

Helper Slides on Abduction-based Minimal Explanations

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Abduction in Logic

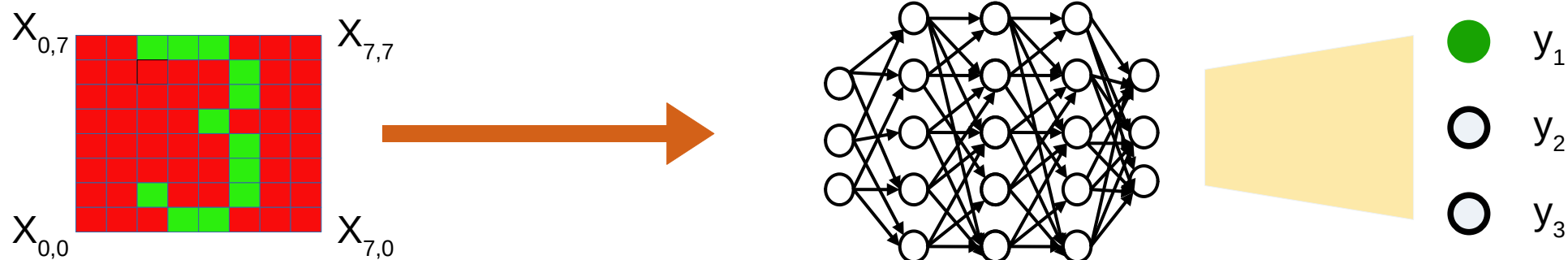
Given a theory (consistent set of sentences) \mathcal{F} and a formula \mathcal{E} in a logic \mathcal{L}
Find a formula α such that

- $\alpha \models \mathcal{F} \Rightarrow \mathcal{E}$
- $\mathcal{F} \wedge \alpha$ is consistent

We often want α to be as weak (permissive) as possible.

α is an "explanation" of \mathcal{E} in theory \mathcal{F}

Formulating Explanation as Abduction



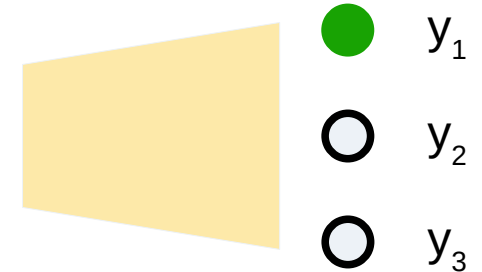
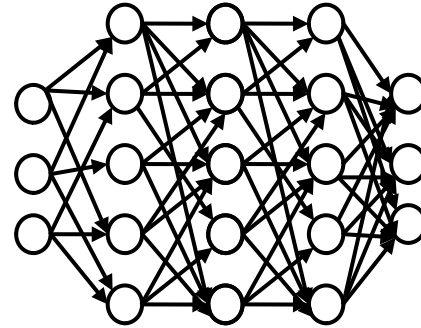
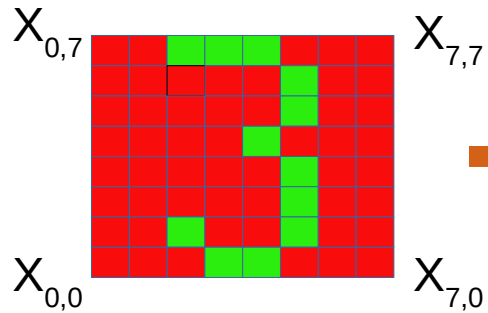
$$\begin{aligned}
 \mathcal{C} = & \\
 & (x_{0,7} = R) \wedge (x_{1,7} = R) \wedge (x_{2,7} = G) \wedge (x_{3,7} = G) \wedge \cdots \wedge (x_{7,7} = R) \wedge \\
 & \quad \vdots \\
 & (x_{0,0} = R) \wedge (x_{1,0} = R) \wedge (x_{2,0} = R) \wedge (x_{3,0} = G) \wedge \cdots \wedge (x_{7,0} = R)
 \end{aligned}$$

 \mathcal{F}

$$\mathcal{E} = (y_1 > y_2) \wedge (y_1 > y_3)$$

Clearly, $\mathcal{C} \wedge \mathcal{F} \wedge \mathcal{E}$ is consistent.

