

CS 716: Introduction to communication networks

- 20th class; 14th Oct 2011

Instructor: Sridhar Iyer
IIT Bombay

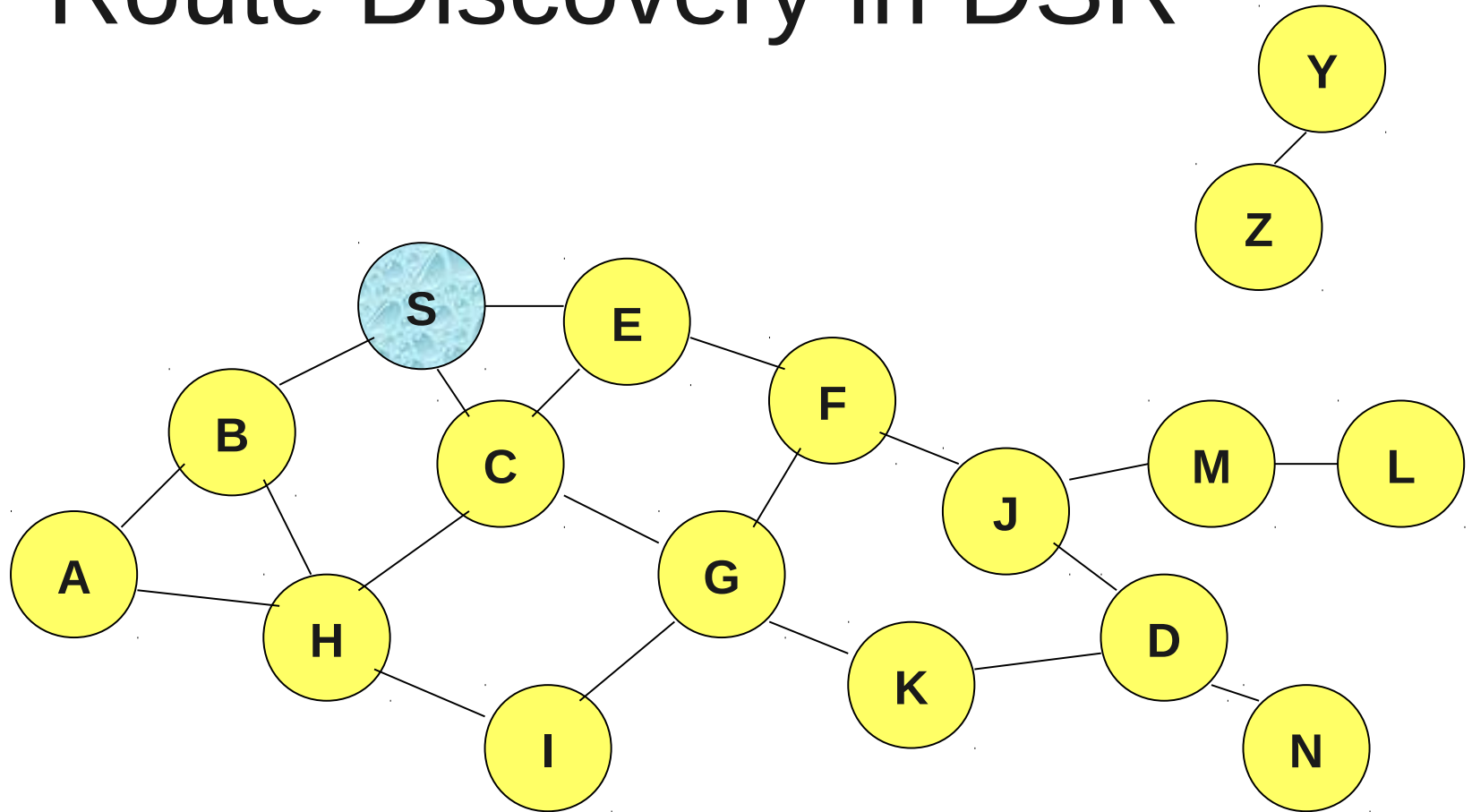
Routing methods

- Static routing: Default routes are specified at boot time
- Dynamic methods:
 - Source-based: Specify route at source (DSR)
 - Distance-vector routing: Set up next-hops to destinations looking at neighbors' routing tables (RIP)
 - Link-state routing: Get map of network in terms of link states and calculate best route (but specify only the next-hop) (OSPF)

Dynamic Source Routing (DSR)

- Source S initiates a **route discovery** by flooding **Route Request (RREQ)**
 - Each node **appends its own identifier** when forwarding **RREQ**
- Destination D on receiving the first **RREQ**, sends a **Route Reply (RREP)**
 - **RREP** sent on route obtained by **reversing** the route appended in **RREQ**
 - **RREP includes the route** from S to D, on which **RREQ** was received by D
- S routes data using “source route” mechanism

