

Tutorial 5

Instructions:

- The following problems are meant for you to practice, so that your understanding of the topic improves.
- You must solve all problems to get the maximum benefit from practice problems.
- You must not submit your solutions to these problems. These are not going to be graded.
- A problem may have multiple solution techniques. Discussion among students is strongly encouraged in order to understand different perspectives.

1. For each of the following sentences, either give a natural deduction proof of its validity, or show that there is a structure over which the sentence evaluates to false.

1. $((\forall x P(x)) \vee (\forall x Q(x))) \rightarrow \forall x (P(x) \vee Q(x))$
2. $\forall x ((\exists y R(x, y)) \vee \neg R(x, x))$
3. $\forall x \forall y (S(x, y) \rightarrow (\exists z (S(x, z) \wedge S(z, y))))$
4. $\exists y ((\forall x P(x)) \rightarrow P(y))$
5. $\forall x \forall y (S(x, y) \rightarrow (x = y)) \rightarrow \exists x S(x, x)$

2. Consider the vocabulary $\mathcal{V} = \{E, =\}$. We have seen earlier that every \mathcal{V} -structure can be viewed as a directed graph. A \mathcal{V} -sentence φ therefore defines a family of directed graphs, namely all those graphs, which when viewed as \mathcal{V} -structures satisfy φ .

For each of the following families of graphs, either give a first-order logic sentence that defines the family precisely, or use the Compactness Theorem to show that there cannot exist any first-order logic sentence that defines the family.

1. $F_1 = \{G \mid G \text{ is a graph with no even length cycle}\}$.
2. $F_2 = \{G \mid G \text{ is a graph in which every vertex has an even number of edges going out of it}\}$.
3. $F_3 = \{G \mid G \text{ is a graph in which the longest path length is even}\}$.
4. $F_4 = \{G \mid \text{every cycle in } G \text{ is of length at most } 2\}$.
5. $F_5 = \{G \mid \text{There is at most one vertex that lies on a cycle of length } 2\}$.