Formulating Explanation as Abduction



$\mathcal{C} =$

$$(x_{0,7} = R) \wedge (x_{1,7} = R) \wedge (x_{2,7} = G) \wedge (x_{3,7} = G) \wedge \cdots \wedge (x_{7,7} = R) \wedge$$

\vdots

$$(x_{0,0} = R) \wedge (x_{1,0} = R) \wedge (x_{2,0} = R) \wedge (x_{3,0} = G) \wedge \cdots \wedge (x_{7,0} = R)$$

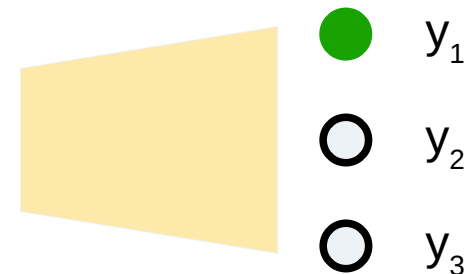
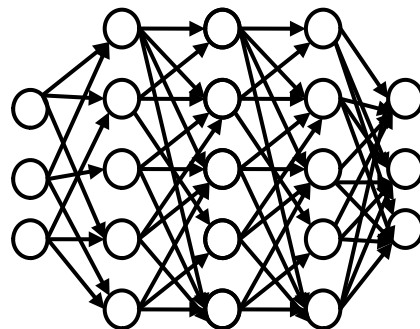
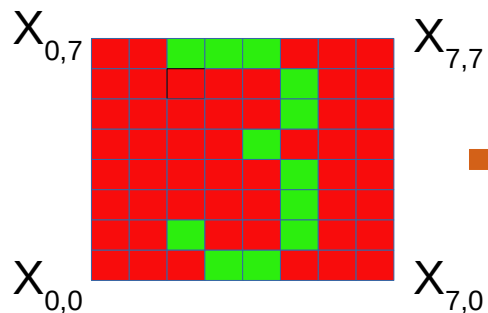
\mathcal{F}

$$\mathcal{E} = (y_1 > y_2) \wedge (y_1 > y_3)$$

Find smallest $\mathcal{C}' \subseteq \mathcal{C}$ s.t.

(a) $\mathcal{C}' \wedge \mathcal{F}$ is consistent, and (b) $\mathcal{C}' \models \mathcal{F} \Rightarrow \mathcal{E}$

Building C' Lazily



$C =$

$$(x_{0,7} = R) \wedge (x_{1,7} = R) \wedge (x_{2,7} = G) \wedge (x_{3,7} = G) \wedge \dots \wedge (x_{7,7} = R) \wedge$$

\vdots

$$(x_{0,0} = R) \wedge (x_{1,0} = R) \wedge (x_{2,0} = R) \wedge (x_{3,0} = G) \wedge \dots \wedge (x_{7,0} = R)$$

\mathcal{F}

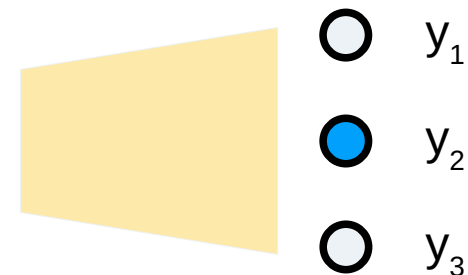
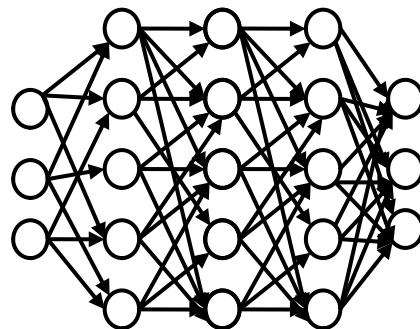
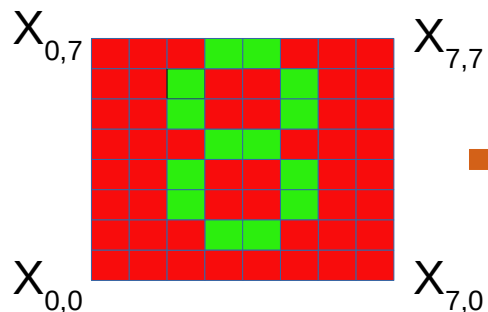
$$\mathcal{E} = (y_1 > y_2) \wedge (y_1 > y_3)$$

Does the empty subset of C suffice?

Does

$\models \mathcal{F} \Rightarrow \mathcal{E}$ hold?

Building C' Lazily



$\hat{C} =$

$$(x_{0,7} = R) \wedge (x_{1,7} = R) \wedge (x_{2,7} = R) \wedge (x_{3,7} = G) \wedge \cdots \wedge (x_{7,7} = R) \wedge$$