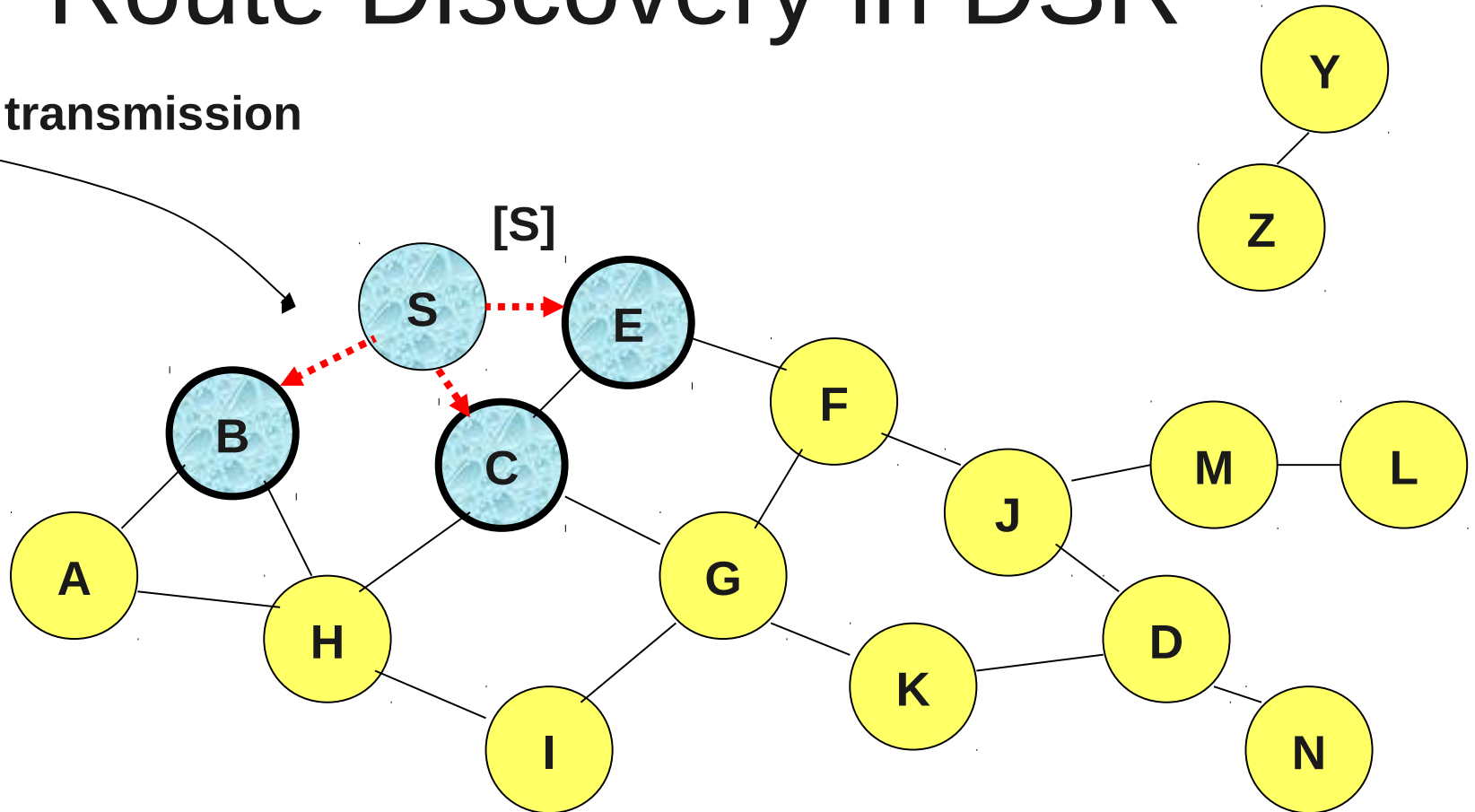
Route Discovery in DSR



Represents a node that has received RREQ for D from S

Route Discovery in DSR

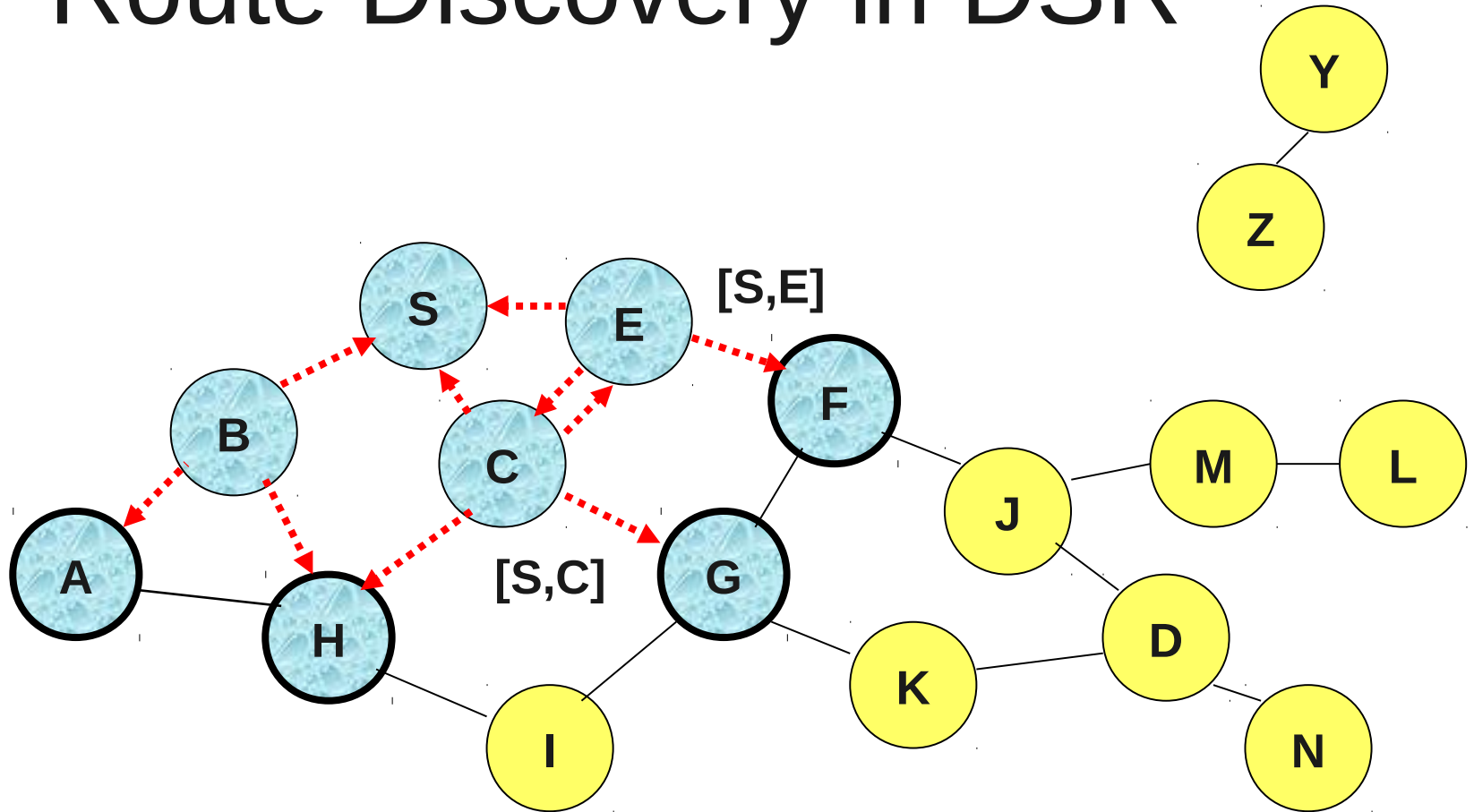
