On minimal odd rankings for Büchi complementation ⋆

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Abstract. We study minimal odd rankings (as defined by Kupferman and Vardi[KV01]) for run-DAGs of words in the complement of a nondeterministic Büchi automaton. We present an optimized version of the ranking based complementation construction of Friedgut, Kupferman and Vardi [FKV06] and Schewe's[Sch09] variant of it, such that every accepting run of the complement automaton assigns a minimal odd ranking to the corresponding run-DAG. This allows us to determine minimally inessential ranks and redundant slices in ranking-based complementation constructions. We exploit this to reduce the size of the complement Büchi automaton by eliminating all redundant slices. We demonstrate the practical importance of this result through a set of experiments using the NuSMV model checker.

1 Introduction

The problem of complementing nondeterministic ω-word automata is fundamental in the theory of automata over infinite words. In addition to the theoretical aspects of the study of complementation techniques, efficient complementation techniques are extremely useful in practical applications as well. Vardi's excellent survey paper on the saga of Büchi complementation spanning more than 45 years provides a brief overview of various such applications of complementation techniques [Var07].

Various complementation constructions for nondeterministic Büchi automata on words (henceforth referred to as NBW) have been developed over the years, starting with Büchi [Büc62]. Büchi's algorithm resulted in a complement automaton with $2^{2^{O(n)}}$ states, starting from an NBW with n states. This upper bound was improved to $2^{O(n^2)}$ by Sistla et al [SVW87]. Safra [Saf88] provided the first asymptotically optimal $n^{O(n)}$ upper bound for complementation that passes through determinization. By a theorem of Michel [Mic88], it was known that Büchi complementation has a $n!$ lower-bound. With this, Löding [Löd99] showed that Safra's construction is asymptotically optimal for Büchi determinization, and hence for complementation. The $O(n!)$ (approximately $(0.36n)^n$) lower bound for Büchi complementation was recently sharpened to $(0.76n)^n$

⋆ This version rectifies a minor error in the version of this paper that appears in Proceedings of ATVA 2009 by Yan [Yan08] using a full-automata technique. The complementation constructions of Klarlund [Kla91], Kupferman and Vardi [KV01] and Kähler and Wilke [KW08] for nondeterministic Büchi automata are examples of determinization free (or *Safraless* as they are popularly called) complementation constructions for Büchi automata. The best known upper bound for the problem until recently was $(0.97n)^n$, given by Kupferman and Vardi. This was recently sharpened by Schewe[Sch09] to an almost tight upper bound of $(0.76n)^n$ modulo a factor of n^2.

NBW complementation techniques based on optimized versions of Safra's determinization construction (see, for example, Piterman's recent work [Pit07]) have been experimentally found to work well for automata of small sizes (typically $8-10$ states) [TCT$^+$08]. However, these techniques are complex and present fewer opportunities for optimized implementations. Ranking-based complementation constructions [KV01,FKV06,Sch09] are comparatively simpler and appear more amenable to optimizations, especially when dealing with larger automaton sizes. For example, several optimization techniques for ranking based complementation constructions have been proposed recently [FKV06,GKSV03]. Similarly, language universality and containment checking techniques that use the framework of ranking-based complementation but avoid explicit construction of complement automata have been successfully applied to NBW with more than 100 states [DR09,FV09]. This leads us to believe that ranking-based complementation constructions hold much promise, and motivates our study of new optimization techniques for such constructions.

The primary contributions of this paper can be summarized as follows : (i) We present an improvement to the ranking based complementation constructions of [FKV06] and [Sch09] for NBW. All accepting runs of our automaton on a word in the complement language correspond to a minimal odd ranking of the run-DAG. (ii) We show how to reduce the size of the complement automaton by efficiently identifying and removing redundant slices without language containment checks. (iii) We present an implementation of the proposed technique using the BDD-based symbolic model checker NuSMV, and experimentally demonstrate the advantages of our technique on a set of examples.

2 Ranking-based NBW complementation

Let $A = (Q, q_0, \Sigma, \delta, F)$ be an NBW, where Q is a set of states, $q_0 \in Q$ is an initial state, Σ is an alphabet, $\delta : Q \times \Sigma \to 2^Q$ is a transition function, and $F \in 2^Q$ is a set of accepting states. An NBW accepts a set of ω-words, where an ω-word α is an infinite sequence $\alpha_0\alpha_1 \ldots$, and $\alpha_i \in \Sigma$ for all $i \geq 0$. A run ρ of A on α is an infinite sequence of states given by $\rho : \mathbb{N} \to Q$, where $\rho(0) = q_0$ and $\rho(i+1) \in \delta(\rho(i), \alpha_i)$ for all $i \geq 0$. A run ρ of A on α is called *accepting* if $\mathit{inf}(\rho) \cap F \neq \emptyset$, where $\mathit{inf}(\rho)$ is the set of states that appear infinitely often along ρ. The run ρ is called *rejecting* if $\mathit{inf}(\rho) \cap F = \emptyset$. An ω-word α is accepted by A if A has an accepting run on it, and is rejected otherwise. The set of all words accepted by A is called the *language* of A, and is denoted $L(A)$. The complementation problem for NBW is to construct an automaton A^c from a given NBW A such that $L(A^c) = \Sigma^\omega \setminus L(A)$. We will henceforth denote the complement language $\Sigma^\omega \setminus L(A)$ by $\overline{L(A)}$. An NBW is said to be *complete* if every state has at least one outgoing transition on every letter in Σ. Every NBW can be made complete without changing the accepted language by adding at most one non-accepting "sink" state. All NBW considered in the remainder of this paper are assumed to be complete.

The (possibly infinite) set of all runs of an NBW $A = (Q, q_0, \Sigma, \delta, F)$ on a word α can be represented by a directed acyclic graph $G_\alpha = (V, E)$, where V is a subset of $Q \times \mathbb{N}$ and $E \subseteq V \times V$. The root vertex of the DAG is $(q_0, 0)$. For all $i > 0$, vertex $(q, i) \in V$ iff there is a run ρ of A on α such that $\rho(i) = q$. The set of edges of G_α is $E \subseteq V \times V$, where $((q, i), (q', j)) \in E$ iff both (q, i) and (q', j) are in V, $j = i + 1$ and $q' \in \delta(q, \alpha_i)$. Graph G_α is called the *run-DAG* of α in A. A vertex $(q, l) \in V$ is called an *F-vertex* if $q \in F$, i.e, q is a final state of A. A vertex (q, l) is said to be *F-free* if there is no F-vertex that is reachable from (q, l) in G_α. Furthermore, (q, l) is called *finite* if only finitely many vertices are reachable from (q, l) in G_α. For every $l \geq 0$, the set of vertices $\{(q, l) \mid (q, l) \in V\}$ constitutes *level* l of G_α. An accepting path in G_α is an infinite path $(q_0, 0), (q_{i_1}, 1), (q_{i_2}, 2) \ldots$ such that $q_0, q_{i_1}, \ldots$ is an accepting run of A. The run-DAG G_α is called rejecting if there is no accepting path in G_α. Otherwise, G_α is said to be accepting.

