

# Learning to Rank Networked Entities

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[www.cse.iitb.ac.in/~soumen/doc/netrank](http://www.cse.iitb.ac.in/~soumen/doc/netrank)

## Learning to rank...

- ...feature vectors, studied in detail
  - $i^{\text{th}}$  entity represented by feature vector  $x_i$
  - Score of  $i^{\text{th}}$  entity is dot product  $\beta' x_i$
  - Want  $\beta' x_i \leq \beta' x_j$  if we say “ $i < j$ ”
  - Max-margin setup
$$\min_{\beta \in \mathbb{R}^d} \beta' \beta \text{ subject to } \beta' x_i + 1 \leq \beta' x_j \text{ for all } i < j$$
  - Other scores, e.g. 2-layer neural net (RankNet)
- ...nodes in a graph, less so
  - Strongly motivated by Pagerank and HITS
  - Changing score of one node influences others

## Edge conductance and Pagerank

- Conductance of edge  $i \rightarrow j$  written as  $C(j,i)$ 
  - $C(j,i) = \Pr(j \text{ @ this step} \mid i \text{ @ previous step})$
  - Pagerank vector  $p$  satisfies  $p = C p$

- Unweighted (standard) Pagerank

$$C(j,i) = \begin{cases} \alpha \frac{[(i,j) \in E]}{\text{OutDegree}(i)} + (1-\alpha)r_j & i \in V_o \\ r_j & \text{otherwise} \end{cases}$$

$i$  is a dead-end  
 $\sum_{j \in V} r_j = 1$

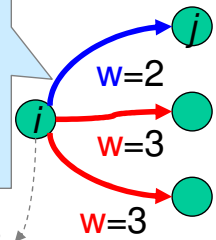
- Weighted Pagerank:  $i \rightarrow j$  edge weight  $w(i,j)$

$$C(j,i) = \begin{cases} \alpha \frac{w(i,j)}{\sum_{j'} w(i,j')} + (1-\alpha)r_j & i \in V_o \\ r_j & \text{otherwise} \end{cases}$$

Pr(teleport)

Prob. of following this edge  
 $2/(2+3+3)$

Teleport?



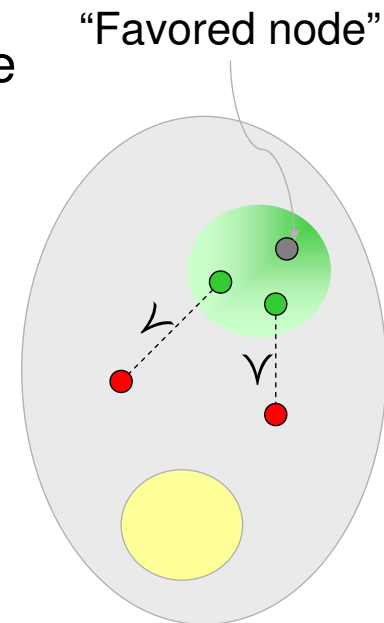
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## Inverse problem

- Traditionally: Given matrix  $C$ , find Pagerank
- Clever design of  $C$  for various applications
  - Hand-tweak teleport vector  $r$  – topic sensitive (Haveliwala), personalized (Jeh+Widom)
  - Hand-tweak  $w(i,j)$  (Intelligent Surfer, ObjectRank)
- Our problem: Given partial order  $<_{\text{train}}$ , find  $C$  (and  $p$ ) such that
  - $p$  satisfies  $p = C p$  approximately
  - $p$  satisfies  $<_{\text{train}}$  and unseen  $<_{\text{test}}$  well: i.e.,  $p_i \leq p_j$  if  $i < j$
  - $<_{\text{train}}$  and  $<_{\text{test}}$  comes from same “hypothesis”

## Preferred community scenario

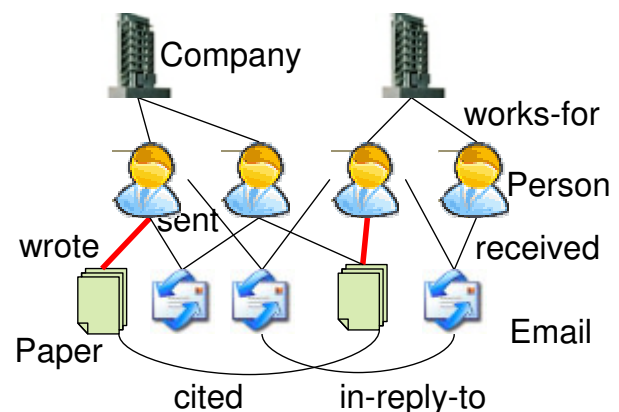
- Ranking papers for Data Mining researcher
- Some subgraphs and citations more important than others
- Revealed via pairwise preferences
- Do not estimate  $C(j,i)$  directly
- Directly estimate  $p_{ij}$ , a constrained “flow” from  $i$  to  $j$
- “BTW”  $C(j,i) = \frac{p_{ij}}{\sum_{(k,i) \in E} p_{ki}}$  Inflow into  $i$
- Local “transductive” setting
- Lots of parameters



