Indexing, Searching, and Ranking in Entity-Relationship Networks with Associated Text

Soumen Chakrabarti IIT Bombay http://www.cse.iitb.ac.in/~soumen (In fewer words) Ranking and Indexing for Semantic Search

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Working notion of "semantic search"

- Extractors and annotators associate structured knowledge with strings
- Neither complete nor perfect
- No complete schema (despite CYC and WordNet)
- Noisy structure must coexist with source text
- Must exploit structured info and uninterpreted strings in conjunction



Figure: Artist's impression of "semantic search"

Annotated corpus and query examples



Search in entity-relationship graphs



Talk outline

Ranking problems: what is "NEAR"?

- Typed proximity search in text + is-a graphs
- Proximity search in typed graphs
- Discovering hidden favorite communities

Indexing and query processing problems

- Typed proximity search in text + is-a graphs
- Dynamic (query-sensitive) Pageranking on typed graphs

Learning proximity scores of token spans

type=person NEAR "television" "invent*"

- Rarity of selectors
- Distance from candidate position to selectors
- Many occurrences of one selector (closest)
- Combining scores from many selectors (sum)



Making up a feature vector x for position 0

- Limit to $\pm W$ window
- x(-4) = 0; x(-1) = IDF(invent*); x(-6) = IDF(television)
- IDF(w) = numDocs/numDocsWith(w), or perhaps IDF(w) = log(1 + numDocs/numDocsWith(w))
- Other features, e.g., is selector noun? candidate has digits?



If in doubt, throw everything into the kitchen sink

The model vector β

- Score of candidate position is βx
- $\beta(j)$ is the value of the decay function at offset j
- ▶ TREC gives us correct *i* and incorrect *j* token spans
- $i \prec j$ means we want $\beta x_i + \text{margin} \leq \beta x_j$
- Want β to be smooth with $\beta(-W-1) = \beta(W+1) = 0$

$$\begin{split} \min_{\beta} \sum_{j=-W}^{W+1} (\beta_{j-1} - \beta_j)^2 + B \sum_{i \prec j} \mathsf{SmoothLoss}(\beta x_i + 1 - \beta x_j) \\ \min_{\beta} \sum_{j=-W}^{W+1} (\beta_{j-1} - \beta_j)^2 + B \sum_{i \prec j} \log \left(1 + \exp(\beta x_i + 1 - \beta x_j) \right) \end{split}$$

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β fit to TREC QA



- Unexpected decay shape!
- Smooth function slightly better than rough function
- Improves beyond flat scoring function with only IDF

Next: Extending beyond chain graphs

Typed entity-relationship graphs



- ▶ Nodes have entity types: Person, Paper, Email, Company
- Edges have relation types: wrote, sent, cited, in-reply-to
- Edge e has type $t(e) \in \{1, \ldots, T\}$
- ► Edge (u, v) of type t(u, v) has weight β(t(u, v)) and conductance C(v, u)

Conductance, teleport, Pagerank

- Create dummy node d
- Create edges (d, u) and (u, d) for each u
- \blacktriangleright Let 0 $< \alpha < 1$ be the teleport probability
- Let r_v be the probability of teleport to v

$$C_{\{\alpha,\beta\}}(v,u) = \begin{cases} \alpha \frac{\beta(t(u,v))}{\sum_{(u,w)\in E} \beta(t(u,w))}, & u \neq d, v \neq d \\ 1-\alpha, & u \neq d, u \in V_o, v = d \\ 1, & u \neq d, u \in V \setminus V_o, v = d \\ 1, & u = d, v \neq d \\ 0, \text{ otherwise} \end{cases}$$

 ${\it C}$ is a function of α and β

Dodging complicated constraints

$$\min_{\substack{0 < \alpha < 1 \\ \beta \ge 0, p}} \mathsf{ModelCost}(\beta) + B \sum_{i \prec j} \mathsf{SmoothLoss}(p_u - p_v)$$

subject to $p = C_{\{\alpha, \beta\}} p$

- Both C and p are variables \Rightarrow complicated constraints
- Following power iterations, approximate p ≈ C^H p₀ where H is a (possibly adaptive) horizon and p₀ is the initial Pagerank vector (say uniform)
- ▶ Also, $C_{\{\alpha,\beta\}}$ does not change if all β are scaled

