 16	Done	75
20	DCS-52, DCS-59, DCS-51A, DCS-40A	CBT	NEHARU NAGAR	6	KCD, SFR	Deeksha, Adhok	32, 33	Done	1
21	DCS-6A	Railway Station	NEW BUS STAND	5.5	CBT			Done	-
22	DCS-16, DCS-21A, DCS-31A, D	CBT	PAVATE NAGAR	5	KCD	Vishnu, Manoj	14, 19	Done	68
23	DCS-41A, DCS-18A	MALAPUR	R.G NAGAR	10.5		Tirumal, Bhanu	23, 24	Done	-
24	DCS-41B	CBT	R.G.NAGAR	8		Tirumal, Bhanu	23, 24	Done	0
25	DCS 39A, DCS-48	CBT	R.K.NAGAR	6		Deeksha, Adhok	32, 33	Done	2
26	DCS-15	CBT	Railway Station	3	CC	Sanjay, Varsha	22, 27	Done	11
27	DCS 19A, DCS-20A, DCS 58A	CBT	Ramanagar	6	YMK	Ishan, Akash	13, 16	Done	98
28	DCS-10, DCS-33A, DCS-63A	CBT	S.R.NAGAR	5	TOLL NAKA	Riya, Shreyas, Prabhakar	07, 10, 9	Done	28
29	DCS-35A	CBT	SATHANAKOTTA	4		Kannath Paul Ashish	04, 05, 17	Done	30

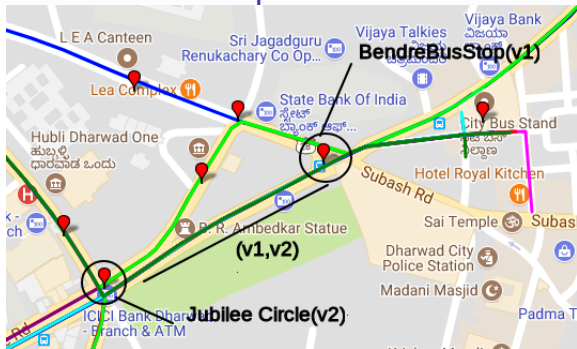
75 distinct routes were allocated to the students, out of which 12 were reported to be dysfunctional or duplicate.

The Representation



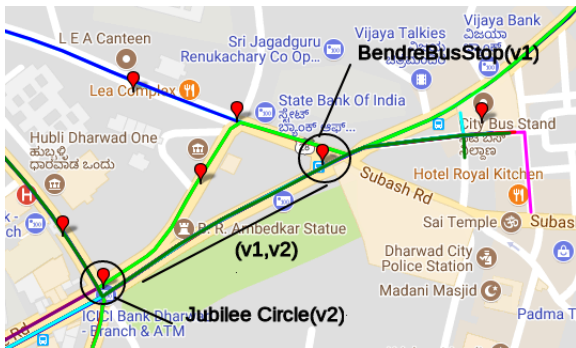
A graph $G(V, E)$ is eminently suitable to represent locations (as vertices) and paths between locations (as edges).

The Representation



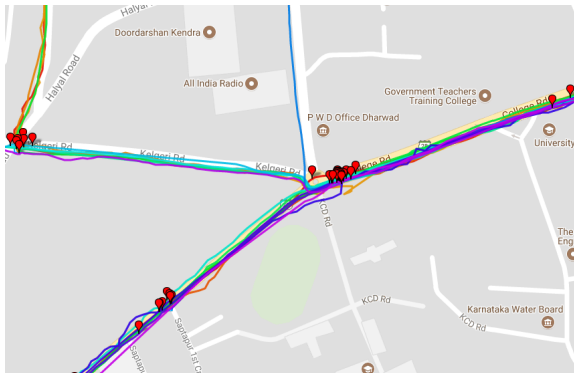
- The **place** (v) represents an actual location and contains location data (latitude-longitude (lat-long)), name and other attributes specific to the place and the route on which it lies.
- An **edge** consists of two places (v_1, v_2) and a track between them.
- A **route** consists of a sequence of places, connected by the edges.

The Representation



- A **service** from v_i to v_j , $S(v_i, v_j, t_0, d)$, consists of a route followed by a bus along with the start time(t_0) and the trip duration (d).
- Finally, a **schedule** is a set of services carried out by a single bus.

Figure: The KML tracks generated from real-time GPS tracking.



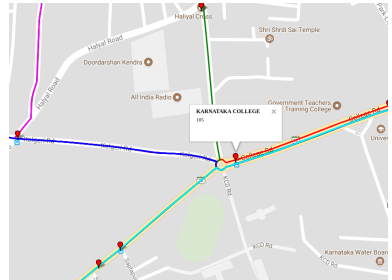
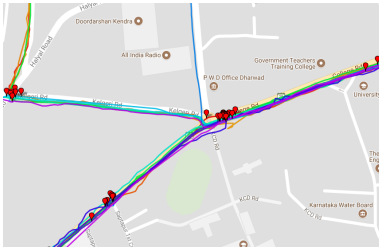
Some Issues

- The KML tracks are irregular and not aligned to official road polylines.
- Multiple coordinates and multiple names for the same stops.
- Names in the schedule are not standardized.
- Data not fit for graph representation.