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## Entity-relationship graph scenario

- Many node and edge types
- Edge  $e$  has type  $t(e) \in \{1, \dots, T\}$
- Weight  $w(i,j) = \beta(t(i,j))$
- Find  $\beta(1), \beta(2), \dots, \beta(T)$  for least violation
- “Global entanglement” but far fewer parameters
- Somewhat “inductive”, can augment graph with objects of known types



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# 1: The constrained flow formulation

$p_{uv}$  is interpreted as a flow along  $(u, v)$

$\min_{\substack{\{0 \leq p_{uv} \leq 1 : (u,v) \in E'\} \\ \{0 \leq s_{uv} : u < v\}}} \sum_{(u,v) \in E'} p_{uv} \log \frac{p_{uv}}{q_{uv}} + B \sum_{u < v} s_{uv}$

$\{q_{uv}\}$  is a parsimonious flow (unweighted Pagerank)

$s_{uv}$  is the slack or violation on preference  $u < v$

$\sum_{(u,v) \in E'} p_{uv} = 1$

Balance

$\forall v \in V': \sum_{(u,v) \in E'} p_{uv} = \sum_{(v,w) \in E'} p_{vw}$

Dummy node implements teleport

Teleport

$\forall v \in V_o: \frac{p_{vd}}{(1-\alpha)} = \frac{\sum_{(v,w) \in E'} p_{vw}}{\alpha}$

Preference

$\forall u < v: (1+\epsilon) \sum_{(w,u) \in E'} p_{wu} \leq s_{uv} + \sum_{(w,v) \in E'} p_{wv}$

$\alpha \times$  inflow

$(1-\alpha) \times$  inflow

## The dual optimization

- $O(2|V|+|<|)$  dual variables
  - Unconstrained  $\beta_v$  for balance,  $\tau_v$  for teleport
  - $0 \leq \pi_{uv} \leq B$  for preference

- Primal flows in familiar log-linear form

$$\forall v \in V \quad p_{dv} = (1/Z) q_{dv} \exp(\beta_v - \beta_d + \text{bias}(v))$$

$$\forall v \in V_o \quad p_{vd} = (1/Z) q_{vd} \exp(\beta_d - \beta_v + \alpha \tau_v)$$

$$\forall v \in V \setminus V_o \quad p_{vd} = (1/Z) q_{vd} \exp(\beta_d - \beta_v)$$

$$\forall (u,v) \in E \quad p_{uv} = (1/Z) q_{uv} \exp(\beta_v - \beta_u - (1-\alpha)\tau_u + \text{bias}(v))$$

- Where  $\text{bias}_\epsilon(v) = \sum_{r < v} \pi_{rv} - (1+\epsilon) \sum_{v < s} \pi_{vs}$

- Dual objective: minimize  $\log Z$

- Can include dual vars gradually

Large bias  $\Rightarrow$   
large flow into  
 $v \Rightarrow$  high rank

# Competition: Teleport learning via QP

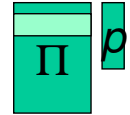
- Let  $A$  be node adjacency matrix with no dead ends and rows scaled to sum to 1

$$p = \alpha A' p + (1 - \alpha)r, \quad \therefore p = (1 - \alpha)(\mathbb{I} - \alpha A')^{-1} r \triangleq M r$$

Fit this

- Preference  $\prec$  expressed as  $\Pi p = \Pi M r \geq 0$

$$\underbrace{(0, \dots, -1, \dots, 1, \dots, 0)}_{\text{Row of } \Pi} \underbrace{(0, \dots, p_i, \dots, p_j, \dots, 0)}_{\text{Column vector } p} \geq 0$$



- Let  $r^U$  be the parsimonious uniform teleport

Deviate from unweighted Pagerank as little as possible...