Broadcast transmission



.....→ Represents transmission of RREQ

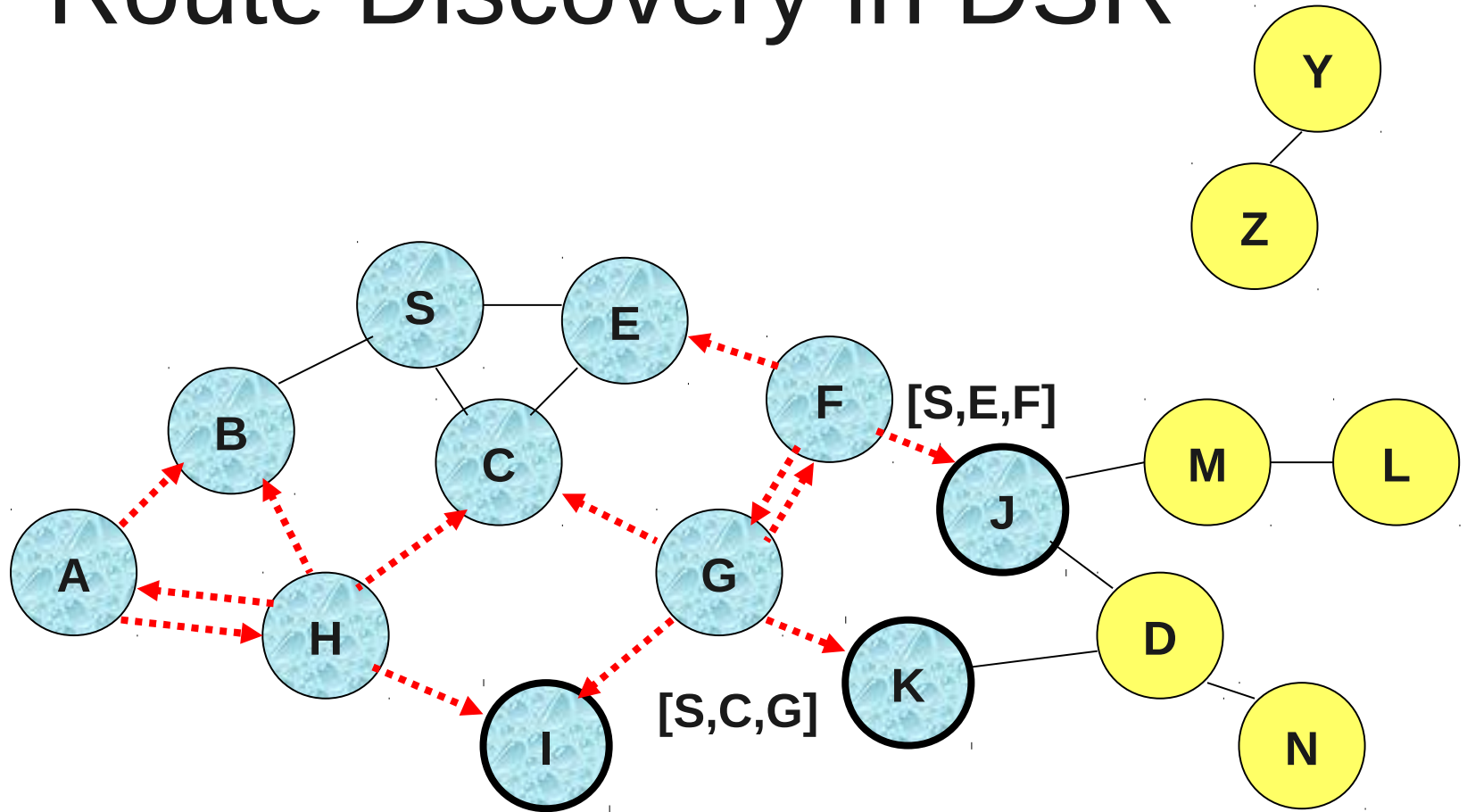
[X,Y] Represents list of identifiers appended to RREQ