\vdots

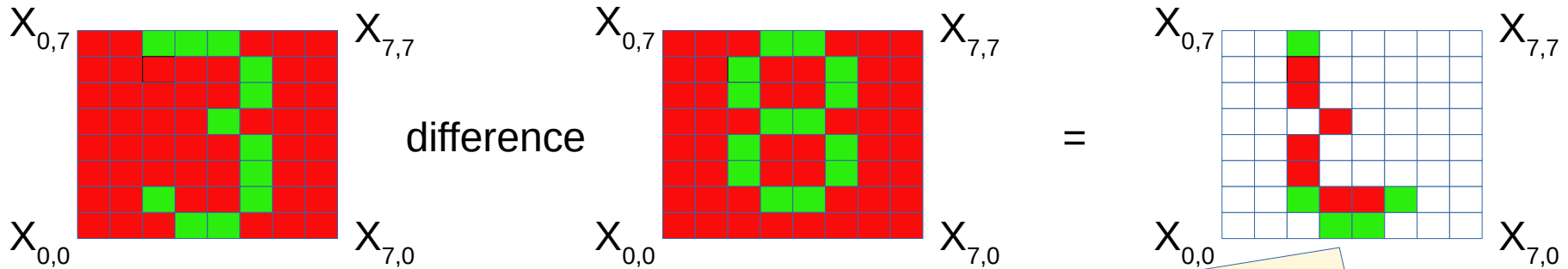
$$(x_{0,0} = R) \wedge (x_{1,0} = R) \wedge (x_{2,0} = R) \wedge (x_{3,0} = R) \wedge \cdots \wedge (x_{7,0} = R)$$

\mathcal{F}

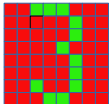
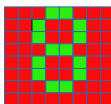
$$\hat{\mathcal{E}} = (y_2 > y_1) \wedge (y_2 > y_3)$$

Certainly $\models \mathcal{F} \Rightarrow \mathcal{E}$ doesn't hold!

How do the two inputs differ?



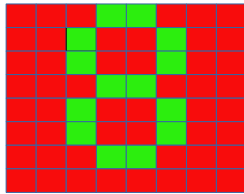
$$S_1 = \{ (x_{2,7} = G), (x_{2,6} = R), (x_{2,5} = R), (x_{3,4} = R), (x_{2,3} = R), (x_{2,2} = R), (x_{2,1} = G), (x_{3,1} = R), (x_{4,1} = R), (x_{5,1} = G), (x_{3,0} = G), (x_{4,0} = G) \}$$

Unless one of the literals in S_1 is included in the explanation C' , we can't distinguish between  and 

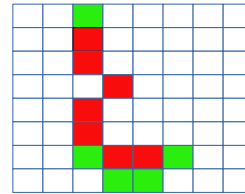
Choosing subset of C



difference


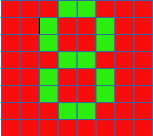


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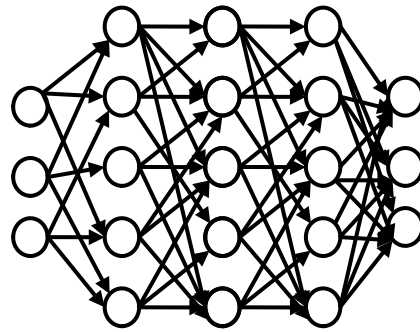
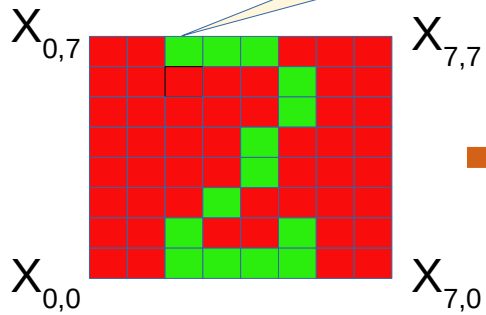
$$S_1 = \{(x_{2,7} = G), (x_{2,6} = R), (x_{2,5} = R), (x_{3,4} = R), (x_{2,3} = R), (x_{2,2} = R), (x_{2,1} = G), (x_{3,1} = R), (x_{4,1} = R), (x_{5,1} = G), (x_{3,0} = G), (x_{4,0} = G)\}$$




Suppose we choose $(x_{2,7} = G)$ for $C' \subseteq C$

Certainly this distinguishes  from 

So, have we found the explanation?

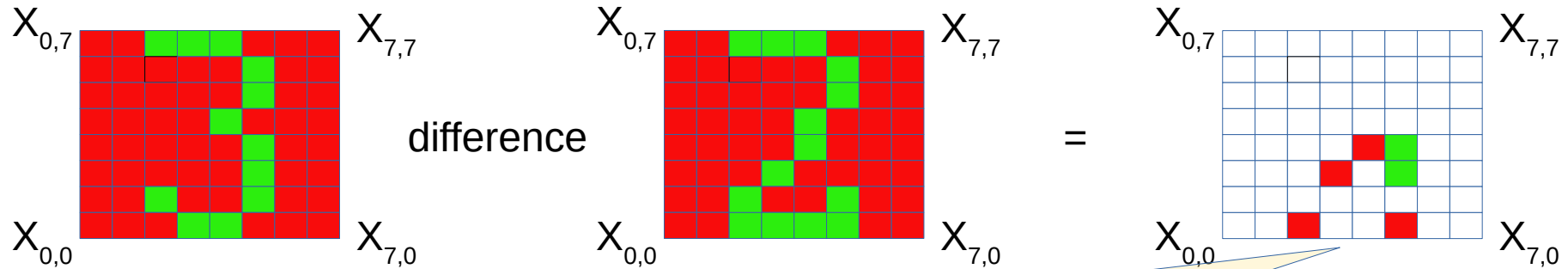
Does $(x_{2,7} = G) \models \mathcal{F} \Rightarrow \mathcal{E}$ hold?