$$\min_{r \in \mathbb{R}^{|V|}, s \geq 0} (Mr - Mr^U)'(Mr - Mr^U) + B\mathbf{1}'s$$

s.t.  $r \geq 0, \mathbf{1}'r = 1, \Pi Mr + s \geq 0$

Makes QP very expensive

...while satisfying  $\prec$

## Data set preparation

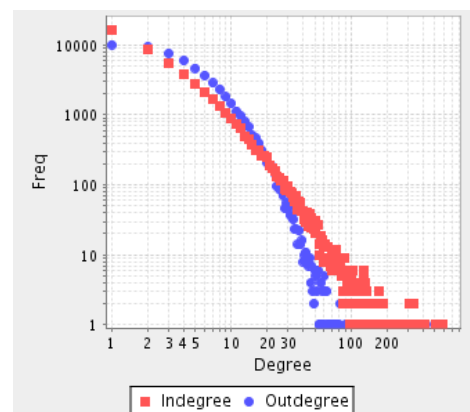
- No open benchmark for this task
  - ☹️ No standardized comparison yet
  - 😊 We will make code and some data available
  - 😊 Synthetic  $G$  and  $\prec$  can explore space thoroughly

- Generating graph  $G$

- RMAT (power-law degree, small dia, clustering)
- Real DBLP+CiteSeer graph

- Generating preference  $\prec$

- Use  $r_{\text{hidden}}$  to compute  $p_{\text{hidden}}$
- Sample  $\prec_{\text{train}}, \prec_{\text{test}}$  from  $p_{\text{hidden}}$
- Measure flips on  $\prec_{\text{test}}$



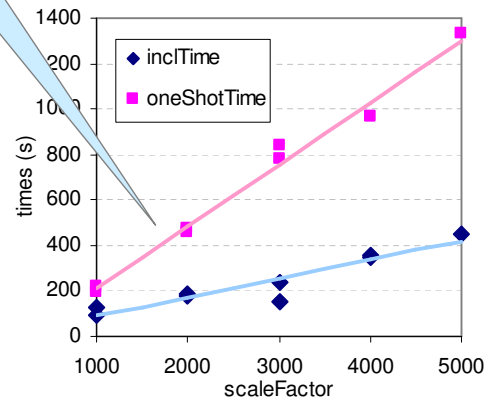
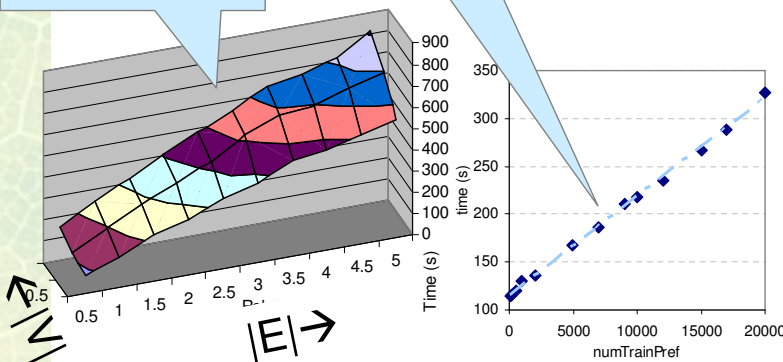
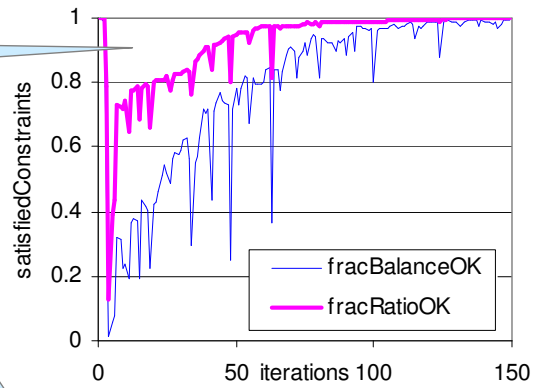
# Optimization dynamics and performance

Gradual jittery decrease in primal violations

Time scales linearly with  $|\mathcal{K}_{\text{train}}|$

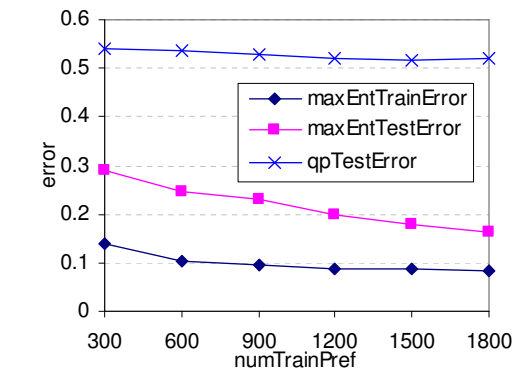
Time scales linearly with  $|V|$  and  $|E|$