Route Discovery in DSR



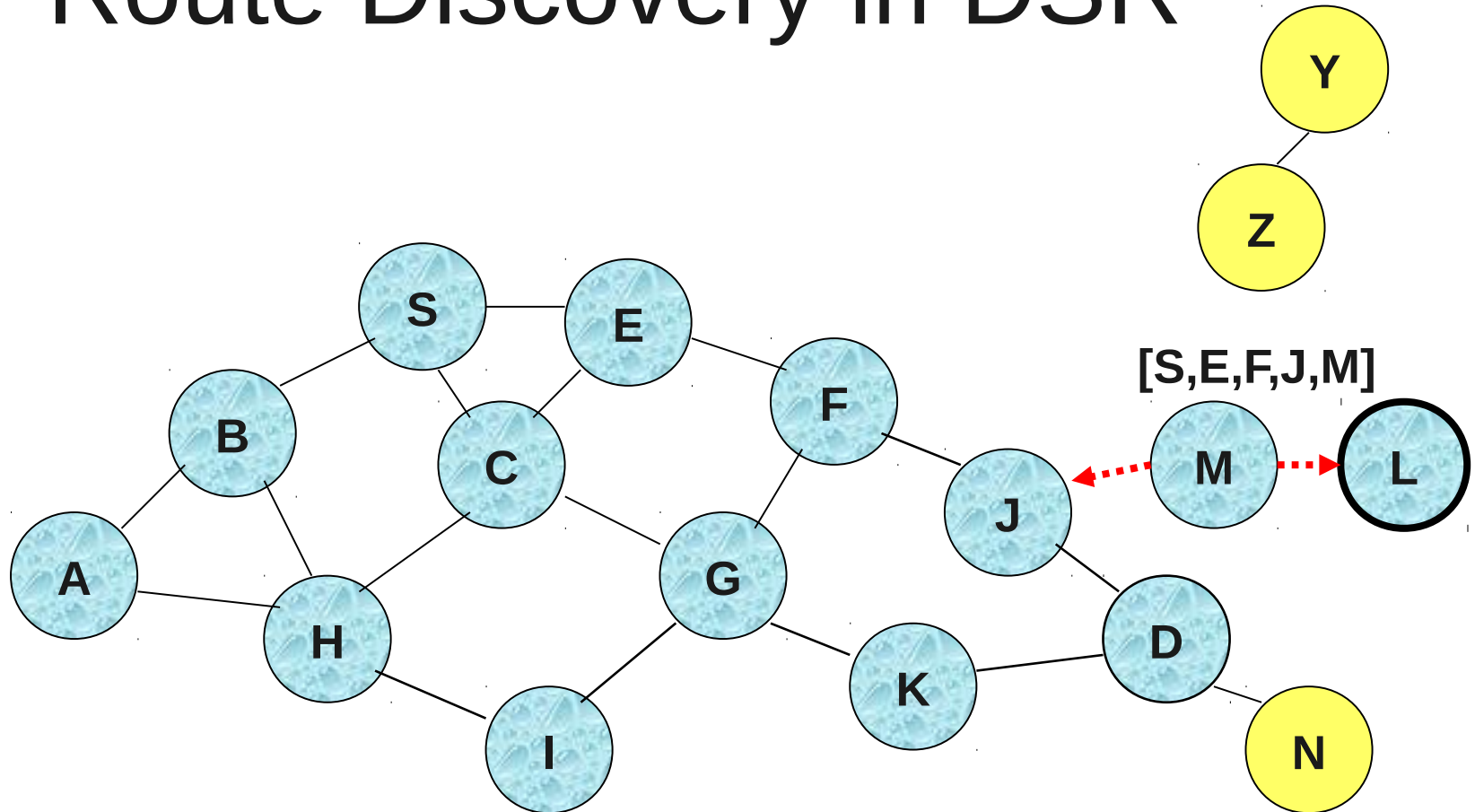
- Node H receives packet RREQ from two neighbors:
potential for collision

Route Discovery in DSR



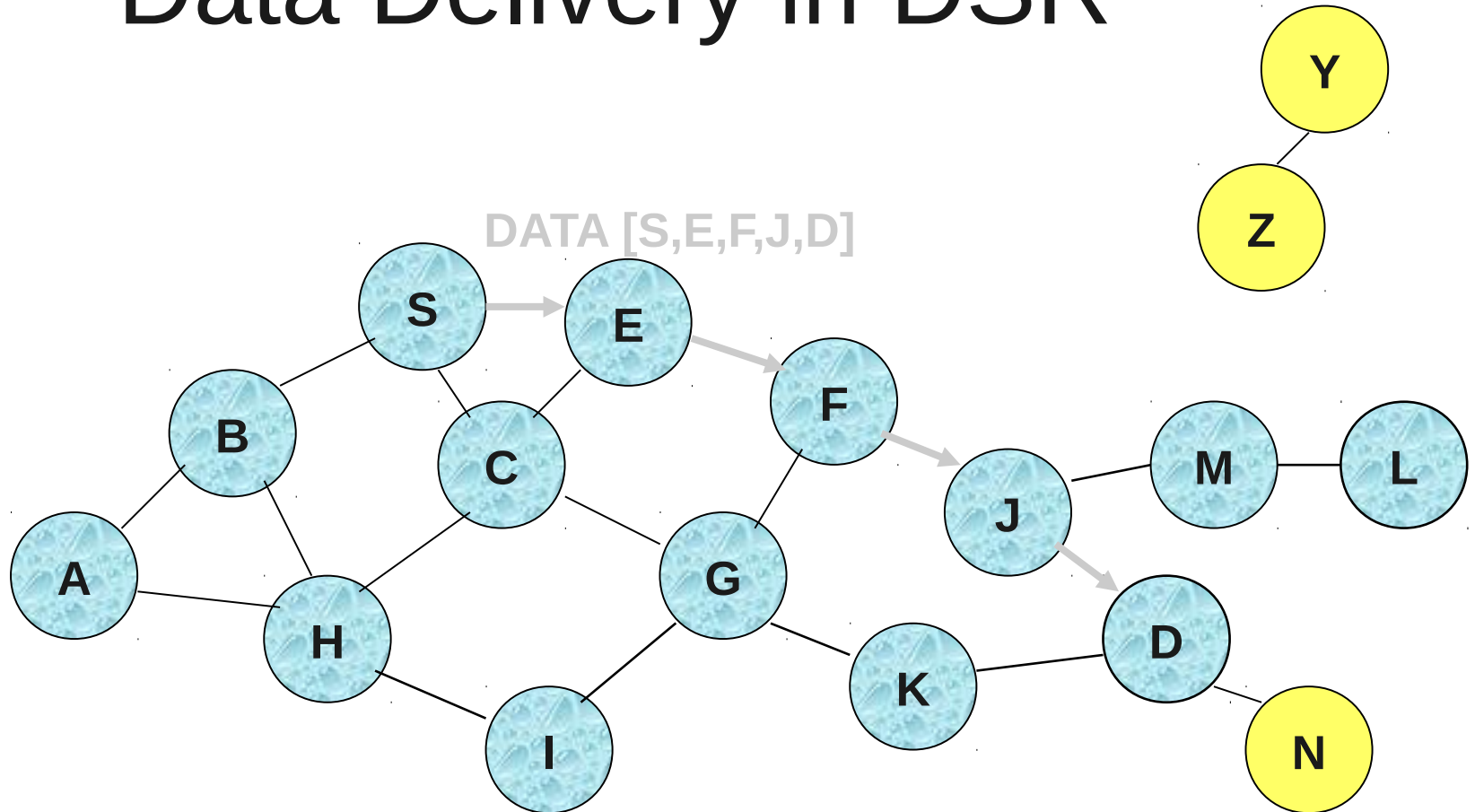
- Node C receives RREQ from G and H, but does not forward it again, because node C has **already forwarded RREQ** once