2.1 Ranking functions and complementation

Kupferman and Vardi [KV01] introduced the idea of assigning ranks to vertices of run-DAGs, and described a rank-based complementation construction for alternating Büchi automata. They also showed how this technique can be used to obtain a ranking-based complementation construction for NBW, that is easier to understand and implement than the complementation construction based on Safra's determinization construction[Saf88]. In this section, we briefly overview ranking-based complementation constructions for NBW.

Let $[k]$ denote the set $\{1, 2, \ldots, k\}$, and $[k]^{odd}$ (respectively $[k]^{even}$) denote the set of all odd (respectively even) numbers in the set $\{1, 2, \ldots, k\}$. Given an NBW A with n states and an ω-word α, let $G_\alpha = (V, E)$ be the run-DAG of α in A. A *ranking* r of G_α is a function $r : V \to [2n]$ that satisfies the following conditions: (i) for all vertices $(q, l) \in V$, if $r((q, l))$ is odd then $q \notin F$, and (ii) for all edges $((q, l), (q', l + 1)) \in E$, we have $r((q', l + 1)) \leq r((q, l))$. A ranking associates with every vertex in G_α a rank in $[2n]$ such that the ranks along every path in G_α are non-increasing, and vertices corresponding to final states always get even ranks. A ranking r is said to be *odd* if every infinite path in G_α eventually gets trapped in an odd rank. Otherwise, r is called an *even ranking.* We use $max\_odd(r)$ to denote the highest odd rank in the range of r.

A *level ranking* for A is a function $g : Q \to [2n] \cup \{\perp\}$ such that for every $q \in Q$, if $g(q) \in [2n]^{odd}$, then $q \notin F$. Let $\mathcal{L}$ be the set of all level rankings for A. Given two level rankings $g, g' \in \mathcal{L}$, a set $S \subseteq Q$ and a letter σ, we say that g' covers (g, S, σ) if for all $q \in S$ and $q' \in \delta(q, \sigma)$, if $g(q) \neq \perp$, then $g'(q') \neq \perp$

and $g'(q') \leq g(q)$. For a level ranking g, we abuse notation and let $max\_odd(g)$ denote the highest odd rank in the range of g. A ranking r of G_α induces a level ranking for every level $l \geq 0$ of G_α. If $Q_l = \{q \mid (q,l) \in V\}$ denotes the set of states in level l of G_α, then the level ranking g induced by r for level l is as follows: $g(q) = r((q,l))$ for all $q \in Q_l$ and $g(q) = \perp$ otherwise. It is easy to see that if g and g' are level rankings induced for levels l and $l+1$ respectively, then g' covers (g, Q_l, α_l), where α_l is the l^{th} letter in the input word α. A level ranking g is said to be *tight* if the following conditions hold: (i) the highest rank in the range of g is odd, and (ii) for all $i \in [max\_odd(g)]^{odd}$, there is a state $q \in Q$ with $g(q) = i$.

Lemma 1 ([KV01]). *The following statements are equivalent:*

(P1) All paths of G_α see only finitely many F-vertices.
(P2) There is an odd ranking for G_α.

Kupferman and Vardi [KV01] provided a constructive proof of (P1) $\Rightarrow$ (P2) in the above Lemma. Their construction is important for some of our subsequent discussions, hence we outline it briefly here. Given an NBW A with n states, an ω-word $\alpha \in \overline{L(A)}$ and the run-DAG G_α of A on α, the proof in [KV01] inductively defines an infinite sequence of DAGs $G_0 \supseteq G_1 \supseteq \ldots$, where (i) $G_0 = G_\alpha$, (ii) $G_{2i+1} = G_{2i} \setminus \{(q,l) \mid (q,l)$ is finite in $G_{2i}\}$, and (iii) $G_{2i+2} = G_{2i+1} \setminus \{(q,l) \mid (q,l)$ is F-free in $G_{2i+1}\}$, for all $i \geq 0$. An interesting consequence of this definition is that for all $i \geq 0$, G_{2i+1} is either empty or has no finite vertices. It can be shown that if all paths in G_α see only finitely many F-vertices, then G_{2n-1} and all subsequent G_is must be empty. A ranking $\mathsf{r}^{\mathsf{KV}}_{\mathsf{A},\alpha}$ of G_α can therefore be defined as follows: for every vertex (q,l) of G_α, $\mathsf{r}^{\mathsf{KV}}_{\mathsf{A},\alpha}((q,l)) = 2i$ if (q,l) is finite in G_{2i}, and $\mathsf{r}^{\mathsf{KV}}_{\mathsf{A},\alpha}((q,l)) = 2i+1$ if (q,l) is F-free in G_{2i+1}. Kupferman and Vardi showed that $\mathsf{r}^{\mathsf{KV}}_{\mathsf{A},\alpha}$ is an odd ranking [KV01]. Throughout this paper, we will use $\mathsf{r}^{\mathsf{KV}}_{\mathsf{A},\alpha}$ to denote the odd ranking computed by the above technique due to Kupferman and Vardi (hence KV in the superscript) for NBW A and $\alpha \in \overline{L(A)}$. When A and α are clear from the context, we will simply use r^{KV} for notational convenience.

The NBW complementation construction and upper size bound presented in [KV01] was subsequently tightened in [FKV06], where the following important observation was made.

