



भारतीय प्रौद्योगिकी संस्थान मुंबई  
Indian Institute of Technology Bombay

# CS 6001: Game Theory and Algorithmic Mechanism Design

Week 5

Swaprava Nath

Slide preparation acknowledgments: Ramsundar Anandanarayanan and Harshvardhan Agarwal

ज्ञानम् परमम् ध्येयम्

Knowledge is the supreme goal



- ▶ Imperfect Information Extensive Form Games

- ▶ Strategies in IIEFGs

- ▶ Equivalence of strategies in IIEFGs

- ▶ Perfect Recall



## The story so far

- Games discussed so far (EFGs) are of perfect information

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<sup>a</sup><https://rbc.jhuapl.edu/>



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- Limited use in certain setups:
  - several games have states that are unknown to certain agents, e.g., card games like poker, reconnaissance blind chess<sup>a</sup>
  - not possible to represent simultaneous move games using EFGs

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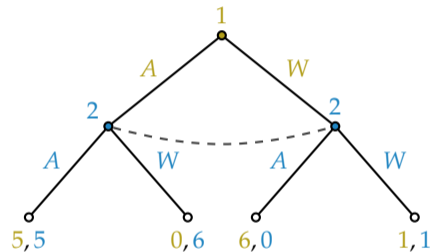
# Games with Imperfect Information



Kingdom 1

		Kingdom 2	
		Agri	War
Kingdom 1	Agri	5,5	0,6
	War	6,0	1,1

Neighboring Kingdom's Dilemma



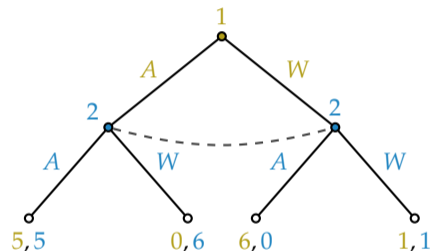


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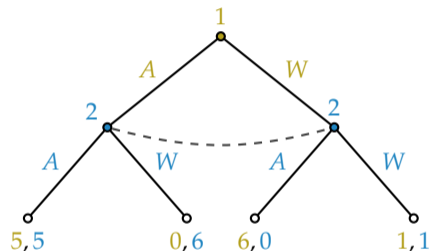
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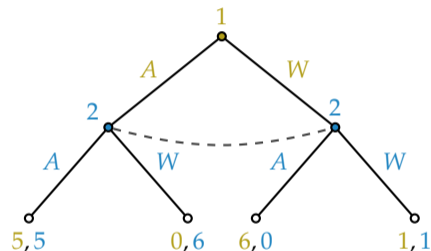
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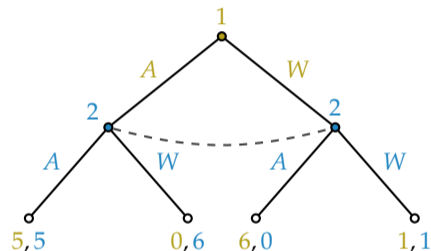
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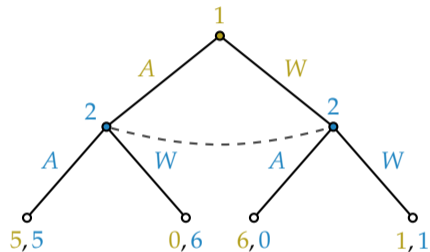
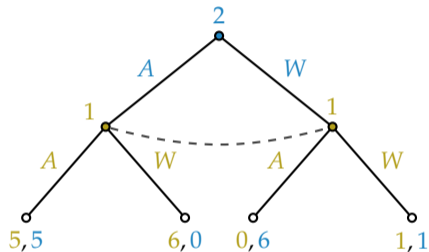
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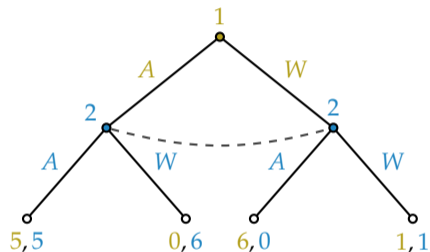
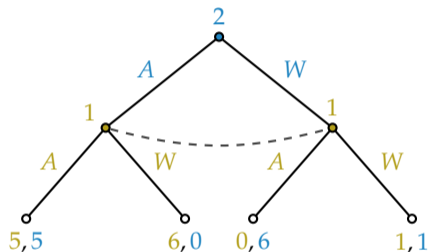