Route Discovery in DSR



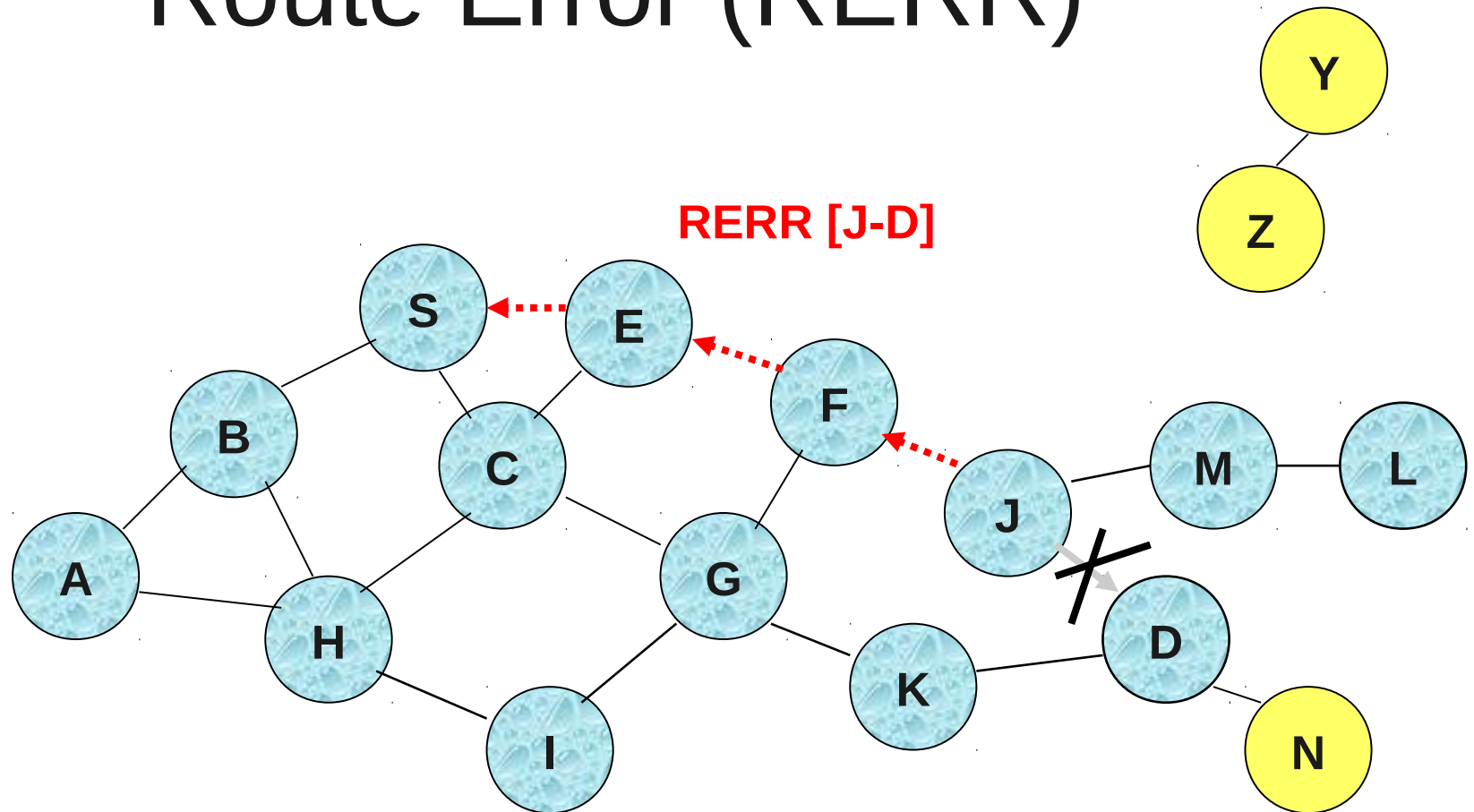
- Node D **does not forward** RREQ, because node D is the **intended target** of the route discovery

Data Delivery in DSR



Packet header size grows with route length

Route Error (RERR)



J sends a route error to S along route J-F-E-S when its attempt to forward the data packet S (with route SEFJD) on J-D fails

DSR: Route caching

- Each node caches a new route it learns by *any means*
- When node S finds route [S,E,F,J,D] to node D, node S also learns route [S,E,F] to node F
- When node K receives **Route Request [S,C,G]** destined for node, node K learns route [K,G,C,S] to node S

Route caching

- When node F forwards **Route Reply RREP** **[S,E,F,J,D]**, node F learns route [F,J,D] to node D
- When node E forwards **Data** [S,E,F,J,D] it learns route [E,F,J,D] to node D
- A node may also overhear Data to learn routes

Route caching

- Uses:
 - Finding alternate routes in case original route breaks
 - Route reply from intermediate nodes
- Problems:
 - Cached routes may become invalid over time and due to host mobility
 - Stale caches can adversely affect performance

DSR: Advantages

- Routes maintained only between nodes who need to communicate
 - reduces overhead of route maintenance
- Route caching can further reduce route discovery overhead
 - A single route discovery may yield many routes to the destination, due to intermediate nodes replying from local caches