Lemma 2 ([FKV06]). *Given a word $\alpha \in \overline{L(A)}$, there exists an odd ranking r of G_α and a level $l_{lim} \geq 0$, such that for all levels $l > l_{lim}$, the level ranking induced by r for l is tight.*

Lemma 2 led to a reduced upper bound for the size of ranking-based complementation constructions, since all non-tight level rankings could now be ignored after reading a finite prefix of the input word. Schewe[Sch09] tightened the construction and analysis further, resulting in a ranking-based complementation construction with an upper size bound that is within a factor of n^2 of the best known lower bound [Yan08]. Hence, Schewe's construction is currently the best known ranking-based construction for complementing NBW. Gurumurthy et al [GKSV03] presented a collection of practically useful optimization techniques for keeping the size of complement automata constructed using ranking techniques under control. Their experiments demonstrated the effectiveness of their optimizations for NBW with an average size of 6 states. Interestingly, their work also highlighted the difficulty of complementing NBW with tens of states in practice. Doyen and Raskin [DR09] have recently proposed powerful anti-chain optimizations in ranking-based techniques for checking universality ($L(A) =^? \Sigma^\omega$) and language containment ($L(A) \subseteq^? L(B)$) of NBW. Fogarty and Vardi [FV09] have evaluated Doyen and Raskin's technique and also Ramsey-based containment checking techniques in the context of proving size-change termination (SCT) of programs. Their results bear testimony to the effectiveness of Doyen and Raskin's anti-chain optimizations for ranking-based complementation in SCT problems, especially when the original NBW is known to have *reverse-determinism* [FV09].

Given an NBW A, let $\mathsf{KVF}(A)$ be the complement NBW constructed using the Friedgut, Kupferman and Vardi construction with tight level rankings [KV01,FKV06]. For notational convenience, we will henceforth refer to this construction as *KVF-construction*. Similarly, let $\mathsf{KVFS}(A)$ be the complement automaton constructed using Schewe's variant[Sch09] of the KVF-construction. We will henceforth refer to this construction as *KVFS-construction*.

The work presented in this paper can be viewed as an optimized variant of Schewe's [Sch09] ranking-based complementation construction. The proposed method is distinct from other optimization techniques proposed in the literature (e.g. those in [GKSV03]), and adds to the repertoire of such techniques. We first show that for every NBW A and word $\alpha \in \overline{L(A)}$, the ranking r^{KV} is *minimal* in the following sense: if r is any odd ranking of G_α, then every vertex (q, l) in G_α satisfies $\mathsf{r}^{\mathsf{KV}}((q, l)) \leq r((q, l))$. We then describe how to restrict the transitions of the complement automaton obtained by the KVFS-construction, such that every accepting run of α assigns the same rank to all vertices in G_α as is assigned by r^{KV}. Thus, our construction ensures that acceptance of α happens only through minimal odd rankings. This allows us to partition the set of states of the complement automaton into *slices* such that for every word $\alpha \in \overline{L(A)}$, all its accepting runs lie in exactly one slice. Redundant slices can then be identified as those that never contribute to accepting any word in $\overline{L(A)}$. Removal of such redundant slices results in a reduced state count, while preserving the language of the complement automaton. The largest $k(> 0)$ such that there is a non-redundant slice with that assigns rank k to some vertex in the run-DAG gives the *rank of* A, as defined by [GKSV03]. Notice that our sliced view of the complement automaton is distinct from the notion of slices as used in [KW08].

Gurumurthy et al have shown [GKSV03] that for every NBW A, there exists an NBW B with $L(B) = L(A)$, such that both the KVF- and KVFS-constructions for B^c require at most 3 ranks. However, obtaining B from A is non-trivial, and requires an exponential blowup in the worst-case [FKV06]. Therefore, ranking-based complementation constructions typically focus on reducing the state count of the complement automaton starting from a given NBW A, instead of first computing B and then constructing B^c. We follow the same approach in this paper. Thus, we do not seek to obtain an NBW with the *minimum* rank for the complement of a given ω-regular language. Instead, we wish to reduce the state count of Kupferman and Vardi's rank-based complementation construction, starting from a given NBW A.

3 Minimal odd rankings

Given an NBW A and an ω-word $\alpha \in \overline{L(A)}$, an odd ranking r of G_α is said to be *minimal* if for every odd ranking r' of G_α, we have $r'((q,l)) \geq r((q,l))$ for all vertices (q,l) in G_α.

Theorem 1. *For every NBW A and ω-word $\alpha \in \overline{L(A)}$, the ranking $\mathsf{r}^{\mathsf{KV}}_{\mathsf{A},\alpha}$ is minimal.*

Proof. Let α be an ω-word in $\overline{L(A)}$. Since A and α are clear from the context, we will use the simpler notation r^{KV} to denote the ranking computed by Kupferman and Vardi's method. Let r be any (other) odd ranking of G_α, and let $V_{r,i}$ denote the set of vertices in G_α that are assigned the rank i by r. Since A is assumed to be a complete automaton, there are no finite vertices in G_α. Hence, by Kupferman and Vardi's construction, $V_{\mathsf{r}^{\mathsf{KV}},0} = \emptyset$. Note that this is consistent with our requirement that all ranking functions have range $[2n] = \{1, \ldots 2n\}$.

We will prove the theorem by showing that $V_{r,i} \subseteq \bigcup_{k=1}^{i} V_{\mathsf{r}^{\mathsf{KV}},k}$ for all $i > 0$. The proof proceeds by induction on i, and by following the construction of DAGs $G_0, G_1, \ldots$ in Kupferman and Vardi's proof of Lemma 1.

Base case: Consider $G_1 = G_0 \setminus \{(q',l') \mid (q',l') \text{ is finite in } G_0\} = G_\alpha \setminus \emptyset = G_\alpha$. Let (q_1,l_1) be a vertex in G_1 such that $r((q_1,l_1)) = 1$. Let (q_f,l_f) be an F-vertex reachable from (q_1,l_1) in G_α, if possible. By virtue of the requirements that F-vertices must get even ranks, and ranks cannot increase along any path, $r((q_f,l_f))$ must be < 1. However, this is impossible given that the range of r must be $[2n]$. Therefore, no F-vertex can be reachable from (q_1,l_1). In other words, (q_1,l_1) is F-free in G_1. Hence, by definition, we have $\mathsf{r}^{\mathsf{KV}}(q_1,l_1) = 1$. Thus, $V_{r,1} \subseteq V_{\mathsf{r}^{\mathsf{KV}},1}$.

Hypothesis: Assume that $V_{r,j} \subseteq \bigcup_{k=1}^{j} V_{\mathsf{r}^{\mathsf{KV}},k}$ for all $1 \leq j \leq i$.

Induction: By definition, $G_{i+1} = G_\alpha \setminus \bigcup_{s=1}^{i} V_{\mathsf{r}^{\mathsf{KV}},s}$. Let (q_{i+1},l_{i+1}) be a vertex in G_{i+1} such that $r((q_{i+1},l_{i+1})) = i+1$. Since r is an odd ranking, all paths starting from (q_{i+1},l_{i+1}) must eventually get trapped in some odd rank $\leq i+1$ (assigned by r). We consider two cases.