- In IIEFG, indistinguishable nodes are connected via a dotted line.
- Kingdom 2 does not know which node/history the game is in
- These indistinguishable histories form an **information set** for player 2.
- More general representation than PIEFG since information sets can be singleton

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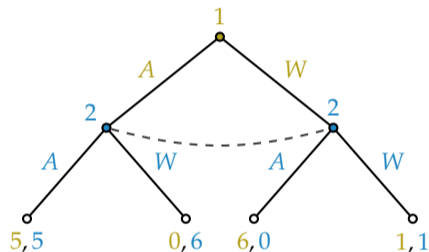
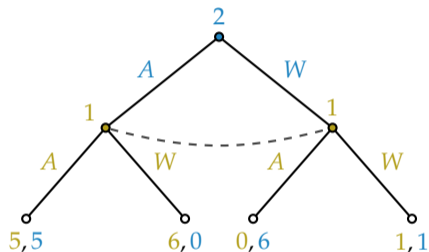


# Games with Imperfect Information



- The Neighboring Kingdom's dilemma can also be represented with the information set of player 1 being non-singleton.

# Games with Imperfect Information



- The Neighboring Kingdom's dilemma can also be represented with the information set of player 1 being non-singleton.
- IIEFGs are not unique for a given simultaneous move game

# Games with Imperfect Information



## Definition (IIEFG)

An IIEFG is tuple  $\langle N, A, H, X, P, (u_i)_{i \in N}, (I_i)_{i \in N} \rangle$



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- $I_i^j$ s are called an **information set** of player  $i$  and  $I_i$  is the collection of information sets of  $i$ .
- At an information set, the player and her available actions are the same.
- The player is uncertain about which history in the information set is reached.

# Games with Imperfect Information (contd.)



## Definition (IIIEFG)

An IIIEFG is tuple  $\langle N, A, H, X, P, (u_i)_{i \in N}, (I_i)_{i \in N} \rangle$  where  $\langle N, A, H, X, P, (u_i)_{i \in N} \rangle$  is a PIEFG and for every  $i \in N$ ,  $I_i = (I_i^1, I_i^2, I_i^3, \dots, I_i^{k(i)})$  is a partition of  $\{h \in H \setminus Z : P(h) = i\}$  with the property that  $X(h) = X(h')$  and  $P(h) = P(h') = i$ , whenever  $\exists j$  s.t.  $h, h' \in I_i^j$ .

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- Some differences with PIEFG

- Since actions at an information set are identical,  $X$  (action set function) can be defined over  $I_i^j$ s i.e.,  
$$X(h) = X(h') = X(I_i^j), \forall h, h' \in I_i^j$$

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- Strategies can also be defined over information sets, i.e., strategy set of a player  $i \in N$  is defined as the Cartesian product of actions available to  $i$  at her information sets

$$S_i = \times_{I' \in I_i} X(I') = \times_{j=1}^{j=k(i)} X(I_i^j)$$



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- IIIEFG is a richer representation than both NFG and PIEFG.



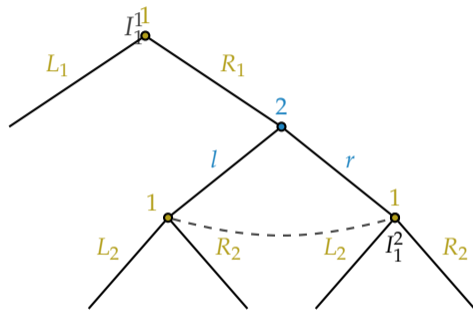
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# Randomized Strategies in IIEFGs



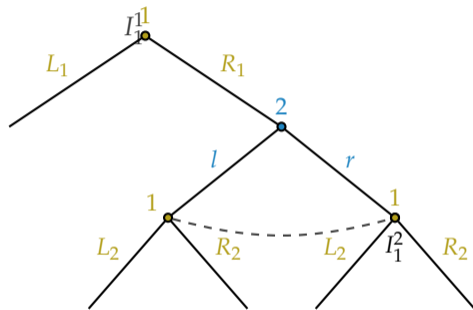
- Strategy set of  $i$ :  $S_i = \times_{j=1}^{j=k(i)} X(I_i^j)$
- In NFGs, mixed strategies randomize over pure strategies
- In EFGs, randomization can happen in different ways
  - randomize over the strategies defined at the beginning of the game
  - randomize over the action at an information set: **behavioral strategy**

