# CS781: <br> A Quick Primer on Abstract Interpretation for Neural Networks 

## Supratik Chakraborty IIT Bombay

## Notion of State in Neural Network



State: $\left(X_{1}, X_{2}, \ldots x_{18}\right)$ in $R^{18}$

## State Change in Feed-Forward Neural Network


$\left(x_{1}^{\prime}, x_{2}^{\prime}, \ldots x_{i-1}^{\prime}, x_{i}^{\prime}\right)=f_{i}\left(x_{1}, x_{2}, \ldots x_{i-1}\right)$, for i in $\{3, \ldots, 18\}$

## State Change in Feed-Forward NN as a sequence of instrns



$$
\begin{array}{ll}
\left(x_{1}^{\prime}, x_{2}^{\prime}, x_{3}^{\prime}\right) & =f_{3}\left(x_{1}, x_{2}\right) ; \\
\left(x_{1}^{\prime \prime}{ }_{1}, x^{\prime \prime}{ }_{2}, x^{\prime \prime \prime}{ }_{3,} x^{\prime \prime \prime}{ }_{4}\right) & =f_{4}\left(x_{1}^{\prime}, x_{2}^{\prime}, x_{3}^{\prime}\right) ;
\end{array}
$$

NN computation: a sequence of state transitions caused by seq of instructions

## Proving Property of a FF NN

Pre-
condition on (x1, x2)


Post-
condition on (x17, x18)
\{Pre-condition on (x1, x2)\}

$$
\begin{array}{ll}
\left(x^{\prime}, x_{2}^{\prime}, x_{3}^{\prime}\right) & =f_{3}\left(x_{1}, x_{2}\right) ; \\
\left(x^{\prime \prime}{ }_{1}, x^{\prime \prime}{ }_{2}, x^{\prime \prime}{ }_{3,} x^{\prime \prime}{ }_{4}\right) & =f_{4}\left(x_{1}^{\prime}, x_{2}^{\prime}, x_{3}^{\prime}\right) ;
\end{array}
$$

\{Post-condition on $(x 17, x 18)\}$

## NN Computation as a State Transition System



$$
\begin{aligned}
& \text { \{Pre-condition on }(x 1, x 2)\} \\
& \begin{array}{ll}
\left(x_{1}^{\prime}, x_{2}^{\prime}, x_{3}^{\prime}\right) & =f_{3}\left(x_{1}, x_{2}\right) ; \\
\left(x^{\prime \prime \prime}, x_{1},{ }_{2}^{\prime \prime}, x^{\prime \prime}{ }_{3} x^{\prime \prime \prime}{ }_{4}\right) & =f_{4}\left(x_{1}^{\prime}, x_{2}^{\prime}, x_{3}^{\prime}\right) ; \\
\ldots & \\
\text { \{Post-condition on }(x 17, x 18)\}
\end{array}
\end{aligned}
$$

## Dealing with State Space Size

- Infinite state space
- Difficult to represent using state transition diagram
- Can we still do some reasoning?
» Solution: Use of abstraction
Concrete states
- Naive view
- Bunch sets of states together"ntelligently"
- Don't talk of individual states, talk of a representation of a set of states
- Transitions between state set representations
- Granularity of reasoning shifted
- Extremely powerful general technique
- Allows reasoning about large/infinite state spaces


## A Generic View of Abstraction

Set of concrete states Set of abstract states


Abstraction ( $\alpha$ )

ン Every subset of concrete states mapped to unique abstract state

- Desirable to capture containment relations
> Transitions between state sets (abstract states)


## The Game Plan



Abstract analysis engine

## The Game Plan



How do we choose the right abstraction? Is there a method beyond domain expertise? Can we learn from errors in abstraction to build better (refined) abstractions? Can refinement be automated?

## The Game Plan

Abstract state spaces can be infinite. What can we do to make abstract analysis practical? Finite ascending chains what beyond?

## A <br> T <br> E



Abstract analysis engine

## Desirable Properties of Abstraction

Set of concrete states

## Set of abstract states



Abstraction ( $\alpha$ )


Concretization ( $\gamma$ )

> Suppose $S_{1} \subseteq S_{2}$ : subsets of concrete states

- Any behaviour starting from $S_{1}$ can also happen starting from $S_{2}$
- If $\alpha\left(S_{1}\right)=a_{1}, \alpha\left(S_{2}\right)=a_{2}$ we want this monotonicity in behaviour in abstr state space too
- Need ordering of abstract states, similar in spirit to $S_{1} \subseteq S_{2}$


## Structure of Concrete State Space

- Set of concrete states: S



## Structure of Abstract State Space

${ }^{\text {' }}$ Abstract lattice $\mathrm{A}=(\mathcal{A}, \sqsubseteq, \sqcup, \sqcap, \top, \perp)$
, Abstraction function $\alpha: \wp(S) \rightarrow \mathcal{A}$

- Monotone: $S_{1} \subseteq S_{2} \Rightarrow \alpha\left(S_{1}\right) \sqsubseteq \alpha\left(S_{2}\right)$ for all $S_{1}, S_{2} \subseteq S$
- $\alpha(S)=$ Т, $\quad \alpha(\emptyset)=\perp$
- Concretization function $\gamma: \mathcal{A} \rightarrow \wp(S)$
- Monotone: $a_{1} \sqsubseteq a_{2} \Rightarrow \gamma\left(a_{1}\right) \subseteq \gamma\left(a_{2}\right)$ for all $a_{1}, a_{2} \in \mathcal{A}$
- $\gamma(\top)=S, \quad \gamma(\perp)=\emptyset$