- Suppose there are no infinite paths starting from (q_{i+1},l_{i+1}) in G_{i+1}. This implies (q_{i+1},l_{i+1}) is a finite vertex in G_{i+1}. We have seen earlier that for all $k \geq 0$, G_{2k+1} must be either empty or have no finite vertices. Therefore, $i+1$ must be even, and $\mathsf{r}^{\mathsf{KV}}((q_{i+1},l_{i+1})) = i+1$ by Kupferman and Vardi's construction.

– There exists a non-empty set of infinite paths starting from (q_{i+1}, l_{i+1}) in G_{i+1}. Since $G_{i+1} = G_\alpha \setminus \bigcup_{s=1}^{i} V_{\mathsf{r}^{\mathsf{KV}},s}$, none of these paths reach any vertex in $\bigcup_{s=1}^{i} V_{\mathsf{r}^{\mathsf{KV}},s}$. Since $V_{r,j} \subseteq \bigcup_{k \in \{1,2,\ldots,j\}} V_{\mathsf{r}^{\mathsf{KV}},k}$ for all $1 \leq j \leq i$, and since ranks cannot increase along any path, it follows that r must assign $i+1$ to all vertices along each of the above paths. This, coupled with the fact that r is an odd ranking, implies that $i+1$ is odd. Since F-vertices must be assigned even ranks by r, it follows from above that (q_{i+1}, l_{i+1}) is F-free in G_α. Therefore, $\mathsf{r}^{\mathsf{KV}}((q_{i+1}, l_{i+1})) = i+1$ by Kupferman and Vardi's construction.

We have thus shown that all vertices in G_{i+1} that are assigned rank $i+1$ by r must also be assigned rank $i+1$ by r^{KV}. Therefore, $V_{r,i+1} \setminus \bigcup_{j=1}^{i} V_{\mathsf{r}^{\mathsf{KV}},j} \subseteq V_{\mathsf{r}^{\mathsf{KV}},i+1}$. In other words, $V_{r,i+1} \subseteq \bigcup_{j=1}^{i+1} V_{\mathsf{r}^{\mathsf{KV}},j}$.

By the principle of mathematical induction, it follows that $V_{r,i} \subseteq \bigcup_{k=1}^{i} V_{\mathsf{r}^{\mathsf{KV}},k}$ for all $i > 0$. Thus, if a vertex is assigned rank i by r, it must be assigned a rank $\leq i$ by r^{KV}. Hence r^{KV} is minimal.

□

Lemma 3. *For every $\alpha \in \overline{L(A)}$, the run-DAG G_α ranked by $\mathsf{r}^{\mathsf{KV}}_{\mathsf{A},\alpha}$ (or simply r^{KV}) satisfies the following properties.*

1. *For every vertex (q, l) that is not an F-vertex such that $\mathsf{r}^{\mathsf{KV}}((q, l)) = k$, there must be at least one immediate successor $(q', l+1)$ such that $\mathsf{r}^{\mathsf{KV}}((q', l+1)) = k$.*
2. *For every vertex (q, l) that is an F-vertex, such that $\mathsf{r}^{\mathsf{KV}}((q, l)) = k$, there must be atleast one immediate successor $(q', l+1)$ such that $\mathsf{r}^{\mathsf{KV}}((q', l+1)) = k$ or $\mathsf{r}^{\mathsf{KV}}((q', l+1)) = k-1$*
3. *For every vertex (q, l) such that $\mathsf{r}^{\mathsf{KV}}((q, l)) = k$, where k is odd and > 1, there is a vertex (q', l') for $l' > l$ such that (q', l') is an F-vertex reachable from (q, l) and $\mathsf{r}^{\mathsf{KV}}((q', l')) = k-1$.*
4. *For every vertex (q, l) such that $\mathsf{r}^{\mathsf{KV}}((q, l)) = k$, where k is even and > 0, every path starting at (q, l) eventually visits a vertex with rank less than k. Furthermore, there is a vertex (q', l') for $l' > l$ such that (q', l') is reachable from (q, l) and $\mathsf{r}^{\mathsf{KV}}((q', l')) = k-1$.*

We omit the proof due to lack of space. The reader is referred to [KC09] for details of the proof.

4 A motivating example

We have seen above the KVFS-construction leads to almost tight worst case bounds for NBW complementation. This is a significant achievement considering the long history of Büchi complementation [Var07]. However, the KVFS-construction does not necessarily allow us to construct small complement automata for every NBW. Specifically, there exists a family of NBW $\mathcal{A} = \{ A_3, A_5, \ldots \}$ such that for every $i \in \{3, 5, \ldots\}$: (i) A_i has i states, (ii) each of the ranking-based complementation constructions in [KV01], [FKV06] and [Sch09] produces a complement automaton with at least $\frac{i^{(\frac{i-1}{2})}}{e^i}$ states, and (iii) a ranking-based complementation construction that assigns minimal ranks to all run-DAGs results in a complement automaton with $\Theta(i)$ states.

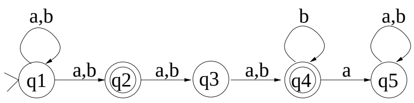

Fig. 1. Automaton A_5

Automata in the family $\mathcal{A}$ can be described as follows. Each A_i is an NBW $(Q_i, q_1, \Sigma, \delta_i, F_i)$, where $Q_i = \{q_1, q_2, \ldots, q_i\}$, $\Sigma = \{a, b\}$, q_1 is the initial state and $F_i = \{q_j \mid q_j \in Q_i, j \text{ is even }\}$. The transition relation δ_i is given by: $\delta_i = \{(q_j, a, q_{j+1}), (q_j, b, q_{j+1}) \mid 1 \leq j \leq i-2\} \cup \{(q_1, a, q_1), (q_1, b, q_1), (q_i, a, q_i), (q_i, b, q_i), (q_{i-1}, b, q_{i-1}), (q_{i-1}, a, q_i)\}$. Figure 1 shows the structure of automaton A_5 defined in this manner. Note that each $A_i \in \mathcal{A}$ is a complete automaton.Furthermore, $b^\omega \in L(A_i)$ and $a^\omega \notin L(A_i)$ for each $A_i \in \mathcal{A}$. Let $\mathsf{KVF}(A_i)$ and $\mathsf{KVFS}(A_i)$ be the complement automata for A_i constructed using the KVF-construction and KVFS-construction respectively.