 $\min_{\substack{0 < \alpha < 1 \\ \beta \ge 1}} \mathsf{ModelCost}(\beta) + B \sum_{i < j} \mathsf{SmoothLoss}\Big((C^H p_0)_u - (C^H p_0)_v \Big)$

ModelCost and SmoothLoss

Parsimonious model: all $\beta(t)$ s equal

$$\mathsf{ModelCost}(eta) = \sum_{t
eq t'} \Bigl(eta(t) - eta(t')\Bigr)^2$$

If β and $\kappa\beta$ (some multiple $\kappa>1)$ are both solutions, optimizer should prefer β



Can compute approximate $\nabla_\beta,\,\partial/\partial\alpha,$ and use gradient descent

Is SmoothLoss approximation good?

- Hinge and Huber essentially identical
- Empirically, wrt β(t), true and Huber error have same minima
- Wrt α Huber has spurious minima, but basic grid search adequate





Learning rate and robustness

- 20000-node, 120000-edge graph
- 100 pairwise training preferences enough to cut down test error to 11 out of 2000
- 20% random reversal of train pairs leads to 5% increase in test error
- Model cost reduces as noise increases — makes sense



Accuracy of estimating β and α

- Assign hidden β , α
- ► Compute weighted Pagerank and sample ≺
- See if algorithm can recover hidden weights
- ► Upward (downward) pressure on small (large) β thanks to ∑_{t,t'}(β(t) − β(t'))² regularizer
- Large patches of β lead to same Pagerank ordering
- α sensitive to *B* setting



Learning a hidden favored community

- Unlike the random surfer, humans are very selective about following links
- Links in some communities matter, other links do not
- Also preferential teleport
- \prec expresses this indirectly, as in $u_1 \prec v_1$, $u_2 \prec v_2$ etc.
- Goal is to generalize and find the boundaries of favored community
- Unlike global β(t) here info is local; does not extend beyond the "teleport radius"



A constrained flow formulation

- Directly estimate p_{uv} instead of $C(v, u) = \Pr(v|u)$
- *q*_{uv} is a "parsimonious" reference flow (unweighted Pagerank)

$$\begin{split} & \min_{\substack{\{0 \leq p_{uv} \leq 1\} \\ \{0 \leq s_{uv}: u \prec v\}}} & \sum_{(u,v) \in E'} p_{uv} \log \frac{p_{uv}}{q_{uv}} + B \sum_{u \prec v} s_{uv} \quad \text{(SoftObj)} \\ & \text{subject to} & \sum_{(u,v) \in E'} p_{uv} = 1 \quad \text{(Total)} \\ & \forall v \in V' & \sum_{(u,v) \in E'} p_{uv} = \sum_{(v,w) \in E'} p_{vw} \quad \text{(Balance)} \\ & \forall v \in V_o \quad (1 - \alpha) \sum_{(v,w) \in E} p_{vw} = \alpha p_{vd} \quad \text{(Teleport)} \\ & \forall u \prec v \quad (1 + \epsilon) \sum_{(w,u) \in E'} p_{wu} \leq s_{uv} + \sum_{(w,v) \in E'} p_{wv} \quad \text{(SoftPref)} \end{split}$$

Experimental results

- We solve the dual via cutting-plane approach
- Time is linear in | ≺ |, |V| + |E|
- Generalization improves with margin $1 + \epsilon$





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Indexing for is-a proximity search

"Which scientist studied whales?" \rightarrow

type=scientist NEAR study|studied whale*

- Open-domain type hierarchies very large: 15000 internal and 80000 leaf types in WordNet (full set A)
- Runtime type expansion too expensive: even WordNet knows 650 scientist, 860 cities, ...

Pre-generalize

- Index a subset $R \subset A$
- ► Query atype a ∉ R, want k answers
- Probe index with g, ask for k' > k



Cost models

 How much space saved by indexing *R* instead of *A*? (Cannot afford to try out many *R*s, need quick estimate)
 What is the average query time bloat owing to *a* → *g*

Post-filter

- Fetch k' high-scoring spans w for g
- Check if w is-a a as well (using forward and reachability index); if not, discard
- If fewer than k survive, restart with larger k' (expensive!)