# Randomized Strategies in IIEFGs



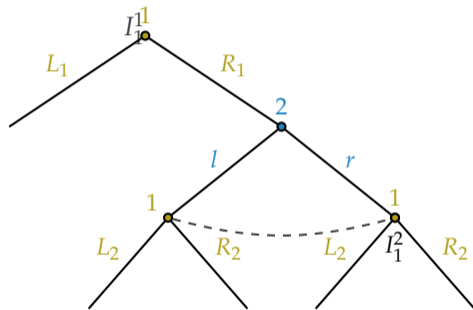
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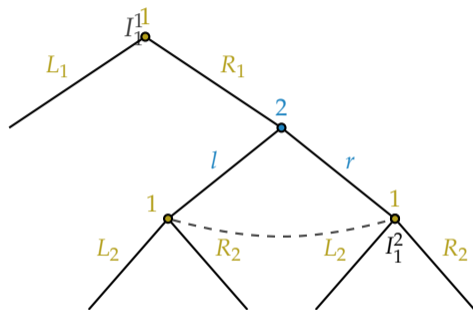
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- Mixed Strategy  $\sigma_1, \sigma_1(L_1L_2), \sigma(L_1R_2), \sigma(R_1L_2), \sigma(R_1R_2)$ .
- Behavioral Strategy  $b_1, b_1(I_1) \in \Delta(L_1, R_1), b_2(I_1) \in \Delta(L_1, R_1)$





## Definition

A behavioral strategy of a player in an IIEFG is a function that maps each of her information sets to a probability distribution over the set of possible actions **at that information set**.



## Question

What is the relation between mixed and behavioral strategies?



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- In this example, MSs live in  $\mathbb{R}^4$ , BSs live in two  $\mathbb{R}^2$  spaces
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# Mixed and Behavioral strategy



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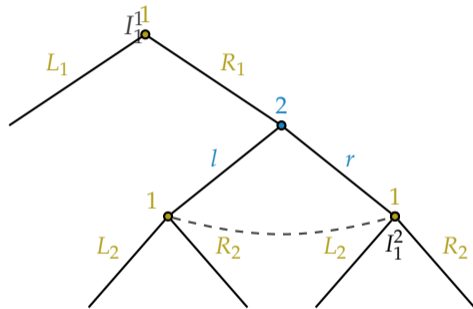
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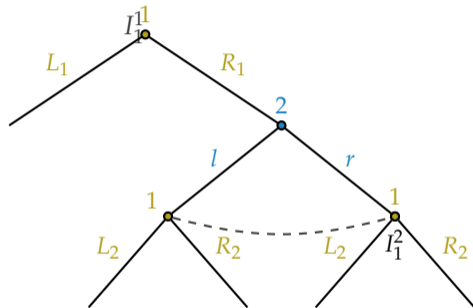
Equivalence in terms of the probability of reaching a vertex/history  $x$

- Say  $\rho(x; \sigma)$  is the probability of reaching a node  $x$  under mixed strategy profile  $\sigma$
- Similarly,  $\rho(x; b)$  is the same for behavioral strategy profile  $b$

# Example



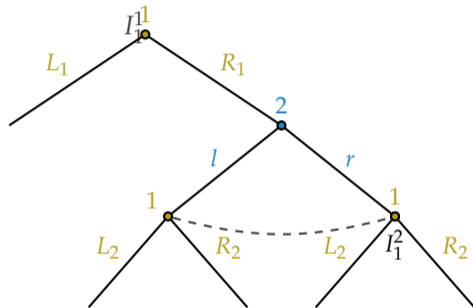
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$$\rho(x; \sigma) = \sigma_1(R_1)\sigma_2(r)$$

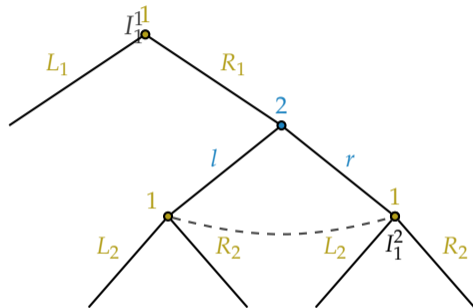


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$$\begin{aligned}\rho(x; \sigma) &= \sigma_1(R_1)\sigma_2(r) \\ &= (\sigma_1(R_1L_2) + \sigma_1(R_1R_2)) \cdot \sigma_2(r)\end{aligned}$$

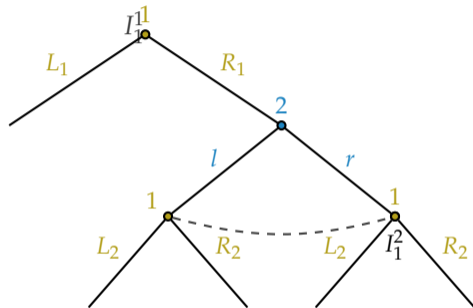
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Players can choose different kind of strategies

$$\rho(x; \sigma_1, b_2) = (\sigma_1(R_1L_2) + \sigma_1(R_1R_2)) \cdot b_2(I_2^1)(r)$$

# Equivalence Definition



## Definition

A mixed strategy  $\sigma_i$  and a behavioural strategy  $b_i$  of a player  $i$  in an IIEFG are **equivalent** if for every mixed/behavioral strategy  $\xi_{-i}$  of the other players and every vertex  $x$  in the game tree.