Lemma 4. *For every $A_i \in \mathcal{A}$, the number of states in $\mathsf{KVF}(A_i)$ and $\mathsf{KVFS}(A_i)$ is atleast $\frac{(i)^{(i-1)/2}}{e^i}$. Furthermore, there exists a ranking-based complementation construction for A_i that gives a complement automaton ${A_i}'$ with $\Theta(i)$ states.*

Proof sketch: The proof is obtained by considering the run-DAG for a^ω (which is $\in \overline{L(A_i)}$), and by showing that at least $\frac{(i)^{(i-1)/2}}{e^i}$ states are required in both the KVF- and KVFS-constructions in order to allow all possible consistent rank assignments to vertices of the run-DAG. On the other hand, a complementation construction that uses the ranking $f(q_i) = 1$, $f(q_1) = 3$ and $f(q_j) = 2$ for all $2 \leq j \leq i-1$ at all levels of the run-DAG can be shown to accept $\overline{L(A)}$ with $\Theta(i)$ states. Details of the proof may be found in [KC09].

This discrepancy in the size of a sufficient rank set and the actual set of ranks used by the KVF- and KVFS-constructions motivates us to ask if we can devise a ranking-based complementation construction for NBW that uses the minimum number of ranks when accepting a word in the complement language. In this paper, we answer this question in the affirmative, by providing such a complementation construction.

5 Complementation with minimal ranks

Motivated by the example described in the previous section, we now describe an optimized ranking-based complementation construction for NBW. Given an NBW A, the complement automaton A' obtained using our construction has the special property that when it accepts an ω-word α, it assigns a rank r to G_α

that agrees with the ranking $\mathsf{r}^{\mathsf{KV}}_{\mathsf{A},\alpha}$ at every vertex in G_α. This is achieved by non-deterministically mimicking the process of rank assignment used to arrive at $\mathsf{r}^{\mathsf{KV}}_{\mathsf{A},\alpha}$. Our construction imposes additional constraints on the states and transitions of the complement automaton, beyond those in the KVF- and KVFS-constructions. For example, if k is the smallest rank that can be assigned to vertex (q,l), then the following conditions must hold: (i) if (q,l) is not an F-vertex then it must have a successor of the same rank at the next level, and (ii) if (q,l) is an F-vertex then it must have a successor at the next level with rank k or $k-1$. The above observations are coded as conditions on the transitions of A', and are crucial if every accepting run of A' on $\alpha \in \overline{L(A)}$ must correspond to the unique ranking $\mathsf{r}^{\mathsf{KV}}_{\mathsf{A},\alpha}$ of G_α.

Recall that in [FKV06], a state of the automaton that tracks the ranking of the run-DAG vertices at the current level is represented as a triple (S,O,f), where S is a set of Büchi states reachable after reading a finite prefix of the word, f is a tight level ranking, and the O-set checks if all even ranked vertices in a (possibly previous) level of the run-DAG have moved to lower odd ranks. In our construction, we use a similar representation of states, although the O-set is much more versatile. Specifically, the O-set is populated turn-wise with vertices of the same rank k, for both odd and even k. This is a generalization of Schewe's technique that uses the O-set to check if even ranked vertices present at a particular level have paths to vertices with lower odd ranks. The O-set in our construction, however, does more. It checks if every vertex of rank k (whether even or odd) in an O-set eventually reaches a vertex with rank $k-1$. If k is even. then it also checks if all runs starting at vertices in O eventually reach a vertex with rank $k-1$. When all vertices in an O-set tracking rank 2 reach vertices with rank 1, the O-set is reset and loaded with vertices that have the maximal odd rank in the range of the current level ranking. The process of checking ranks is then re-started. This gives rise to the following construction for the complement automaton A'.

Let $A' = (Q', Q'_0, \Sigma, \delta', F')$, where

- $Q' = \{2^Q \times 2^Q \times \mathcal{R}\}$ is the state set, such that if $(S,O,f) \in Q'$, then $S \subseteq Q$ and $S \neq \emptyset$, $f \in \mathcal{R}$ is a level ranking, $O \subseteq S$, and either $O = \emptyset$ or $\exists k \in [2n-1]\ \forall q \in O,\ f(q) = k$.
- $Q'_0 = \bigcup_{i \in [2n-1]} \{(S,O,f) \mid S = \{q_0\},\ f(q_0) = i, O = \emptyset\}$ is the set of initial states.
- For every $\sigma \in \Sigma$, the transition function δ' is defined such that if $(S',O',f') \in \delta'((S,O,f),\sigma)$, the following conditions are satisfied.
  1. $S' = \delta(S,\sigma)$, f' covers (f,S,σ).
  2. For all $q \in S \setminus F$, there is a $q' \in \delta(q,\sigma)$ such that $f'(q') = f(q)$.
  3. For all $q \in S \cap F$ one of the following must hold
     (a) There is a $q' \in \delta(q,\sigma)$, such that $f'(q') = f(q)$.
     (b) There is a $q' \in \delta(q,\sigma)$, such that $f'(q') = f(q) - 1$.
  4. In addition, we have the following restrictions:
     (a) If f is not a tight level ranking, then $O' = O$
     (b) If f is a tight level ranking and $O \neq \emptyset$, then

i. If for all $q \in O$, we have $f(q) = k$, where k is even, then
   - Let $O_1 = \delta(O, \sigma) \setminus \{q \mid f'(q) < k\}$.
   - If $(O_1 = \emptyset)$ then $O' = \{q \mid q \in Q \wedge f'(q) = k - 1\}$, else $O' = O_1$.

ii. If for all $q \in O$, we have $f(q) = k$, where k is odd, then
   - If $k = 1$ then $O' = \emptyset$
   - If $k > 1$ then let $O_2 = O \setminus \{q \mid \exists q' \in \delta(q, \sigma),\ (f'(q') = k - 1)\}$.
     - If $O_2 \neq \emptyset$ then $O' \subseteq \delta(O_2, \sigma)$ such that $\forall q \in O_2\, \exists q' \in O',\ \ q' \in \delta(q, \sigma)\ \wedge\ f'(q') = k$.
     - If $O_2 = \emptyset$ then $O' = \{q \mid f'(q) = k - 1\}$.

(c) If f is a tight level ranking and $O = \emptyset$, then $O' = \{q \mid f'(q) = max\_odd(f')\}$.