Cleanup

- **Clustering of Stops**

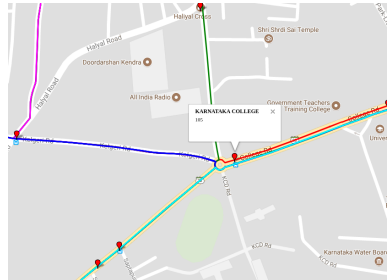
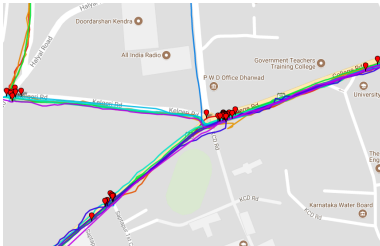
We used a clustering algorithm (see breadth-first search) using connected components to cluster stops within 50 m into unique stops.



Cleanup

- Fitting the GPS generated tracks with the Google Map Roads

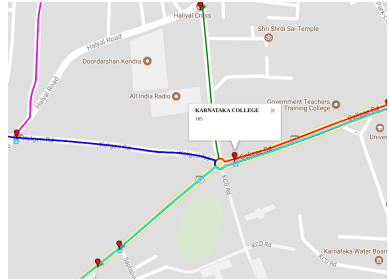
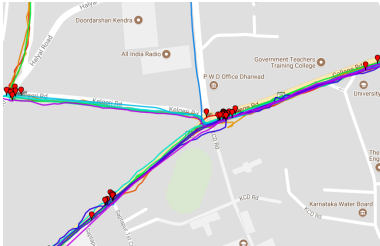
It was decided that the edges between places (whose preliminary lat-longs and names are now available) will follow existing roads, i.e., polylines as shown by Google Maps. To align the tracks with the official road polylines, we used a Google API



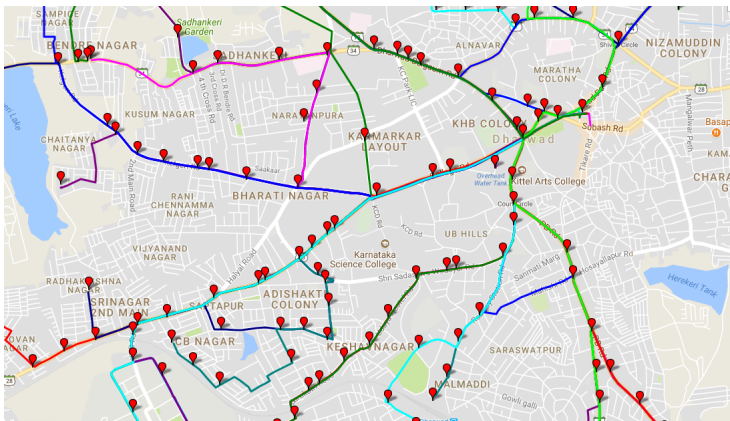
Cleanup

- Inserting stops into polylines

As the data was generated from one trip only, many buses did not stop at all the stops which lie on the route depending on the ridership that day. So, we then wrote a program which took all the stops, computed its minimum distance from the track and inserted it into the track O_i if the distance was less than some chosen epsilon.



After Cleanup



Question Revisited

Question: How can you reach Srinagar from High Court in the minimum possible time starting at time t_0 ?

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This can be modelled as a **graph theory** problem.

- Here the **vertices**(V) are the bus stops.
- Let v_0 = Srinagar and v_n =High Court.
- There **edges** are the services.
- The **edge weights** are the trip length of the service $S(v_i, v_j, t, d)$.

Question Revisited

- **Initialization:** The edge cost of $S(v_0, v_i, t, d) = (t - t_0) + d$ where $(t - t_0) > 0, \infty$ otherwise.
- **Constraint:** (v_i, v_j) can be followed by (v_j, v_k) iff arrival time($S(v_i, v_j, t_1, d_1)$) < departure time($S(v_j, v_k, t_2, d_2)$).
- Run **Dijkstra's algorithm** and get the shortest path from v_0 to v_n .

Ref: https://en.wikipedia.org/wiki/Dijkstra's_algorithm

Can we do better?

- Suppose there are k routes with n stops each and 1 route has m trips in a day, the number of edges becomes $(n - 1) \times k \times m$.
- So the complexity of above algorithm is $O((n - 1) \times k \times m + |V| \log |V|)$!
- This is **significant amount of computation** if done on the average mobile phone processors.

Can we do better?