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## Example (in the game above)

Equivalent strategies induce same probability of reaching a vertex.

$$b_1(I_1^1)(L_1) = \sigma_1(L_1L_2) + \sigma_1(L_1R_2)$$

$$b_1(I_1^1)(R_1) = \sigma_1(R_1L_2) + \sigma_1(R_1R_2)$$

$$b_1(I_1^2)(L_2) = \sigma_1(L_2|R_1)$$

$$b_1(I_1^2)(R_2) = \sigma_1(R_2|R_1)$$

We call  $b_1$  and  $\sigma_1$  are equivalent.

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## Claim

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**Reason:** Pick an arbitrary non-leaf node, the probability of reaching that node is equal to the sum of the probabilities of reaching the leaf nodes in its subtree.



# More on Equivalent Strategies



This argument can be extended further

## Theorem (Utility Equivalence)

*If  $\sigma_i$  and  $b_i$  are equivalent, then for every mixed/behavioural strategy vector of the other players  $\xi_{-i}$ , the following holds,*

$$u_j(\sigma_i, \xi_{-i}) = u_j(b_i, \xi_{-i}) \forall j \in N.$$

Repeat the argument for any equivalent mixed and behavioral strategy profiles.

## Corollary

Let  $\sigma$  and  $b$  be equivalent, i.e.,  $\sigma_i$  and  $b_i$  are equivalent  $\forall i \in N$ , then  $u_i(\sigma) = u_i(b)$ .



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# Equivalence of strategies in IIEFGs



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## Answer

- More natural in large IIEFGs

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  - Consider a player having 4 information sets with 2 actions each

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# Equivalence of strategies in IIEFGs



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Why behavioral strategies are desirable?

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Can we construct one from another?

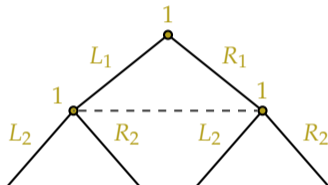
OR

Does equivalence always hold?

# Equivalence of strategies in IIEFGs (Example 1)



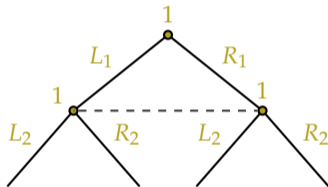
Player remembers that it made a move but forgets which move



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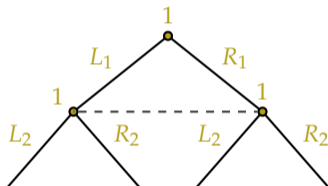


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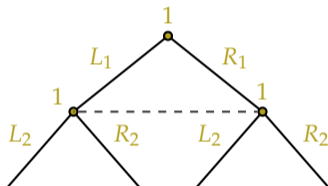


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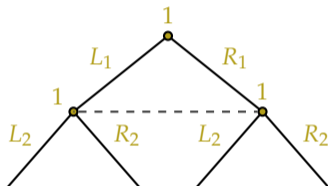


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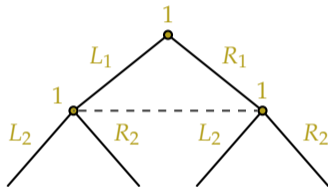


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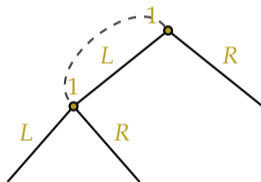
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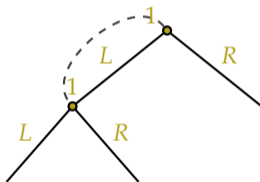
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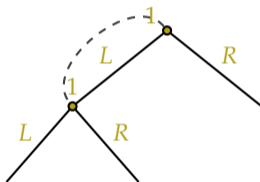


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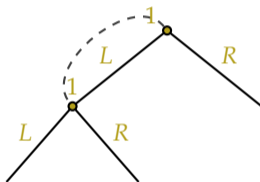
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Answer

The equivalence does not hold if the players are forgetful

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When does behavioral strategy have no equivalent mixed strategy?