DSR: Disadvantages

- Packet header size grows with route length due to source routing
- Flood of route requests may potentially reach all nodes in the network
- An intermediate node may send Route Reply using a stale cached route, thus polluting other caches

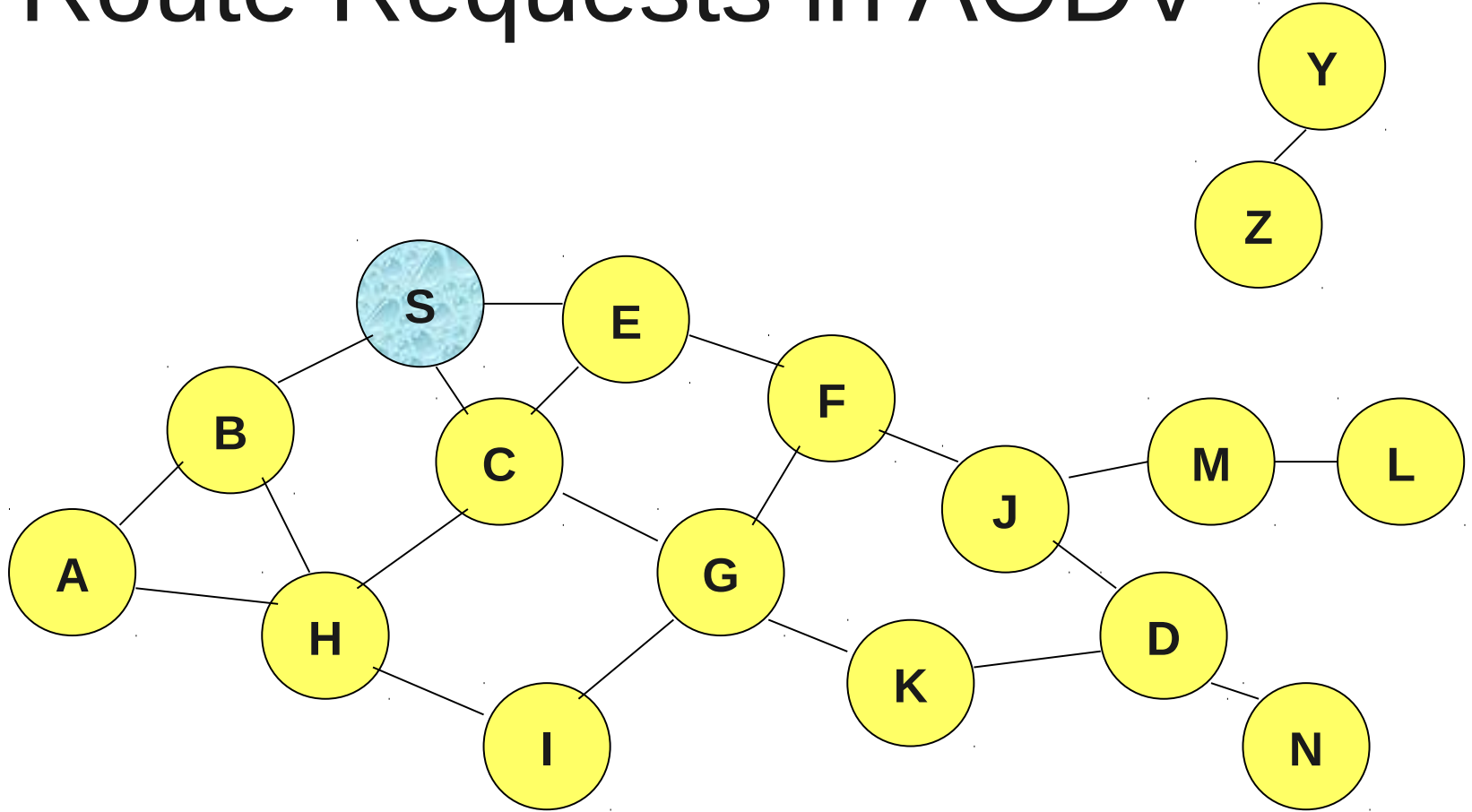
Ad Hoc On-Demand Distance Vector Routing (AODV)

- DSR includes source routes in packet headers
- Resulting large headers can sometimes degrade performance
 - particularly when data contents of a packet are small
- AODV attempts to improve on DSR by maintaining routing tables at the nodes, so that data packets do not have to contain routes