-  y_1
-  y_2
-  y_3


Clearly not!

How do the two inputs differ again?



$$S_2 = \{ (x_{4,3} = R), (x_{5,3} = G), (x_{3,2} = R), (x_{5,2} = G), (x_{2,0} = R), (x_{5,0} = R) \}$$

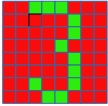
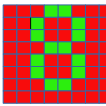
Unless one of the literals in S_2 is included in the explanation C' ,


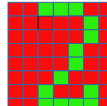
we can't distinguish between  and 

Finding updated C'

$$S_1 = \{(x_{2,7} = G), (x_{2,6} = R), (x_{2,5} = R), (x_{3,4} = R), (x_{2,3} = R), \\ (x_{2,2} = R), (x_{2,1} = G), (x_{3,1} = R), (x_{4,1} = R), \\ (x_{5,1} = G), (x_{3,0} = G), (x_{4,0} = G)\}$$

$$S_2 = \{(x_{4,3} = R), (x_{5,3} = G), (x_{3,2} = R), (x_{5,2} = G), \\ (x_{2,0} = R), (x_{5,0} = R)\}$$

Unless one of the literals in S_1 is included in the explanation C' ,
we can't distinguish between  and 

Unless one of the literals in S_2 is included in the explanation C' ,
we can't distinguish between  and 

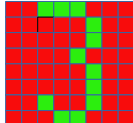
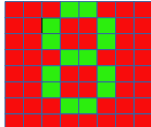
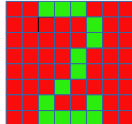
Finding updated C'

$$S_1 = \{(x_{2,7} = G), (x_{2,6} = R), (x_{2,5} = R), (x_{3,4} = R), (x_{2,3} = R), \\ (x_{2,2} = R), (x_{2,1} = G), (x_{3,1} = R), (x_{4,1} = R), \\ (x_{5,1} = G), (x_{3,0} = G), (x_{4,0} = G)\}$$

$$S_2 = \{(x_{4,3} = R), (x_{5,3} = G), (x_{3,2} = R), (x_{5,2} = G), \\ (x_{2,0} = R), (x_{5,0} = R)\}$$

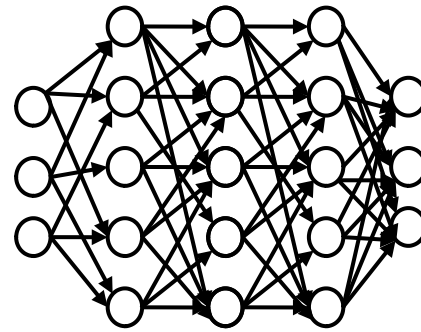
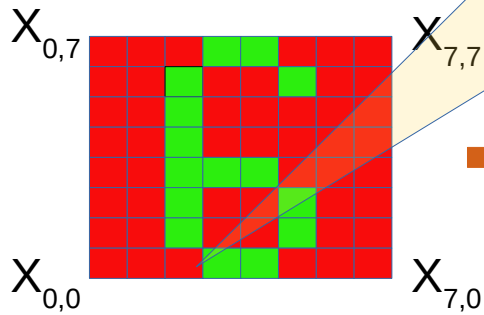
Find a minimum hitting set of S_1 and S_2

$$C' = (x_{3,0} = G) \wedge (x_{2,0} = R)$$

Certainly distinguishes  from both  and 

So, have we found the explanation?

Does $(x_{2,0} = R) \wedge (x_{3,0} = G) \models \mathcal{F} \Rightarrow \mathcal{E}$ hold?



- \bigcirc y_1
- \bullet y_2
- \bigcirc y_3

Clearly not!

Continuing the process

- Find difference with current counterexample
- Find another set S_3 from which we must choose a literal
- Find hitting set C' of S_1, S_2, S_3, \dots
- Check if C' serves as an abductive explanation
 - Does $C' \models \mathcal{F} \Rightarrow \mathcal{E}$?
- If not, repeat above steps
- If yes, output C' as minimal explanation