- **Solution:** Edge generation on demand.
- Note that we are only interested in the **next bus**.
- So instead of explicitly storing the edges for all buses, we simply compute the cost of next bus from the current time.
- This brings the complexity of the algorithm by a multiplicative factor of m : **Significant Improvement**

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Our algorithm is ready to be implemented on a mobile phone!

Algorithm 1 Algorithm for finding bus route

```
procedure ROUTEPLANNER(Graph, source,
destination, start time)
  create vertex set Q
  for each vertex v in Graph do           ▷ Initialization
     $dist[v] \leftarrow \infty$  ▷ Unknown distance from source
  to v
     $prev[v] \leftarrow UNDEFINED$  ▷ Previous edge in
  optimal path from source
     $arrival\_time[v] \leftarrow \infty$  ▷ Stores the best arrival
  time at each vertex
  end for
  add v to Q      ▷ All nodes initially in Q (unvisited
  nodes)
   $dist[source] \leftarrow 0$  ▷ Distance from source to source
   $arrival\_time[source] \leftarrow start\ time$ 
  while Q is not empty do
     $u \leftarrow$  vertex in Q with min  $dist[u]$  ▷ Node with
  the least distance will be selected first
    remove u from Q
    if u is destination then
      break           ▷ no need of further computation
    end if
```

```
    for each edge starting from u do ▷ where v is still in Q
       $length(u, v) \leftarrow$ 
       $waiting\_time(arrival\_time[u]) + length(u, v)$            ▷ re
    compute edge cost
       $alt \leftarrow dist[u] + length(u, v)$ 
      if  $alt < dist[v]$  then           ▷ A shorter path to v has
    been found
         $dist[v] \leftarrow alt$ 
         $prev[v] \leftarrow e(u, v)$ 
         $arrival\_time[v] \leftarrow arrival\_time[u] + length(u, v)$ 
      end if
    end for

  return  $dist[ ], prev[ ]$ 
```

Development of an android app for passengers

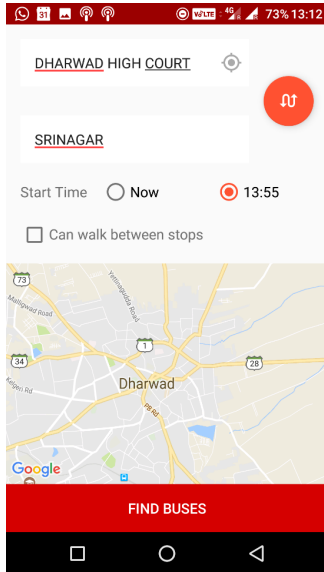


Figure: Route generated by shortest path algorithm

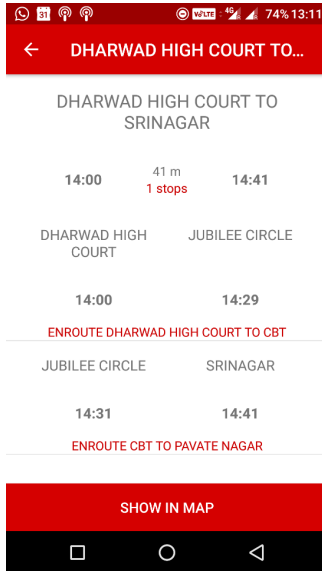
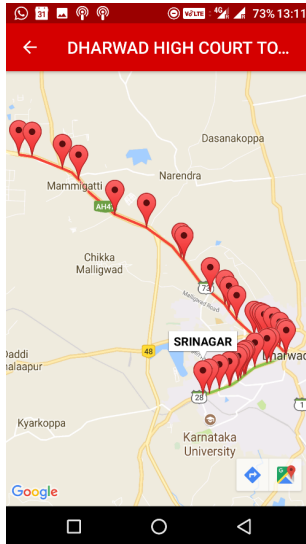


Figure: Route viewed on map



GIS

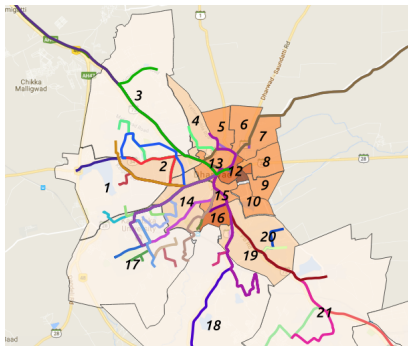


Figure: Ward map of Dharwad (classified by population density) superimposed on the bus network.

The way ahead...

- Implement a similar strategy in your town/taluka to generate a database.
- Scale the database over multiple talukas.
- Encourage addition of **development** course projects into the curriculum.
- There are other fields that need work also:
Energy sector, public transport, more water, town planning,...
- See : www.ctara.iitb.ac.in for project ideas.

For full report on Dharwad Bus Depot:

<https://www.cse.iitb.ac.in/~sohoni/dharwadbus.pdf>

Thanks