AODV

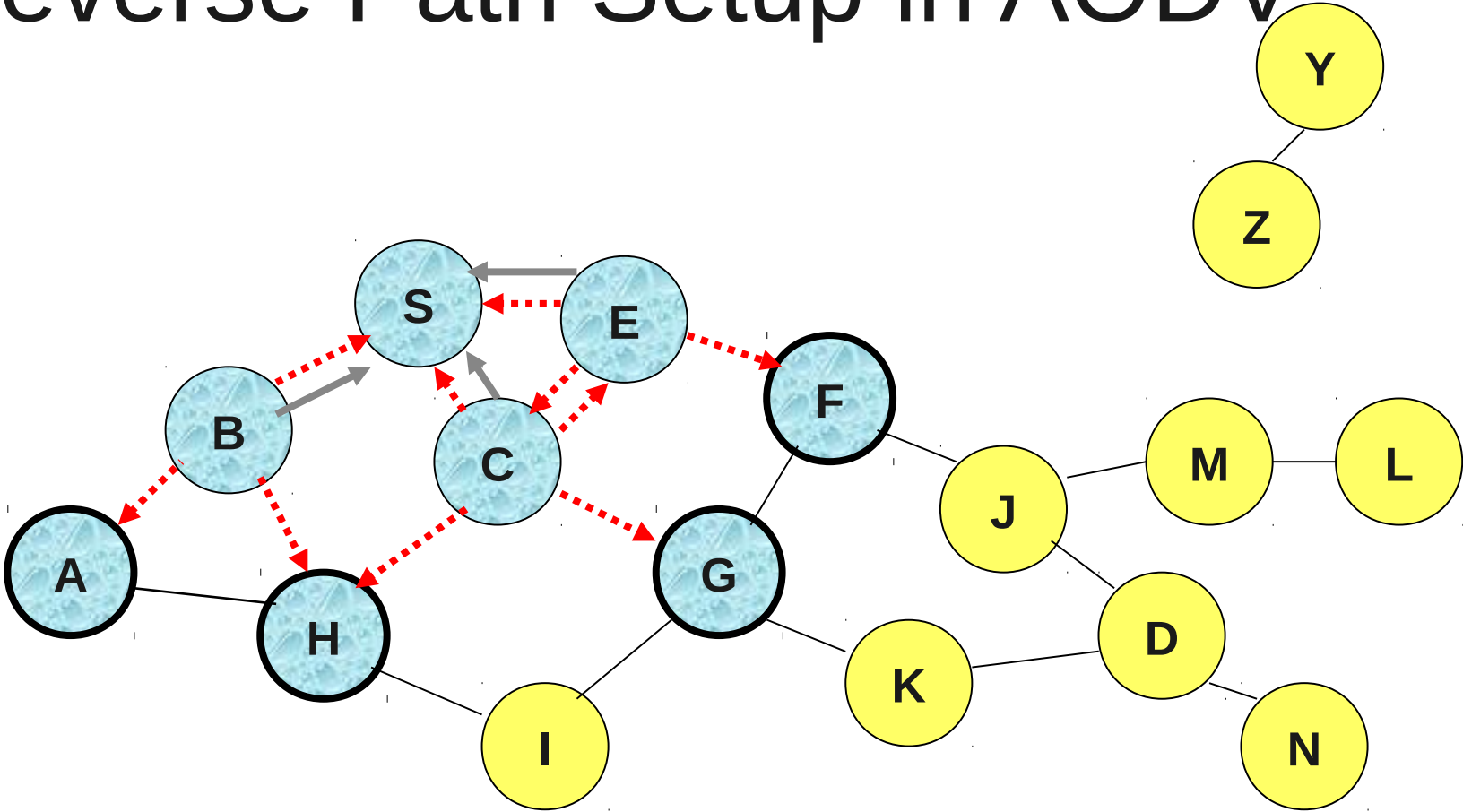
- **Route Requests (RREQ)** are forwarded in a manner similar to DSR
- When a node re-broadcasts a Route Request, it sets up a reverse path pointing towards the source
- **Route Reply (RREP)** travels along the reverse path set-up when Route Request is forwarded

Route Requests in AODV



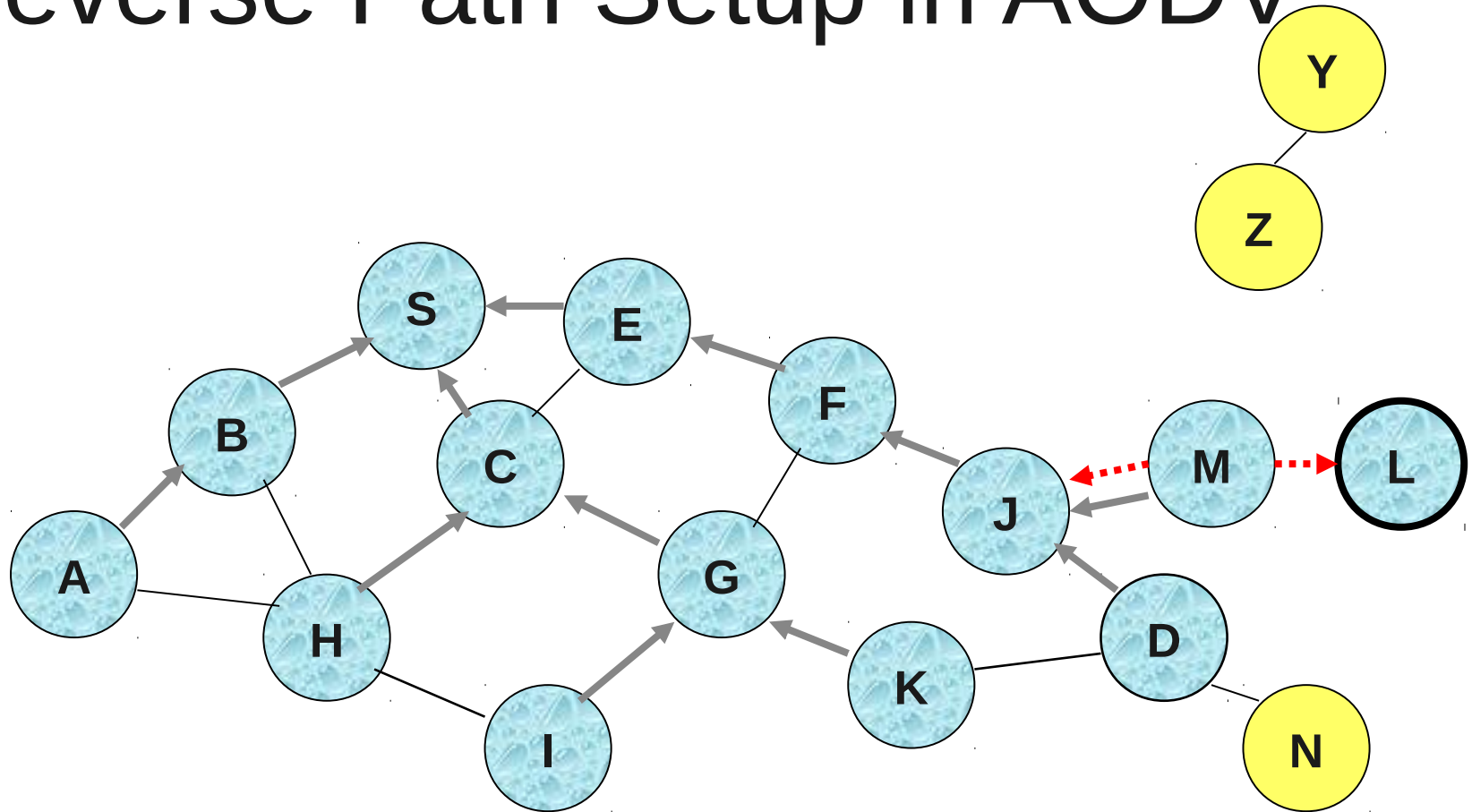
Represents a node that has received RREQ for D from S

