

Mechanism Design (Inverse Game Theory)

The objectives/desired outcomes are set - task is to set the rules of the game

E.g., Election, license scarce resource (spectrum, cloud), matching students to universities

General model:

N : set of players

X : set of outcomes, e.g., winner in an election, which resource allocated to whom etc.

Θ_i : set of private information of agent i (type). A type $\theta_i \in \Theta_i$

The type may manifest in the preferences over the outcomes in different ways

① Ordinal: θ_i defines an ordering over the outcomes

② Cardinal: an utility function u_i maps an (outcome, type) pair to real numbers, $u_i: X \times \Theta_i \rightarrow \mathbb{R}$ (private value model)
or $u_i: X \times \Theta \rightarrow \mathbb{R}$ (interdependent value model)

Examples: Voting: X is the set of candidates

θ_i is a ranking over these candidates, e.g., $\theta_i = (a, b, c)$, i.e., a is more preferred than b which in turn is more preferred than c .

Single object allocation: an outcome is $x = (\underline{a}, \underline{p}) \in X$

$\underline{a} = (a_1, a_2, \dots, a_n)$, $a_i \in \{0, 1\}$, $\sum_{i \in N} a_i \leq 1$, allocations

$\underline{p} = (p_1, p_2, \dots, p_n)$, p_i is the payment charged to i

θ_i : value of i for the object

$$u_i(x, \theta_i) = a_i \theta_i - p_i$$

But The designer has an objective

This is captured through a Social Choice Function (SCF)

$$f : \Theta_1 \times \Theta_2 \times \dots \times \Theta_n \rightarrow X$$

E.g., in voting, if there is a candidate who beats everyone else in pairwise contests must be chosen as a winner.

in public project choice, where $\theta_i : X \rightarrow \mathbb{R}$, value for each project pick $f(\theta) \in \operatorname{argmax}_{a \in X} \sum_{i \in N} \theta_i(a)$.

Q: How can we create a game where $f(\theta)$ emerges as an outcome of an equilibrium?

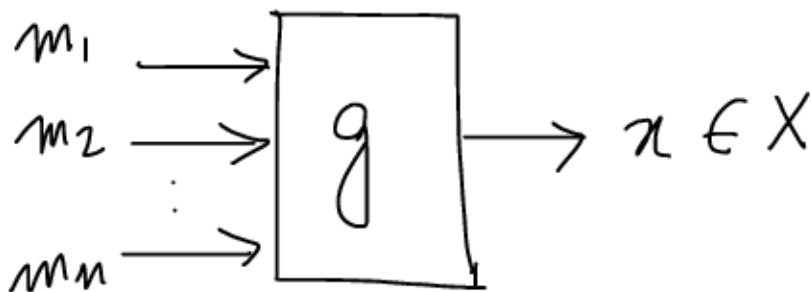
A: we need mechanisms.

Defn. An (indirect) mechanism is a collection of message spaces and a decision rule $\langle M_1, M_2, \dots, M_n, g \rangle$

- M_i is the message space of agent i

- $g : M_1 \times M_2 \times \dots \times M_n \rightarrow X$

A direct mechanism is same as above with $M_i = \Theta_i, \forall i \in N, g \equiv f$.



The message space is similar to equipping every agent with a card deck and asking to pick some

Q: Why these are not so commonplace?

A: due to a result that will follow.

Defn. In a mechanism $\langle M_1, \dots, M_n, g \rangle$, a message m_i is **weakly dominant** for player i at θ_i if

$$u_i(g(m_i, \tilde{m}_{-i}), \theta_i) \geq u_i(g(m'_i, \tilde{m}_{-i}), \theta_i), \forall \tilde{m}_{-i}, \forall m'_i$$

[all subsequent definitions assume cardinal preferences, however they can be replaced with ordinal, e.g., the above one could be defined as

$$g(m_i, \tilde{m}_{-i}) \theta_i g(m'_i, \tilde{m}_{-i}) \forall m'_i, \forall \tilde{m}_{-i} \quad]$$

↑ this outcome is preferred at least as much as the latter

Defn. An SCF $f: \Theta \rightarrow X$ is **implemented** in dominant strategies by $\langle M_1, \dots, M_n, g \rangle$ if

- ① \exists message mappings $s_i: \Theta_i \rightarrow M_i$, s.t., $s_i(\theta_i)$ is a dominant strategy for agent i at θ_i , $\forall \theta_i \in \Theta_i$, $\forall i \in N$, and
- ② $g(s_1(\theta_1), \dots, s_n(\theta_n)) = f(\theta)$, $\forall \theta \in \Theta$.

We call this an **indirect implementation**, i.e., SCF f is **dominant strategy implementable (DSI)** by $\langle M_1, \dots, M_n, g \rangle$.

Defn. A direct mechanism $\langle \Theta_1, \dots, \Theta_n, f \rangle$ is **dominant strategy incentive compatible (DSIC)** if

$$u_i(f(\theta_i, \tilde{\theta}_{-i}), \theta_i) \geq u_i(f(\theta'_i, \tilde{\theta}_{-i}), \theta_i), \forall \theta_i, \theta'_i, \tilde{\theta}_{-i} \\ \forall i \in N.$$

To find if an SCF f is dominant strategy implementable, we need to search over all possible indirect mechanisms $\langle M_1, \dots, M_n, g \rangle$. But luckily, there is a result that reduces the search space.