Gradual dual var inclusion speedup

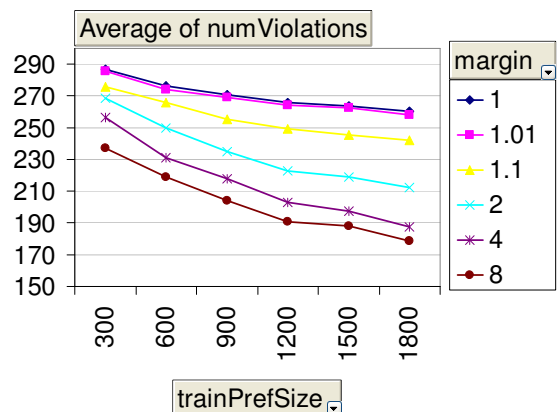


# Learning rate and effect of margin

- Without  $\|r\|_1 = 1$  constraint QP fails to learn from  $\prec$
- Enforcing  $\|r\|_1 = 1$  improves learning
- $|V| \times |V|$  inversion impractical, QP slow
  - Days vs. minutes
- Flow formulation with margin is much better
- Margin needs tuning, not scale-insensitive

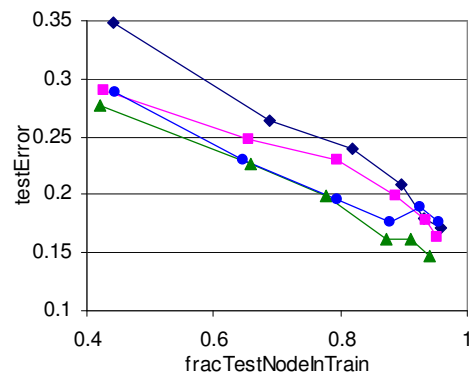
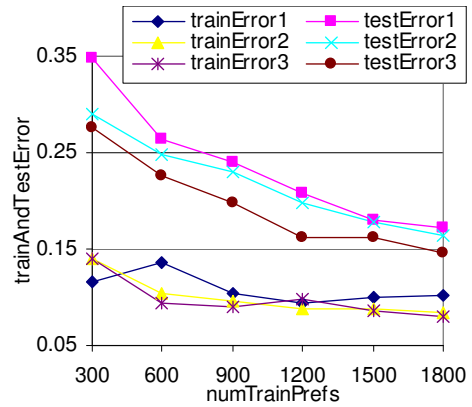


mode test\_ cee 10\_



## Effect of node overlap

- Nodes involved in  $\prec$   
 $V(\prec) = \{w : w \prec v \text{ or } u \prec w\}$
- What if  $V(\prec_{\text{train}})$  and  $V(\prec_{\text{test}})$  overlap?
- Note,  $\prec_{\text{train}}$  and  $\prec_{\text{test}}$  do not overlap!
- Well-motivated in relevance feedback settings
  - Train and test communities overlap
- Test error drops fast



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## 2: The typed conductance formulation

- Edge  $e$  has type  $t(e) \in \{1, \dots, T\}$
- Weight  $w(i, j) = \beta(t(i, j))$ , params  $\beta(1), \dots, \beta(T)$
- Matrix  $C$  is now a function of  $\beta$ , denoted  $C(\beta)$
- Find  $\beta$  so that the  $p$  satisfying  $p = C(\beta)p$  also satisfies  $\prec$

Scaling all  $\beta$  preserves  $p$ , so we can demand all  $\beta(t) \geq 1$

min  $\beta' \beta$  subject to:  
 $\beta \geq 1$

$p = C(\beta)p$   
 $p_i \leq p_j$  for all  $i \prec j$

Both  $\beta$  and  $p$  are variables, leading to nasty quadratic equality constraints

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## Two approximations

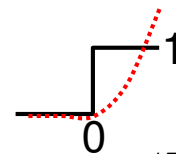
- Breaking the quadratic constraints
  - Approximate  $p \approx C(\beta)^H p^0$  where
    - $p^0$  is the initial Pagerank vector in power iteration
    - $H$  is a finite horizon (or, stop at convergence)
- Design of a loss function

- Training loss (not convex or differentiable)

$$\sum_{u < v} \text{step}(p_u - p_v) = \sum_{u < v} \text{step}((C^H p^0)_u - (C^H p^0)_v)$$

