## **Friends of Friends**

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A well-known 'paradox' in (social) networks is informally stated as "your friends are more popular than you are, on average". A first step to make this eye-catching statement more rigorous is to restate it as "the average degree of a randomly selected node in a network is less than the average degree of neighbors of a randomly selected node", but this leaves unspecified the exact mechanism of averaging. There are at least two interpretations, one well-known, the other, apparently not.

Given an undirected graph G = (V, E) with |V| = N and |E| = M and no isolated nodes.

The set of neighbors of node u is nbr(u).

The average degree is  $\mu = \frac{1}{N} \sum_{u \in V} |\operatorname{nbr}(u)| = \frac{2M}{N} \ge 1.$ 

Let the number of "friends of friends" of node u be denoted  $FF(u) = \sum_{v \in nbr(u)} |nbr(v)|$ .

We have  $FF(u) \ge |\operatorname{nbr}(u)| \ge 1$ .

Consider the quantities

$$\operatorname{MicroAvg} = \frac{\sum_{u \in V} \operatorname{FF}(u)}{\sum_{u \in V} |\operatorname{nbr}(u)|} \quad \text{and} \quad \operatorname{MacroAvg} = \frac{1}{N} \sum_{u \in V} \frac{\operatorname{FF}(u)}{|\operatorname{nbr}(u)|}$$

Claim 1. MicroAvg  $\geq \mu$ .

Consider the numerator  $\sum_{u \in V} FF(u) = \sum_{u \in V} \sum_{v \in nbr(u)} |nbr(v)|$ . This is effectively an enumeration over edges, accumulating on both incident nodes:

- 1: for each node  $u \in V$  do
- 2: initialize  $FF(u) \leftarrow 0$
- 3: for each edge  $\{u, v\} \in E$  do
- 4:  $FF(u) \leftarrow FF(u) + |\operatorname{nbr}(v)|$
- 5:  $FF(v) \leftarrow FF(v) + |\operatorname{nbr}(u)|$

Note that v contributes the term  $|\operatorname{nbr}(v)|$  to each of its  $|\operatorname{nbr}(v)|$  neighbors, which then add up in  $\sum_{u \in V} \operatorname{FF}(u)$ . Thus the total contribution of v to the numerator is  $|\operatorname{nbr}(v)|^2$ . Thus we get

$$MicroAvg = \frac{\sum_{u \in V} |\operatorname{nbr}(u)|^2}{\sum_{u \in V} |\operatorname{nbr}(u)|} = \frac{\sum_{u \in V} |\operatorname{nbr}(u)|^2}{2M} = \frac{\sum_{u \in V} |\operatorname{nbr}(u)|^2}{N} \frac{N}{2M}$$

The standard deviation of node degree is defined as

$$\sigma^2 = \frac{\sum_{u \in V} |\operatorname{nbr}(u)|^2}{N} - \mu^2 \quad \Longrightarrow \quad \frac{\sum_{u \in V} |\operatorname{nbr}(u)|^2}{N} = \sigma^2 + \mu^2$$

from which we get

MicroAvg = 
$$\frac{\sigma^2 + \mu^2}{\mu} = \frac{\sigma^2}{\mu} + \mu \ge \mu.$$

Claim 2. MacroAvg  $\geq \mu$ .

Here the computation has to be organized slightly differently.

1: for each node  $u \in V$  do 2: initialize  $Q(u) \leftarrow 0$ 3: for each edge  $\{u, v\} \in E$  do 4:  $Q(u) \leftarrow Q(u) + \frac{|\operatorname{nbr}(v)|}{|\operatorname{nbr}(u)|}$ 5:  $Q(v) \leftarrow Q(v) + \frac{|\operatorname{nbr}(u)|}{|\operatorname{nbr}(v)|}$ return  $\frac{1}{N} \sum_{u \in V} Q(u)$ 

So this time, each edge contributes to MacroAvg the quantity  $\frac{|\operatorname{nbr}(v)|}{|\operatorname{nbr}(u)|} + \frac{|\operatorname{nbr}(u)|}{|\operatorname{nbr}(v)|} \ge 2$ , because  $\min_{a,b>0} \frac{a}{b} + \frac{b}{a} = 2$ . We thus get

MacroAvg = 
$$\frac{1}{N} \sum_{u \in V} Q(u) \ge \frac{1}{N} \cdot M \cdot 2 = \frac{2M}{N} = \mu.$$

The remaining question is whether some inequality holds between MacroAvg and MicroAvg.

Claim 3. MicroAvg > MacroAvg is possible.

Consider a 4-node clique on nodes 1, 2, 3, 4, followed by a chain 4, 5, 6.

| Node $u$ | $\operatorname{nbr}(u)$  | $\mathrm{FF}(u)$                     | $ \operatorname{FF}(u)/ \operatorname{nbr}(u) $                  |
|----------|--------------------------|--------------------------------------|--|
| 1        | $2,\!3,\!4$              | 3 + 3 + 4 = 10                       | 10/3 = 3.3   |
| 2        | $1,\!3,\!4$              | 3 + 3 + 4 = 10                       | $10/3 = 3.\dot{3}$   |
| 3        | $1,\!2,\!4$              | 3 + 3 + 4 = 10                       | $10/3 = 3.\dot{3}$   |
| 4        | 1,2,3,5                  | 3 + 3 + 3 + 2 = 11                   | 11/4 = 2.75  |
| 5        | 4,6                      | 4 + 1 = 5                            | 5/2 = 2.5  |
| 6        | 5                        | 2                                    | 2/1 = 2  |
|          | $\mu = 16/6 = 2.\dot{6}$ | $\sum_{u} \operatorname{FF}(u) = 48$ | $\sum_{u  \frac{\text{FF}(u)}{ \operatorname{nbr}(u) }} = 17.25$ |

MacroAvg =  $\frac{1}{N} \sum_{u} \frac{FF(u)}{|nbr(u)|} = \frac{17.25}{6} = 2.875 < MicroAvg = \frac{48}{16} = 3$ 

Claim 4. MicroAvg < MacroAvg is possible.

Consider the square on nodes 1, 2, 3, 4 with diagonal 2, 4 connected.

| Node $u$ | $\operatorname{nbr}(u)$ | $\mathrm{FF}(u)$                     | $\operatorname{FF}(u)/ \operatorname{nbr}(u) $                         |
|----------|-------------------------|--------------------------------------|--|
| 1        | 2,4                     | 3+3=6                                | 6/2=3  |
| 2        | $1,\!3,\!4$             | 2+2+3=7                              | $7/3 = 2.\dot{3}$  |
| 3        | 2,4                     | 3+3=6                                | 6/2=3  |
| 4        | 1,2,3                   | 2+3+2=7                              | $7/3 = 2.\dot{3}$  |
|          | $\mu = 10/4 = 2.5$      | $\sum_{u} \operatorname{FF}(u) = 26$ | $\sum_{u} \frac{\mathrm{FF}(u)}{ \operatorname{nbr}(u) } = 10.\dot{6}$ |

 $MacroAvg = \frac{1}{N} \sum_{u} \frac{FF(u)}{|\operatorname{nbr}(u)|} = \frac{10.\dot{6}}{4} = 2.\dot{6} > \operatorname{MicroAvg} = \frac{26}{10} = 2.6$