Cost models

- How much space saved by indexing R instead of A? (Cannot afford to try out many Rs, need quick estimate)
- What is the average query time bloat owing to a → g pre-generalize and post-filter?

Post-filter

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Index space estimate

- Let corpusCount(a) be the count of tokens w in the corpus such that w is-a a
- One posting entry for each count of each type a
- Therefore our space estimate is (proportional to) ∑_{a∈R} corpusCount(a)
- Surprisingly accurate despite index compression



Characterizing a query workload

$$\widetilde{\mathsf{Pr}}(a) = rac{\mathsf{queryLogCount}(a) + \lambda}{\sum_{a' \in \mathcal{A}} \left(\mathsf{queryLogCount}(a) + \lambda
ight)}$$

- Heavy-tailed type distribution in queries
- Many test types never seen in training types and vice versa



- $\lambda = 0$ would give these types zero probability
- Danger of allocating no g close to these as
- Build multinomial model over a with positive λ
- Cross-validate likelihood of held-out test log

Query time bloat estimate

- t_{scan} time to scan one candidate position while merging postings
- t_{filter} time to check if w is-a a
- If R = A (all types indexed), query takes time roughly t_{scan} corpusCount(a)
- ▶ If $a \notin R$, the price paid for generalization to g consists of
 - Longer scans: t_{scan} corpusCount(g)
 - Post-filtering k' responses: $k' t_{\text{filter}}$

$$\mathsf{expected \ bloat} = \sum_{a \in A} \widetilde{\mathsf{Pr}}(a) \frac{t_{\mathsf{scan}} \operatorname{corpusCount}(g) + k' t_{\mathsf{filter}}}{t_{\mathsf{scan}} \operatorname{corpusCount}(a)}$$

Now we have a greedy cost-benefit analysis of every type a

Result of greedy knapsack

Estimated query bloat	Corpus/Index	GBytes
reasonably accurate	Original corpus	5.72
	Gzipped corpus	1.33
With only 520 MB index	Stem index	0.91
	Full type A index	4.30
only 1.9 average bloat	Type subset <i>R</i> index	0.52
	Query Bloat	1.90
Space comparable to	Reachability index	0.01
inverted index on stems	Forward index	1.16



Searching a TypedWordGraph



- Attach query words W to preloaded entities N
- Set teleport r > 0 only for word nodes
- Compute personalized Pagerank vector (PPV) p_r slow!

Options to date

Query-time Pagerank

- \blacktriangleright For ~75000 nodes, ~200000 edges, ~11 sec/query
- Impractical except for very small graphs

Combine per-word PPVs

- \blacktriangleright Proposed by $\operatorname{OBJECTRANK}$
- ▶ For ~ 75000 nodes, ~ 175000 words, ~ 526 CPU-hours
- ▶ Full PPV index has size 102 GB
- Compare with text index: 56 MB
- Truncating down to 56 MB leads to serious loss of accuracy



KTau=Kendall's Tau

$\operatorname{HUBRANK}$ query execution

- Input: query words W, abandon/trim threshold δ
- Max priority queue, priority of node u is estimate of path conductance from d to u
- Grow active set A



A is bordered with

Blockers: Hub nodes with indexed (approx) PPVs Losers ℓ : Conductance from d to ℓ is too small to matter

- ► Load trimmed PPV for blockers, trivial PPV for losers
- Iteratively compute PPVs of all active nodes including d
- Sort PPV of d and return results

$\operatorname{HubRank}\ \operatorname{\textbf{results}}$



Query in 250 ms (vs. 11 s)



Conclusion

- Searching typed entity-relationship graphs
- Perhaps attached to mentions in unstructured text
- Many interesting challenges, surprisingly unexplored
- Architecture for flexible space-time-accuracy tradeoffs
- Ranking
 - ► Far from vector-space territory, no guidance
 - Learning linear proximity functions
 - Learning edge conductance parameters
- Indexing and query processing
 - Combining large is-a graphs with linear proximity
 - Dynamic personalized Pageranking
 - Anytime preprocessing, anytime query processing