- Approximate using a smooth Huber loss
- Gradient descent search for

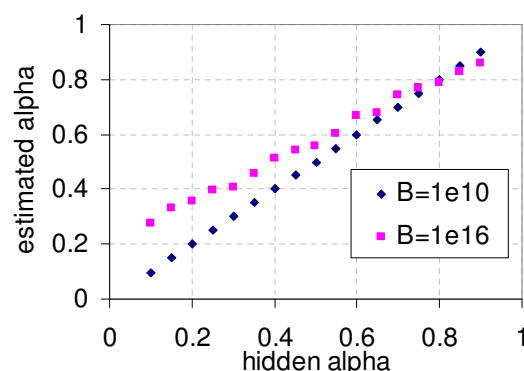
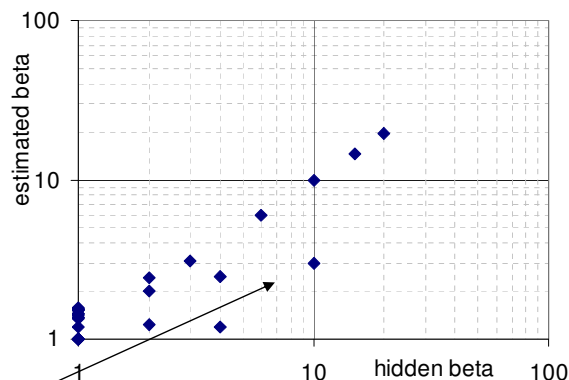
$$\min_{\beta \geq 1} \beta' \beta + B \sum_{u < v} \text{huber}((C^H p^0)_u - (C^H p^0)_v)$$



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## Discovering hidden edge weights

- Assign hidden edge weights to edge types
- Compute weighted Pagerank and sample  $\prec$
- Can recover hidden weights fairly well
  - Penalty on  $\beta' \beta$  shrinks elements toward 0
  - Does not hurt prediction of  $\prec_{\text{test}}$
- Can also find hidden  $\alpha$
- Time scales as  $(|V| + |E|)^{1.34}$

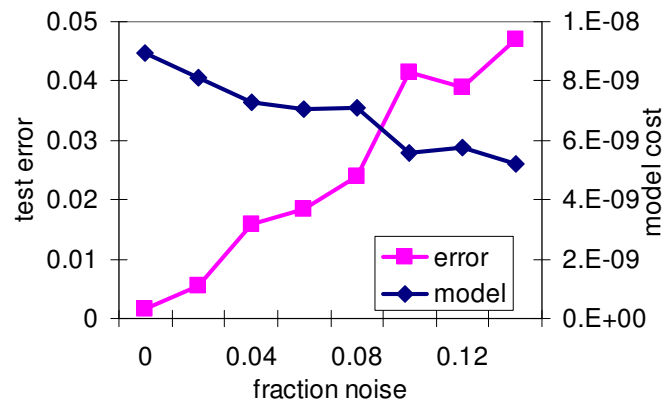
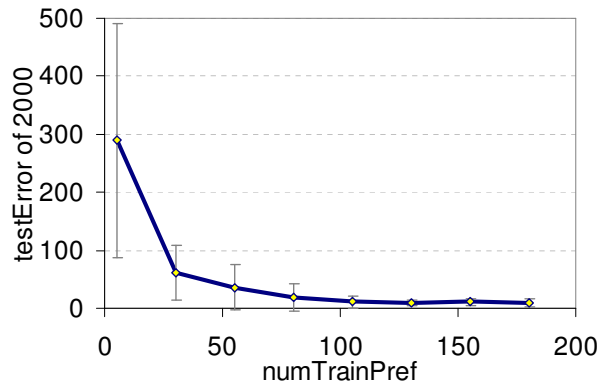


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## Learning rate and robustness

- 20000-node, 120000-edge graph
  - 100 pairwise training preferences enough to cut down test error to 11 out of 2000
  - Training and test preferences node-disjoint
- 20% random reversal of train pairs → 5% increase in test error
  - Model cost  $\beta\beta$  reduces



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## Summary and ongoing work

- Learning to rank nodes in graph from pairwise preferences: surprisingly unexplored
- I.e., design edge conductance so that dominant eigenvector satisfies preferences
- Two design paradigms: constrained flows and typed edge conductance
- New algorithms to learn design parameters
- Integrating queries and node features into models and algorithms (in PKDD 2006)
- Rank-sensitive score learning

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