– As in the KVF- and KVFS-constructions, the set of accepting states of A' is $F' = \{(S, O, f) \mid f$ is a tight level ranking, and $O = \emptyset\}$

Note that unlike the KVF- and KVFS-constructions, the above construction does not have an initial phase of unranked subset construction, followed by a non-deterministic jump to ranked subset construction with tight level rankings. Instead, we start directly with ranked subsets of states, and the level rankings may indeed be non-tight for some finite prefix of an accepting run. The value of O is inconsequential until the level ranking becomes tight; hence it is kept as $\emptyset$ during this period. Note further that the above construction gives rise to multiple initial states in general. Since an NBW with multiple initial states can be easily converted to one with a single initial state without changing its language, this does not pose any problem, and we will not dwell on this issue any further.

Theorem 2. $L(A') = \overline{L(A)}$

The proof proceeds by establishing three sub-results: (i) every accepting run of A' on word α assigns an odd ranking to the run-DAG G_α and hence corresponds to an accepting run of $\mathsf{KVF}(A)$, (ii) the run corresponding to the ranking $\mathsf{r}^{\mathsf{KV}}_{\mathsf{A},\alpha}$ is an accepting run of A' on α, and (iii) $L(\mathsf{KVF}(A)) = \overline{L(A)}$[KV01]. Details of the proof may be found in [KC09].

Given an NBW A and the complement NBW A' constructed using the above algorithm, we now ask if A' has an accepting run on some $\alpha \in \overline{L(A)}$ that induces an odd ranking r different from $\mathsf{r}^{\mathsf{KV}}_{\mathsf{A},\alpha}$. We answer this question negatively in the following lemma.

Lemma 5. *Let $\alpha \in \overline{L(A)}$, and let r be the odd ranking corresponding to an accepting run of A' on α. Let $V_{r,i}$ (respectively, $V_{\mathsf{r}^{\mathsf{KV}}_{\mathsf{A},\alpha},i}$) be the set of vertices in G_α that are assigned rank i by r (respectively, $\mathsf{r}^{\mathsf{KV}}_{\mathsf{A},\alpha}$). Then $V_{r,i} = V_{\mathsf{r}^{\mathsf{KV}}_{\mathsf{A},\alpha},i}$ for all $k > 0$.*

Proof. We prove the claim by induction on the rank i. Since A and α are clear from the context, we will use r^{KV} in place of $\mathsf{r}^{\mathsf{KV}}_{\mathsf{A},\alpha}$ in the remainder of the proof. *Base case:* Let $(q, l) \in V_{\mathsf{r}^{\mathsf{KV}},1}$. By definition, (q, l) is F-free. Suppose $r((q, l)) = m$, where $m > 1$, if possible. If m is even, the constraints embodied in steps (2), (3) and (4(b)i) of our complementation construction, coupled with the fact that the O-set becomes $\emptyset$ infinitely often, imply that (q, l) has an F-vertex descendant (q', l') in G_α. Therefore, (q, l) is not F-free – a contradiction! Hence m cannot be even.

Suppose m is odd and > 1. The constraint embodied in step (4(b)ii) of our construction, and the fact that the O-set becomes $\emptyset$ infinitely often imply that (q, l) has a descendant (q'', l'') that is assigned an even rank by r. The constraints embodied in steps (2), (3) and (4(b)i), coupled with the fact that the O-set becomes $\emptyset$ infinitely often, further imply that (q'', l'') has an F-vertex descendant in G_α. Hence (q, l) has an F-vertex descendant, and is not F-free. This leads to a contradiction again! Therefore, our assumption must have been incorrect, i.e. $r((q, l)) \leq 1$. Since 1 is the minimum rank in the range of r, we finally have $r((q, l)) = 1$. This shows that $V_{\mathsf{r}^{\mathsf{KV}},1} \subseteq V_{r,1}$.

Now suppose $(q, l) \in V_{r,1}$. Since r corresponds to an accepting run of A', it is an odd-ranking. This, coupled with the fact that ranks cannot increase along any path in G_α, imply that all descendants of (q, l) in G_α are assigned rank 1 by r. Since F-vertices must be assigned even ranks, this implies that (q, l) is F-free in G_α. It follows that $\mathsf{r}^{\mathsf{KV}}((q, l)) = 1$. Therefore, $V_{r,1} \subseteq V_{\mathsf{r}^{\mathsf{KV}},1}$. From the above two results, we have $V_{r,1} = V_{\mathsf{r}^{\mathsf{KV}},1}$.

Hypothesis: Assume that $V_{r,j} = V_{\mathsf{r}^{\mathsf{KV}},j}$ for $1 \leq j \leq i$.

Induction: Let $(q, l) \in V_{\mathsf{r}^{\mathsf{KV}},i+1}$. Then by the induction hypothesis, (q, l) cannot be in any $V_{r,j}$ for $j \leq i$. Suppose $r((q, l)) = m$, where $m > i + 1$, if possible. We have two cases.

1. $i + 1$ is odd: In this case, the constraints embodied in steps (2), (3), (4(b)i) and (4(b)ii) of our construction, coupled with the fact that the O-set becomes $\emptyset$ infinitely often, imply that vertex (q, l) has an F-vertex descendant (q', l') (possibly its own self) such that (i) $r((q', l')) = i + 2$, and (ii) vertex (q', l') in turn has a descendant (q'', l'') such that $r((q'', l'')) = i+1$. The constraint embodied in (2) of our construction further implies that there must be an infinite path π starting from (q'', l'') in G_α such that every vertex on π is assigned rank $i+1$ by r, and none of these are F-vertices. Since $\mathsf{r}^{\mathsf{KV}}((q, l)) = i+1$ is odd, and since (q', l') is an F-vertex descendant of (q, l), we must have $\mathsf{r}^{\mathsf{KV}}((q', l')) \leq i$, where i is even. Furthermore, since r^{KV} is an odd ranking, every path in G_α must eventually get trapped in an odd rank assigned by r^{KV}. Hence, eventually every vertex on π is assigned an odd rank $\leq i$ by $\mathsf{r}^{\mathsf{KV}}_{\mathsf{A},\alpha}$. However, we already know that the vertices on π are eventually assigned the odd rank $i + 1$ by r. Hence $V_{r,j} \neq V_{\mathsf{r}^{\mathsf{KV}},j}$ for some $j \in \{1, \ldots i\}$. This contradicts the inductive hypothesis!
2. $i + 1$ is even: In this case, the constraints embodied in steps (2), (3), (4(b)i) and (4(b)ii) of our construction, and the fact that the O-set becomes $\emptyset$ infinitely often, imply that (q, l) has a descendant (q', l') in G_α such that $r((q', l')) = i+2$ (which is odd), and there is an infinite path π starting from (q', l') such that all vertices on π are assigned rank $i+2$ by r. However, since r^{KV} is an odd ranking and since $\mathsf{r}^{\mathsf{KV}}((q, l)) = i + 1$ (which is even), vertices on π must eventually get trapped in an odd rank $\leq i$ assigned by r^{KV}. This

implies that $V_{r,j} \neq V_{\mathsf{r}^{\mathsf{KV}},j}$ for some $j \in \{1, \ldots i\}$. This violates the inductive hypothesis once again!