Reverse Path Setup in AODV



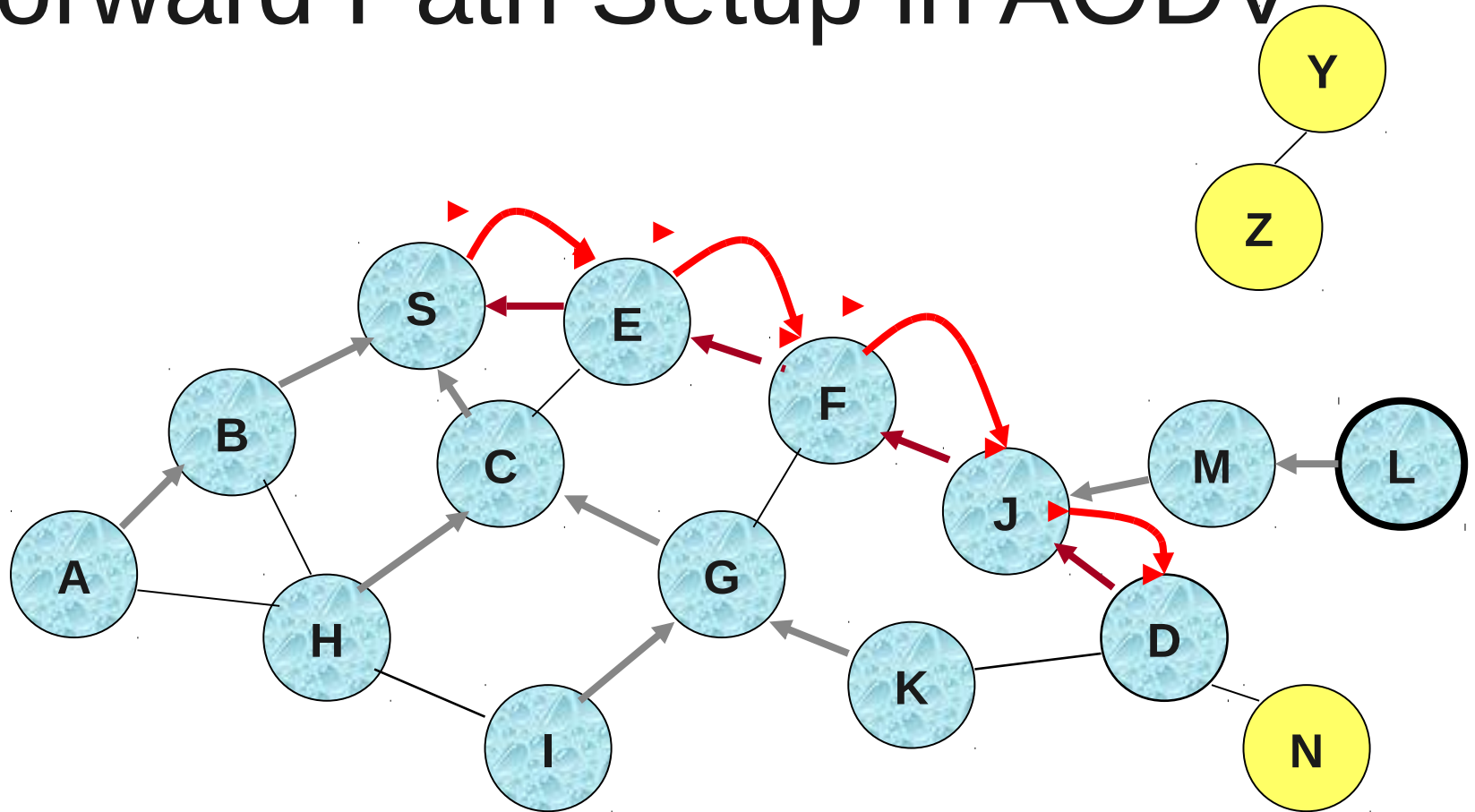
← Represents links on Reverse Path

Reverse Path Setup in AODV



- Node D **does not forward** RREQ, because node D is the **intended target** of the RREQ

Forward Path Setup in AODV



Forward links are setup when RREP travels along the reverse path



Represents a link on the forward path

Route Request and Route Reply

- Route Request (RREQ) includes the last known **sequence number** for the destination
- An intermediate node may also send a Route Reply (RREP) provided that it knows a **more recent path** than the one previously known to sender
- Intermediate nodes that forward the RREP, also record the next hop to destination

AODV: Timeouts

- Neighboring nodes periodically exchange **hello** message
- A routing table entry maintaining a **reverse path** is purged after a timeout interval
- A routing table entry maintaining a **forward path** is purged if *not used* for a ***active_route_timeout*** interval

AODV: Link failure

- Absence of hello message is used as an indication of link failure
- When the next hop link in a routing table entry breaks, all **active** neighbors are informed
- Link failures are propagated by means of **Route Error (RERR)** messages, which also update destination sequence numbers

AODV: Expanding ring search

- Route Requests are initially sent with small Time-to-Live (TTL) field, to limit their propagation
 - DSR also includes a similar optimization
- If no Route Reply is received, then larger TTL tried

AODV: Summary

- Routes need not be included in packet headers
- Nodes maintain routing tables containing entries only for routes that are in active use
- At most one next-hop per destination maintained at each node
- Sequence numbers are used to avoid old/broken routes and prevent routing loops