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- 5 Since mixed strategy is a randomization over pure strategies, every mixed strategy will put zero probability mass on  $x$  but behavioral strategy randomizes on every vertex **independently**, hence  $x$  may be reached in behavioral strategies with a positive probability

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## Theorem (6.11 of MSZ)

*Consider an IIEFG such that every vertex has at least two actions. Every behavioral strategy has an equivalent mixed strategy for a player iff each information set of that player intersects every path emanating from the root at most once.*

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## Proof.

Homework. Reading exercise from MSZ.





- ▶ Imperfect Information Extensive Form Games
- ▶ Strategies in IIEFGs
- ▶ Equivalence of strategies in IIEFGs
- ▶ Perfect Recall



# Behavioral Strategy equivalent to Mixed Strategy



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To formalize (i.e., to set the conditions when the equivalence holds), we need to formalize the **forgetfulness** of a player.

- saw few examples of players' forgetfulness.
- our conditions need to ensure that none of the previous types of forgetfulness happens.

# Behavioral Strategy equivalent to Mixed Strategy



Definition (Choice of **same action at an information set**)

Let  $X = (x^0, x^1, \dots, x^K)$  and  $\hat{X} = (x_0, \hat{x}^1, \dots, \hat{x}^L)$  be two paths in the game tree.

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**'Leading to'** may not be a relation between parent and child nodes, it can be any descendant of the former since the path is unique in a tree.





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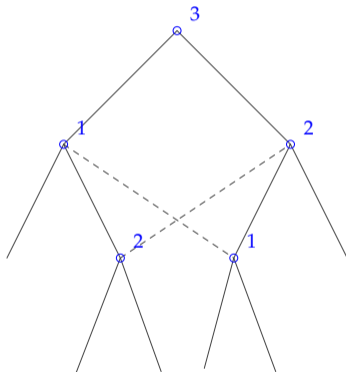


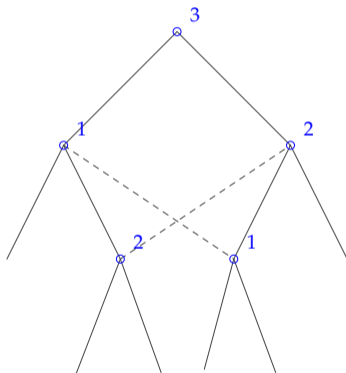
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Note: Definition of perfect recall subsumes the condition where every behavioral strategy has equivalent mixed strategy

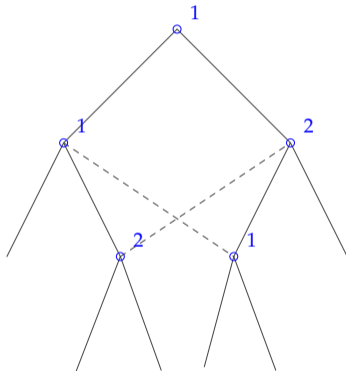
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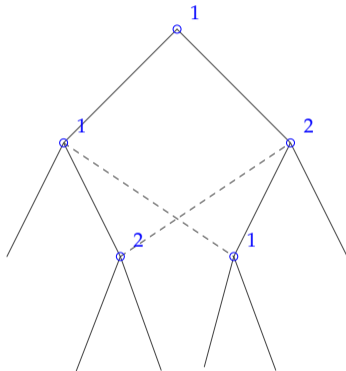




**Game with Perfect Recall:** This example satisfies the conditions of the definitions.

# Examples





**Game with Imperfect Recall:** Player 1 takes two different actions at the first information set to reach two different vertices of the second information set.

# Implications of Perfect Recall



Let  $S_i^*(x)$  be the set of pure strategies of player  $i$  at which he chooses actions leading to  $x$ , i.e., intersections of members of  $S_i$  with the path from root to  $x$ .

## Theorem

*If  $i$  is a player with perfect recall and  $x$  and  $x'$  are the two vertices in the same information set of  $i$ , then  $S_i^*(x) = S_i^*(x')$ .*

The above conclusion comes from the same sequence of information sets and same actions. The next implication of mixed and behavioral strategies.

# Implications of Perfect Recall



## Theorem (Kuhn 1957)

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- The proof is constructive. It starts with the mixed strategy and constructs the behavioral strategies such that the probabilities of reaching a leaf are same. The arguments show that such a construction is always possible because of perfect recall.



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