It follows from the above cases that $r((q,l)) \leq i+1$. However, $(q,l) \notin V_{\mathsf{r}^{\mathsf{KV}},j}$ for $1 \leq j \leq i$ (since $(q,l) \in V_{\mathsf{r}^{\mathsf{KV}},i+1}$), and $V_{r,j} = V_{\mathsf{r}^{\mathsf{KV}},j}$ for $1 \leq j \leq i$ (by inductive hypothesis). Therefore, $(q,l) \notin V_{r,j}$ for $1 \leq j \leq i$. This implies that $r((q,l)) = i+1$, completing the induction.

By the principle of mathematical induction, we $V_{\mathsf{r}^{\mathsf{KV}}_{\mathsf{A},\alpha},i} = V_{r,i}$ for all $i > 0$. □

Theorem 3. *Every accepting run of A' on $\alpha \in \overline{L(A)}$ induces the unique minimal ranking $\mathsf{r}^{\mathsf{KV}}_{\mathsf{A},\alpha}$.*

Proof. Follows from Lemma 5.

5.1 Size of complement automaton

The states of A' are those in the set $\{2^Q \times 2^Q \times \mathcal{R}\}$. While some of these states correspond to tight level rankings, others do not. We first use an extension of the idea in [Sch09] to encode a state (S, O, f) with tight level ranking f as a pair (g, i), where $g : Q \rightarrow \{1, \ldots, r\} \cup \{-1, -2\}$ and $r = \mathit{max\_odd}(f)$. Thus, for all $q \in Q$, we have $q \notin S$ iff $g(q) = -2$. If $q \in S$ and $q \notin O$, we have $g(q) = f(q)$. If $q \in O$ and $f(q)$ is even, then we let $g(q) = -1$ and $i = f(q)$. This part of the encoding is exactly as in [Sch09]. We extend this encoding to consider cases where $q \in O$ and $f(q) = k$ is odd. There are two sub cases to consider: (i) $O \subsetneq \{q \mid q \in S \wedge f(q) = k\}$, and (ii) $O = \{q \mid q \in S \wedge f(q) = k\}$. In the first case, we let $i = k$ and $g(q) = -1$ for all $q \in O$. In the second case, we let $i = k$ and $g(q) = f(q) = k$ for all $q \in O$. Since, f is a tight level ranking, the O-set cannot be empty when we check for states with an odd rank in our construction. Therefore, there is no ambiguity in identifying the set O in both cases (i) and (ii) above. It is now easy to see that g is always onto one of the three sets $\{1, 3, \ldots, r\}$, $\{-1\} \cup \{1, 3, \ldots, r\}$ or $\{-2\} \cup \{1, 3, \ldots, r\}$. By Schewe's analysis [Sch09], the total number of such (g, i) pairs is in $O(\mathit{tight}(n+1))$.

Now, let us consider states with non-tight level rankings. Our construction ensures that once an odd rank i appears in a level ranking g along a run ρ, all subsequent level rankings along ρ contain every rank in $\{i, i+2, \ldots \mathit{max\_odd}(g)\}$. The O-set in states with non-tight level ranking is inconsequential; hence we ignore this. Suppose a state with non-tight level ranking g contains the odd ranks $\{i, i-2, \ldots, j\}$, where $1 < j \leq i = \mathit{max\_odd}(g)$. To encode this state, we first replace g with a level ranking g' as follows. For all $k \in \{j, \ldots, i\}$ and $q \in Q$, if $g(q) = k$, then $g'(q) = k - j + c$, where $c = 0$ if j is even and 1 otherwise. Effectively, this transforms g to a tight level ranking g' by shifting all ranks down by $j - c$. The original state can now be represented as the pair $(g', -(j-c))$. Note that the second component of a state represented as (g, i) is always non-negative for states with tight level ranking, and always negative for states with non-tight level ranking. Hence, there is no ambiguity in decoding the state representation. Clearly, the total no. of states with non-tight rankings is in

$O((n+1) \cdot \mathit{tight}(n))$, i.e., $O(\mathit{tight}(n+1))$. Hence, the size of A' is upper bounded by $O(\cdot \mathit{tight}(n+1))$ which differs from the lower bound of $\Omega(\mathit{tight}(n-1))$ given by [Yan08] by only a factor of n^2.

6 Slices of complement automaton

The transitions of the complement automaton A' obtained by our construction have the property that a state (S, O, f) has a transition to (S', O', f') only if $\mathit{max\_odd}(f) = \mathit{max\_odd}(f')$. Consequently, the set of states Q' can be partitioned into *slices* $Q_1, Q_3, \ldots, Q_{2n-1}$, where the set $Q_i = \{(S, O, f) | (S, O, f) \in Q' \wedge \mathit{max\_odd}(f) = i\}$ is called the i^{th} slice of A'. It is easy to see that $Q' = \bigcup_{i \in [2n-1]^{odd}} Q_i$. Let ρ be an accepting run of A' on $\alpha \in \overline{L(A)}$. We say that ρ is *confined* to a slice Q_i of A' iff ρ sees only states from Q_i. If ρ is confined to slice i, and if r is the odd ranking induced by ρ, then $\mathit{max\_odd}(r) = i$.

Lemma 6. *All accepting runs of A' on $\alpha \in \overline{L(A)}$ are confined to the same slice.*

Proof. Follows from Theorem 3. □

The above results indicate that if a word α is accepted by the i^{th} slice of our automaton A', then it cannot be accepted by $\mathsf{KVF}(A)$ using a tight ranking with max odd rank $< i$. It is however possible that the same word is accepted by $\mathsf{KVF}(A)$ using a tight ranking with max odd rank $> i$. Figure 1 shows an example of such an automaton, where the word a^ω is accepted by $\mathsf{KVF}(A)$ using a tight ranking with max odd rank 5, as well as with a tight ranking with max odd rank 3. The same word is accepted by only the 3^{rd} slice of our automaton A'. This motivates the definition of *minimally inessential ranks*. Given an NBW A with n states, odd rank i ($1 \leq i \leq 2n-1$) is said to be minimally inessential if every word α that is accepted by $\mathsf{KVF}(A)$ using a tight ranking with max odd rank i is also accepted by $\mathsf{KVF}(A)$ using a tight ranking with max odd rank $j < i$. An odd rank that is not minimally inessential is called minimally essential. As the example in Figure 1 shows, neither the KVF-construction nor the KVFS-construction allows us to detect minimally essential ranks in a straightforward way. Specifically, although the 5^{th} slice of $\mathsf{KVF}(A)$ for this example accepts the word a^ω, 5 is not a minimally inessential rank. In order to determine if 5 is minimally inessential, we must isolate the 5^{th} slice of $\mathsf{KVF}(A)$, and then check whether the language accepted by this slice is a subset of the language accepted by $\mathsf{KVF}(A)$ sans the 5^{th} slice. This involves complementing $\mathsf{KVF}(A)$ sans the 5^{th} slice, requiring a significant blowup. In contrast, the properties of our automaton A' allow us to detect minimally (in)essential ranks efficiently. Specifically, if we find that the i^{th} slice of A' accepts a word α, we can infer than i is minimally essential. Once all minimally essential ranks have been identified in this manner, we can prune automaton A' to retain only those slices that correspond to minimally essential ranks. This gives us a way of eliminating redundant slices (and hence states) in $\mathsf{KVF}(A)$.

7 An implementation of our algorithm

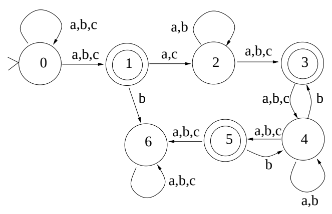

Fig. 2. Example automaton with gaps

We have implemented the complementation algorithm presented in this paper as a facility on top of the BDD-based model checker NuSMV. In our implementation, states of the complement automaton are encoded as pairs (g, i), as explained in Section 5.1. Our tool takes as input an NBW A, and generates the state transition relation of the complement automaton A' using the above encoding in NuSMV's input format. Generating the NuSMV file from a given description of NBW takes negligible time (< 0.01s). The number of boolean constraints used in expressing the transition relation in NuSMV is quadratic in n. We use NuSMV's fair CTL model checking capability to check whether there exists an infinite path in a slice of A' (corresponding to a maximum odd rank k) that visits an accepting state infinitely often. If NuSMV responds negatively to such a query, we disable all transitions to and from states in this section of the NuSMV model. This allows us to effectively detect and eliminate redundant slices of our complement automaton, resulting in a reduction of the overall size of the automaton. For purposes of comparison, we have also implemented the algorithm presented in [Sch09] in a similar manner using a translation to NuSMV. We also compared the performance of our technique with those of Safra-Piterman [Saf88,Pit07] determinization based complementation technique and Kupferman-Vardi's ranking based complementation technique [KV01], as implemented in the GOAL tool [TCT$^+$08]. Table 1 shows the results of some of our experiments. We used the CUDD BDD library with NuSMV 2.4.3, and all our experiments were run on an Intel Xeon 3GHz with 2 GB of memory and a timeout of 10 minutes. The entry for each automaton [1] lists the number of states, transitions and final states of the original automaton, the number of states of the complement automaton computed by the KVFS-construction and by our construction, and the set of minimally inessential ranks (denoted as "MI Ranks") identified by each of these techniques. In addition, each row also lists the number of states computed by the "Safra-Piterman" (denoted as "SP") technique and "Weak Alternating Automata" (denoted as "WAA") technique in GOAL.

A significant advantage of our construction is its ability to detect "gaps" in slices. As an example, the automaton in Figure (2) has ranks $1, 5$ as minimally inessential, while ranks 3 and 7 are minimally essential (see Table (1)). In this case, if we compute the *rank* of the NBW (as suggested in [GKSV03]) and then

[1] All example automata from this paper and the translator from automaton description to NuSMV are available at http://www.cfdvs.iitb.ac.in/reports/minrank

consider slices only upto this *rank*, we will fail to detect that rank 5 is minimally inessential. Therefore, eliminating redundant slices is a stronger optimization than identifying and eliminating states with ranks greater than the rank of the NBW.

8 Conclusion

In this paper, we presented a complementation algorithm for nondeterministic Büchi automata that is based on the idea of ranking functions introduced by Kupferman and Vardi[KV01]. We showed that the ranking assignment presented in [KV01] always results in a minimal odd ranking for run-DAGs of words in the complement language. We then described a complementation construction for NBW such that the complement NBW accepts only the run-DAG with the minimal odd ranking for every word in the complement. We observed that the states of the complement NBW are partioned into *slices*, and that each word in the complement is accepted by exactly one such slice. This allowed us to check for redundant slices and eliminate them, leading to a reduction the size of the complement NBW. It is noteworthy that this ability to reduce the size of the final complement NBW comes for free since the worst case bounds coincide with the worst case bounds of the best known NBW complementation construction. In the future, we wish to explore techniques to construct unambiguous automata and deterministic Rabin automata from NBW, building on the results presented here.

Automaton	KVFS algorithm		Our algorithm		GOAL	
(states,trans.,final)	States	MI Ranks	States	MI Ranks	WAA	SP
michel4(6,15,1)	157	{9,7}	117	{9,7}	XX	105
g15 (6,17,1)	324	{9,7,5}	155	{9,7,5}	XX	39
g47 (6,13,3)	41	{5}	29	{5,3}	XX	28
ex4 (8,9,4)	3956	∅	39	{7,5}	XX	7
ex16 (9,13,5)	10302	∅	909	{7}	XX	30
ex18 (9,13,5)	63605	∅	2886	∅	XX	71
ex20 (9,12,5)	17405	∅	156	{7}	XX	24
ex22 (11,18,7)	258	{7,5}	23	{7,5}	XX	6
ex24 (13,16,8)	4141300	∅	19902	{9,7}	XX	51
ex26 (15,22,11)	1042840	∅	57540	{7,5}	XX	XX
gap1 (15,22,11)	99	{5,1}	26	{5,1}	XX	16
gap2 (15,22,11)	532	{1}	80	{5,1}	XX	48

Table 1. Experimental Results. XX:timeout (after 10 min)

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